

Compiler Practical 2013

Methods with Parameters

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Cartesium 2.48



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1. Method Call and Stack
2. Extending Methods with Parameters
3. Bonus Task: Evaluation of Constant Sub-expressions

- Syntax analysis
 - Extend the grammar
 - *Extend MethodDeclaration and VarOrCall*
- Context analysis
 - For formal and actual parameters
 - Pairwise comparison of formal and actual parameter types
 - Assigning relative addresses in the stack frame

Methods with Parameters (2)

- Synthesis
 - Evaluate expressions of actual parameters one by one
 - Results stay on the stack
 - Call method
 - Remove parameters from stack at end of method execution

memberdecl ::=

vardecl ';'

| METHOD *identifier* ['(' *vardecl* { ';' *vardecl* } ')']

IS *methodbody*

varorcall ::=

identifier ['(' *disjunction* { ',' *disjunction* } ')']

- Parameter passing corresponds to an assignment

- Actual to formal parameter

```
METHOD proc(a, b: Integer) IS
```

isA ↑ ↑ isA

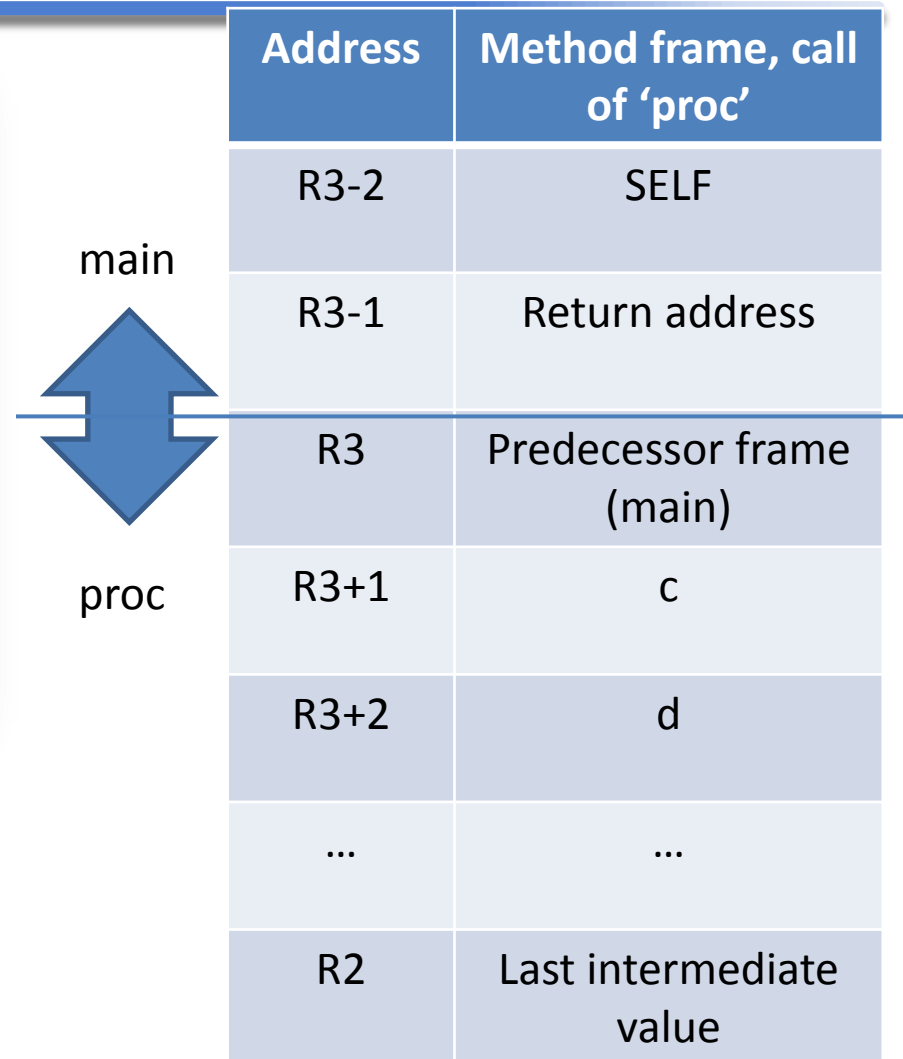
```
proc(5, 7);
```

- Actual parameters are references
 - *Boxing or Dereferencing*, if needed
- Numbers of actual and formal parameters must match
- Types of actual parameters must be compatible to those of formal parameters (*isA*)

Stack Frames without Parameters

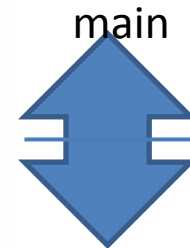
```

CLASS Main
  METHOD proc IS
    c, d: Integer;
  BEGIN END METHOD
  METHOD main IS BEGIN
    proc; | Aufruf von 'proc'
  END METHOD
END CLASS
    
```



Stack Frames with Parameters

```
CLASS Main
METHOD proc(a,b: Integer) IS
  c, d: Integer;
BEGIN END METHOD
METHOD main IS BEGIN
  proc(8, 15); | Aufruf von 'proc'
END METHOD
END CLASS
```



Address	Method frame, call of 'proc'
R3-4	SELF
R3-3	a
R3-2	b
R3-1	Return address
R3	Predecessor frame (main)
R3+1	c
R3+2	d
...	...
R2	Last intermediate value

10%

- Evaluating constant sub-expressions at compile time (runtime reduction)
 - After context analysis, before code generation
- Expressions
 - Attempt to evaluate the operands of an operator
 - If all operands of an operator are literals, evaluate the operator and replace the expression by the result (a literal)
 - Apply transformations of expressions otherwise
 - This concerns *UnaryExpression* and *BinaryExpression*

5%

- This works with some statements as well:
 - IF TRUE THEN s ELSE $t \rightarrow s$
 - IF FALSE THEN s ELSE $t \rightarrow t$
 - WHILE TRUE DO $s \rightarrow$ FOREVER s
 - WHILE FALSE DO $s \rightarrow$ (nothing)

Transformations: +, -

- $- -x \rightarrow x$
- $0 + x \rightarrow x$
- $x + 0 \rightarrow x$
- $x + -y \rightarrow x - y$
- $c \pm x \pm .. \pm y \pm d \pm .. \rightarrow (c \pm d) \pm x \pm .. \pm y \pm ..$
- $0 - x \rightarrow -x$
- $x - 0 \rightarrow x$
- $x - -y \rightarrow x + y$

Transformations: Multiplikation

- $0 * x \rightarrow 0$
- $x * 0 \rightarrow 0$
- $1 * x \rightarrow x$
- $x * 1 \rightarrow x$
- $c * x * .. * y * d * .. \rightarrow (c * d) * x * .. * y * ..$
- $c * -x \rightarrow (-c) * x$
- $-x * c \rightarrow (-c) * x$
- $-x * y \rightarrow -(x * y)$
- $x * -y \rightarrow -(x * y)$
- $-x * -y \rightarrow x * y$

Transformationen: Division

- $0 / x \rightarrow 0$, error if $x = 0$
- $x / 1 \rightarrow x$
- $c / x / .. / y / d / .. \rightarrow (c / d) / x / .. / y / ..$
- $x / c / .. / y / d / .. \rightarrow x / (c * d) / .. / y / ..$
- $-x / c \rightarrow x / (-c)$
- $-x / y \rightarrow -(x / y)$
- $x / -y \rightarrow -(x / y)$
- $-x / -y \rightarrow x / y$

Transformations: AND, OR, NOT

- NOT NOT x $\rightarrow x$
- FALSE AND x \rightarrow FALSE (error in x is ignored!)
- x AND FALSE \rightarrow FALSE (error in x is ignored!)
- TRUE AND x $\rightarrow x$
- x AND TRUE $\rightarrow x$
- FALSE OR x $\rightarrow x$
- x OR FALSE $\rightarrow x$
- TRUE OR x \rightarrow TRUE (error in x is ignored!)
- x OR TRUE \rightarrow TRUE (error in x is ignored!)