

# Theorem Proving in Isabelle

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Try this in Isabelle

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i.e.

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'Paper' proof:

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(first two steps by  $\forall E$ )

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Check it out!

# Lifting

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What about resolution with

$$(\rightarrow I) \frac{P \implies Q}{P \rightarrow Q} \quad \text{or} \quad (\forall I) \frac{\bigwedge x. P}{\forall x. P} \quad ?$$

Problem: The premises contain meta-logical symbols ( $\implies$ ,  $\bigwedge$ ), hence do not match conclusion of any other rule!

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$$\frac{\phi \implies \psi}{(\theta \implies \phi) \implies (\theta \implies \psi)}$$

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**Resolution with  $\rightarrow I$** , i.e. with  $(?P \Longrightarrow ?Q) \Longrightarrow ?P \rightarrow ?Q$ :

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- Introduce quantification over **parameter**  $x$  in premises and conclusion
- Make all unknowns  $?a$  depend on  $x$ : replace by  $?a(x)$ .
- Meta-rule:

$$\frac{\phi \implies \psi}{\bigwedge x. \phi^x \implies \bigwedge x. \psi^x}$$

( $\phi^x$  is  $\phi$  with parametrized unknowns)

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Want to resolve  $\wedge E1$ , i.e.  $?P \wedge ?Q \implies ?P$ , with  $\forall I$ , i.e.  $\bigwedge x. ?P(x) \implies \forall x. ?P(x)$ .

Lift  $\wedge E1$  over parameter  $x$ :

$$\bigwedge x. ?P(x) \wedge ?Q(x) \implies \bigwedge x. ?P(x)$$

Resolve with  $\forall I$ :

$$\bigwedge x. ?P(x) \wedge ?Q(x) \implies \forall x. ?P(x)$$