

Logik für Informatiker

Logic for computer scientists

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LPL book detailed introduction into first-order logic
with many exercises

Boole construct truth tables

Tarski's world evaluate logical formulas within a blocks world

Fitch construct proofs

Grinder gives automatic feedback to your solutions
→ requires purchase of the CD (ca. 13 EUR) or the
book (ca. 25 EUR, with CD)

Platform for exercises: `logic.informatik.uni-bremen.de`
also reachable via
`www.informatik.uni-bremen.de/agbkb/lehre/ws11-12/Logik/`

The formal language PL1

PL1 is the formal language of **first-order predicate logic**

Why do we need a formal language?

⇒ Slides from Prof. Barbara König, Universität Duisburg-Essen

<http://jordan.inf.uni-due.de/teaching/ss2010/logik/folien/folien.pdf>

The language of PL1: individual constants

- **Individual constants** are symbols that denote a person, thing, object
- Examples:
 - Numbers: 0, 1, 2, 3, ...
 - Names: Max, Claire
 - Formal constants: a, b, c, d, e, f, n1, n2
- Each individual constant must denote an existing object
- No individual constant can denote more than one object
- An object can have 0, 1, 2, 3 ... names

The language of PL1: predicate symbols

- **Predicate symbols** denote a property of objects, or a relation between objects
- Each predicate symbol has an **arity** that tell us how many objects are related
- Examples:
 - Arity 0: Gate0_is_low, A, B, ...
 - Arity 1: Cube, Tet, Dodec, Small, Medium, Large
 - Arity 2: Smaller, Larger, LeftOf, BackOf, SameSize, Adjoins
 - ...
 - Arity 3: Between

The interpretation of predicate symbols

- In **Tarski's world**, predicate symbols have a **fixed interpretation**, that not always completely coincides with the natural language interpretation
- In other PL1 languages, the interpretation of predicate symbols may **vary**. For example, \leq may be an ordering of numbers, strings, trees etc.
- Usually, the binary symbol $=$ has a fixed interpretation: **equality**

Atomic sentences

- in propositional logic (Boole):
 - propositional symbols: a, b, c, \dots
- in PL1 (Tarski's world):
 - application of predicate symbols to constants: $\text{Larger}(a,b)$
 - the **order** of arguments **matters**: $\text{Larger}(a,b)$ vs. $\text{Larger}(b,a)$
 - Atomic sentences denote **truth values** (true, false)

- **Function symbols** lead to more complex **terms** that denote objects. Examples:
 - father, mother
 - +, -, *, /
- This leads to new terms denoting objects:
 - father(max) mother(father(max))
 - $3*(4+2)$
- This also leads to new atomic sentences:
 - Larger(father(max),max)
 - $2 < 3*(4+2)$

A sentence A is a **logically valid**, if it is true in all circumstances.

A sentence A is a **satisfiable**, if it is true in at least one circumstance.

A **circumstance** is

- in propositional logic: a valuation of the atomic formulas in the set $\{ \text{true}, \text{false} \}$
- in Tarski's world: a block world

Consequences ...



A sentence B is a **logical consequence** of A_1, \dots, A_n , if all circumstances that make A_1, \dots, A_n true also make B true.

In symbols: $A_1, \dots, A_n \models B$.

A_1, \dots, A_n are called **premises**, B is called **conclusion**.

In this case, it is a **valid argument** to infer B from A_1, \dots, A_n . If also A_1, \dots, A_n are true, then the valid argument is **sound**.

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Logical consequence — examples

- All men are mortal. Socrates is a man. So, Socrates is mortal. (valid, sound)
- All rich actors are good actors. Brad Pitt is a rich actor. So he must be a good actor. (valid, but not sound)
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