## Homework 2

Finite Automata on Infinite Words and Trees Winter semester, 2009-2010

Note: Graphical representations of automata are accepted.

**Exercise 1** (25 points). Consider the alphabet  $\Sigma = \{0, 1\}$ . Construct Büchi automata which accept each of the following languages:

- 1.  $\{w \mid \text{the symbol 0 appears in } w \text{ exactly twice}\}$
- 2.  $\{w \mid \text{every } 0 \text{ which appears in } w \text{ is followed immediately by } 11\}$
- 3.  $\{w \mid w \text{ does not contain the substring } 000 \}$
- 4.  $\{w \mid w \text{ contains finitely many substrings } 11\}$
- 5.  $\{w \mid w \text{ contains finitely many substrings 11 but infinitely many 1's} \}$

Exercise 2 (25 points). Prove or disprove the following two statements:

- 1. The set of languages recognized by deterministic Büchi automata is closed under union.
- 2. The set of languages recognized by deterministic Büchi automata is closed under intersection.

**Exercise 3** (25 points). Suppose that L is a Büchi recognizable language over the alphabet

$$\Sigma = \Sigma_1 \times \Sigma_2 = \{(a, b) \mid a \in \Sigma_1, b \in \Sigma_2\}$$

Prove that the languages  $PR_1(L)$  and  $PR_2(L)$  defined by

 $PR_{1} = \{ u \in \Sigma_{1}^{\omega} \mid \exists v \in \Sigma_{2}^{\omega} \text{ such that } (u[0], v[0])(u[1], v[1])(u[2], v[2]) \dots \in L \}$   $PR_{2} = \{ v \in \Sigma_{2}^{\omega} \mid \exists u \in \Sigma_{1}^{\omega} \text{ such that } (u[0], v[0])(u[1], v[1])(u[2], v[2]) \dots \in L \}$ are both Büchi recognizable.

**Exercise 4** (25 points). Given a Büchi recognizable language SELECT over alphabet  $\{1, 2\}$  and two Büchi recognizable languages  $L_1$ ,  $L_2$  over alphabet  $\{a, b\}$ , prove that the following language FUSION is also Büchi recognizable:

FUSION = {
$$\sigma$$
 | there exists  $\alpha \in L_1$  and  $\beta \in L_2$  and  $\gamma \in \text{SELECT}$  such that  $\sigma[i] = \alpha[i]$  if  $\gamma[i] = 1$ , and  $\sigma[i] = \beta[i]$  if  $\gamma[i] = 2$ }