Modularity in Ontologies: Introduction (Part B)

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Plan for the rest of today



A case for modularity of ontologies



Overview and comparison of modularisation approaches



Overview of the remainder of this course



And now ...



A case for modularity of ontologies

Overview and comparison of modularisation approaches





Ontology users and engineers want to use ontologies to

- represent and archive knowledge
- compute inferences from archived knowledge e.g., classification, query answering
- explain inferences justifications = pinpointing, abduction
- reuse knowledge to build other ontologies import
- impose the logical structure of the represented knowledge comprehension



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- compute inferences from archived knowledge (M) e.g., classification, query answering
- explain inferences (M) justifications = pinpointing, abduction
- reuse knowledge to build other ontologies (M) import
- impose the logical structure of the represented knowledge (M) comprehension

Building and using an ontology often requires

• fast reasoning

expressivity \leftrightarrow complexity, optimisations, incremental reasoning

- collaborative development
- version control
- efficient reuse
- an understanding of the ontology's content and structure comprehension



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- an understanding of the ontology's content and structure (M) comprehension

(M) = modularity helps



An import/reuse scenario

"Borrow" knowledge from external ontologies



- Provides access to well-established knowledge
- Doesn't require expertise in external disciplines

This scenario is well-understood and implemented.

➤ Wednesday's lecture



A collaboration scenario

Collective ontology development



- Developers work (edit, classify) locally
- Extra care at re-combination
- Prescriptive/analytic behaviour

This approach is mostly understood, but not implemented yet.



Understanding and/or structuring an ontology

Compute the modular structure of an ontology





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This is work in progress. → Friday's lecture



And now ...



Overview and comparison of modularisation approaches



A priori vs. a posteriori

A priori

- At first, a modular structure is decided on.
- Then, the ontology is developed and used according to that structure.

A posteriori

- After the ontology has been built, a module is extracted or the ont. is decomposed into modules.
- The ontology is regarded as a monolithic entity.



A-priori modularisation approaches

- Provide a framework to develop an ontology modularly from the start
- Provide means to "bridge" between the modules dependency of modules/signature, flow of knowledge
- Often consist of extensions of (description) logics
- Sometimes allow for distributed reasoning
- Generally, don't guarantee that modules are logically closed in some cases, this is deliberately so



A case for modularity

A-priori: different files with imports

- Used to develop large ontologies about different domains
- Each domain expert (team) maintains "their" file \mathcal{F}_i
- The overall ontology \mathcal{O} imports all files: $\mathcal{O} = \mathcal{F}_1 \cup \cdots \cup \mathcal{F}_n$
- \bullet Example: $\mathcal{F}_1, \mathcal{F}_2, \mathcal{F}_3$ about diseases, anatomy and drugs
- Problems?



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- \bullet Example: $\mathcal{F}_1, \mathcal{F}_2, \mathcal{F}_3$ about diseases, anatomy and drugs
- Problems?
 - The \mathcal{F}_i are not necessarily logically closed
 - Experts' knowledge interferes with each other, e.g.: diseases are located in body parts and treated by drugs
 - $\rightsquigarrow\,$ Maintenance of ${\cal O}$ as difficult as in the monolithic case
 - Reasoning or reuse might still require the whole ontology
- Still used to develop and maintain, e.g., some bio-medical ontologies!



Package-based description logics (PB-DLs)

[Bao et al. 2006, 2009]

- Extension of standard DLs
- Domain-specific files are called packages
- Semantic import links between packages (explicit dependency)
- Terms annotated with "home package"
- Semantics local w.r.t. each package
- Reasoning controlled by the links
- Translation to "plain" DLs yields implicit decision procedures
- Problems?
 - Reasoning or reuse might still require the whole ontology



Distributed description logics

[Borgida and Serafini, 2003] [Serafini and Tamilin, 2009]

- Similar to PB-DLs
- Replace import links by "bridge rules": subconcept relations between (complex) concepts from different packages
- Distributed decision procedures exist

Related notion: E-connections [Kutz et al. 2001]



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A-posteriori modularisation approaches

- Regard an ontology \mathcal{O} as a monolithic entity remember: \mathcal{O} is a set of axioms
- Module: subset $\mathcal{M} \subseteq \mathcal{O}$
- Extract one module (e.g., for reuse) or decompose O into several modules (e.g., for comprehension)
- Often, a signature (set of terms) Σ ⊆ sig(O) is specified and the module extracted using Σ as a parameter
- Ideally, modules encapsulate knowledge in some form e.g., all consequences of ${\cal O}$ in Σ
- Not all module notions guarantee encapsulation



Graph-based a-posteriori modularisation approaches

- Are based on a graph representation of the ontology usually concept/role hierarchy, sometimes enriched with disjointness
- Start with a signature $\boldsymbol{\Sigma}$
- Traverse the graph and "harvest" entities and axioms follow subconcept relation and/or restrictions (∃, domain, range)
- Resulting module = set of harvested axioms
- Examples
 - Ontology segmentation [Seidenberg and Rector 2006, 2009]
 - Traversals [Noy and Musen 2003, 2009]
 - More general framework [d'Aquin et al. 2007]



Pro and contra graph-based approaches

Pro

- Modules can usually be extracted efficiently time polynomial in the size of *O* → robustly scalable
- Easy to implement
- Applicable to many logics

Contra

- Heuristic, no characterisation of the expected module contents
- In particular, no logical guarantees such as entailment preservation
- \rightsquigarrow Modules typically lose knowledge from ${\cal O}$



A-posteriori approaches with coverage

Coverage

 $\mathcal{M}\subseteq \mathcal{O} \text{ covers } \mathcal{O} \text{ for } \Sigma$ if

all $\Sigma\text{-}\mathsf{consequences}$ of $\mathcal O$ already follow from $\mathcal M.$

- i.e., *M* preserves all knowledge in *O* about α
 Tuesday's lecture
- This guarantee is needed, e.g., for ontology reuse or reasoning
 >> Wednesday's lecture

Problems

- \bullet Of course, ${\cal O}$ is always covering
- Minimal covering modules are, in general, hard to extract
 - ➤ Tuesday's lecture



Coverage-providing module notions

- Restricted to logics where coverage can be decided efficiently e.g., MEX for acyclic *EL* ⇒ *Thursday's lecture* [Konev et al. 2008]
- Or use a tractable condition sufficient for coverage, leading to modules that always contain minimal modules Examples:
 - Locality-based modules ⇒ Wednesday's lecture [Cuenca Grau et al. 2007, 2009]
 - Modules obtained from partitions based on E-connections [Cuenca Grau et al. 2006]



Comparison of a-posteriori module extraction approaches

Module notion	Covrg.	Min.	Covered DLs	Complexity
All axioms referencing Σ	×		any	easy
Graph-based	×		any	easy



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The whole ontology	1	xx	any	easy
Min. covproviding mod.* MEX*	1 1	\ \	few acyclic \mathcal{EL}	<mark>hard</mark> easy
Locality-based mod.*	1	×	OWL	easy
E-connections based mod.	1	×	OWL	easy

*Will be covered here >> Tuesday's and Wednesday's lecture



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E-connections based mod.	1	×	OWL	easy
Modules with rewriting	√?	√√?	few?	hard?

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And now ...



Overview and comparison of modularisation approaches



Overview of the remainder of this course



Course overview

- Formal foundations of modularity (Dirk)
 - Conservative extensions
 - Inseparability
- Module extraction (Thomas)
 - Locality classes and locality-based modules
 - Module extraction algorithms and experiments
- Ontology versioning and forgetting of vocabulary (Dirk)
 - Logical difference
 - Forgetting/uniform interpolants
- Secent advances/current work (Thomas)
 - Atomic decomposition
 - Signature decomposition, relevance of terms

