



## Lecture 07:

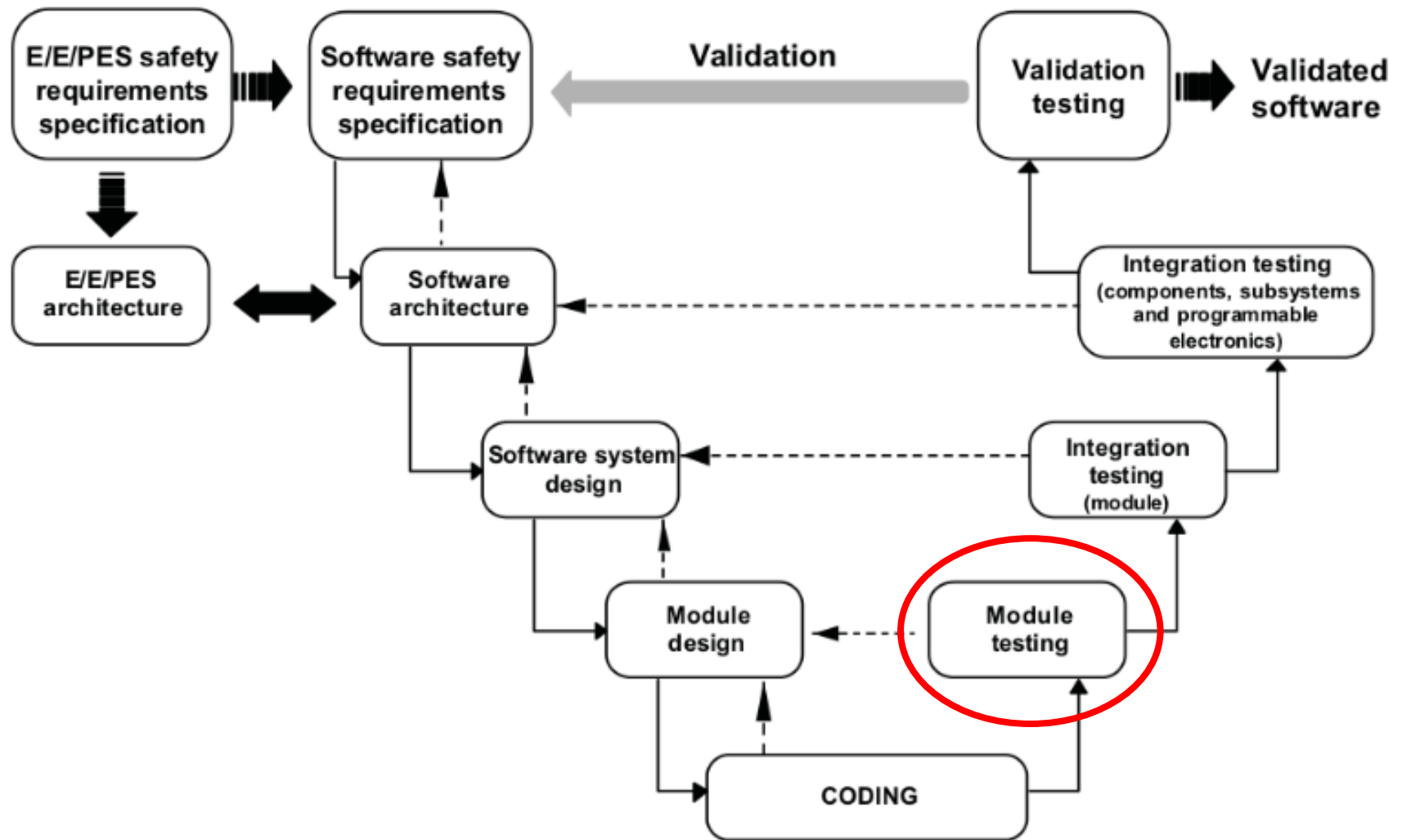
## Testing

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# Where are we?

- ▶ 01: Concepts of Quality
- ▶ 02: Legal Requirements: Norms and Standards
- ▶ 03: The Software Development Process
- ▶ 04: Hazard Analysis
- ▶ 05: High-Level Design with SysML
- ▶ 06: Formal Modelling with OCL
- ▶ 07: Testing
- ▶ 08: Static Program Analysis
- ▶ 09-10: Software Verification
- ▶ 11-12: Model Checking
- ▶ 13: Conclusions

# Testing in the Development Cycle



# What is Testing?

Testing is the process with the intent of finding errors

- ▶ In our sense, testing is a process
- ▶ The **aim** of testing is to find errors
  - ▶ derivation of test cases compared to the specification
  - ▶ inconsistency between specification and program
  - ▶ structural features of program

BUT: testing can prove the **absence** of errors under certain hypotheses – so-called **complete** test methods

see <http://www.informatik.uni-bremen.de/agbs/jp/papers/test-automation-huang-peleska.pdf>

This concept is closely related to model checking

It is well known that Dijkstra hated model checking and frowned upon testing ...

... that cause a faulty behavior of a program

Program testing can be used to show the presence of bugs, but never to show their absence.

*E.W. Dijkstra, 1972*

# Why is testing so important?

- ▶ Even if one day code can be completely verified using formal methods, tests will still be required because--
- ▶ for embedded systems, the correctness of the HW/SW integration must be verified by testing, because--
- ▶ as of today, it is infeasible to provide a correct and complete formal model for complete HW/SW systems, comprising:
  - ▶ source code,
  - ▶ machine code,
  - ▶ CPU micro code,
  - ▶ firmware on interface hardware,
  - ▶ CPUs, busses, caches, memory, and interface boards.
- ▶ This will stay infeasible in the foreseeable future

# The Testing Process

- ▶ Test cases, test plan, etc.
- ▶ System-under-test (s.u.t.)
  - ▶ Aka. TOE (target-of-evaluation) in CC
  - ▶ Aka. Implementation-under-test
- ▶ Warning -- test literature is quite expansive:

Testing is any activity aimed at evaluating an attribute or capability of a program or system and determining that it meets its required results.

*Hetzel, 1983*

# Test Levels

- ▶ **Component** and **unit tests**
  - ▶ test at the interface level of single components (modules, classes)
- ▶ **Integration test**
  - ▶ testing interfaces of components fit together
- ▶ **System test**
  - ▶ functional and non-functional test of the complete system from the user's perspective
- ▶ **Acceptance test**
  - ▶ testing if system implements contract details

# Test Methods

- ▶ Static vs. dynamic
  - ▶ With **static** tests, the code is **analyzed** without being run. We cover these methods as static program analysis later
  - ▶ With **dynamic** tests, we **run** the code under controlled conditions, and check the results against a given specification
- ▶ Central question: where do the **test cases** come from?
  - ▶ **Black-box**: the inner structure of the s.u.t. is opaque, test cases are derived from specification **only**.
  - ▶ **Grey-box**: some inner structure of the s.u.t. is known, e.g. module architecture.
  - ▶ **White-box**: the inner structure of the s.u.t. is known, and tests cases are derived from the source code **and** coverage objectives for the source code



# Black-Box Tests

- ▶ Limit analysis:
  - ▶ If the specification limits input parameters, then values **close** to these limits should be chosen
  - ▶ Idea is that programs behave **continuously**, and errors occur at these limits
- ▶ Equivalence classes:
  - ▶ If the input parameter values can be decomposed into **classes** which are treated equivalently, test cases have to cover all classes
- ▶ Smoke test:
  - ▶ “Run it, and check it does not go up in smoke.”

# Example: Black-Box Testing

- ▶ Equivalence classes or limits?

## **Example: A Company Bonus System**

The loyalty bonus shall be computed depending on the time of employment. For employees of more than three years, it shall be 50% of the monthly salary, for employees of more than five years, 75%, and for employees of more than eight years, it shall be 100%.

- ▶ Equivalence classes or limits?

## **Example: Air Bag**

The air bag shall be released if the vertical acceleration  $a_v$  equals or exceeds  $15 \text{ m/s}^2$ . The vertical acceleration will never be less than zero, or more than  $40 \text{ m/s}^2$ .

# Black-Box Tests

- ▶ Quite typical for **GUI tests**, or **functional testing**
- ▶ Testing **invalid input**: depends on programming language – the stronger the typing, the less testing for invalid input is required
  - ▶ Example: consider lists in C, Java, Haskell
  - ▶ Example: consider object-relational mappings<sup>1</sup> (ORM) in Python, Java

1) Translating e.g. SQL-entries to objects

# Complete Model-based Black-box Testing

- ▶ Create a model  $M$  of the expected system behaviour
- ▶ Specify a **fault model**  $(M, \leq, Dom)$  with reference model  $M$ , **conformance relation**  $\leq$  and **fault domain**  $Dom$  (a collection of models that may or may not conform to  $M$ )
- ▶ Derive test cases from fault model
- ▶ The resulting test suite is **complete** if
  - ▶ Every conforming SUT will pass all tests (**soundness**)
  - ▶ Every non-conforming SUT whose true behavior is reflected by a member of the fault domain fails at least on test case (**exhaustiveness**)
  - ▶ (nothing is guaranteed for SUT behaviors outside the fault domain)

# Example: the W-Method

- ▶ The W-Method specifies a recipe for constructing complete test suites for finite state machines (FSMs) with conformance relation " $\sim$ "  
**language equivalence (I/O-equivalence):**

- ▶ Create a state cover  $V$
- ▶ Create a characterization set  $W$
- ▶ Assume that implementation has at most  $m \geq n$  states ( $n$  is the number of states in the observable, minimized reference model)
- ▶ Create test suite according to formula

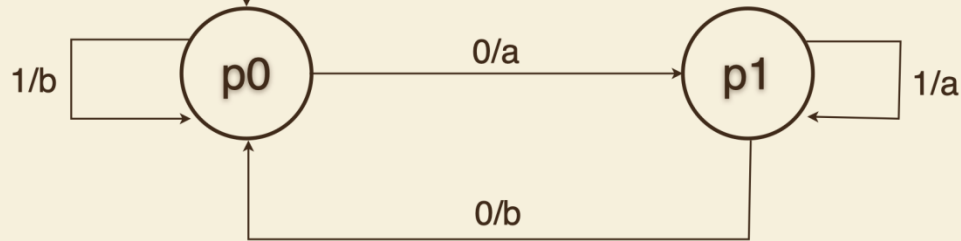
$$W = V. \left( \bigcup_{i=0}^{m-n+1} I^i \right). W$$

$I$  : input alphabet

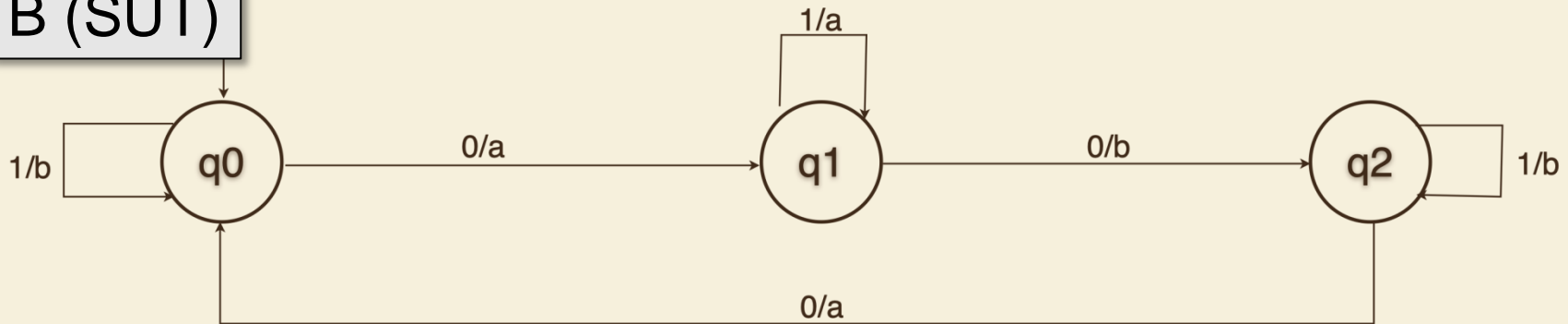
$I^i$  : input traces of length  $i$

$A.B$  : all traces of  $A$  concatenated with all traces from  $B$

## A (Reference Model)



## B (SUT)



$$n = 2 \quad \bigcup_{i=0}^{m-n+1} I^i = \{\varepsilon, 0, 1, 0.0, 0.1, 1.0, 1.1\}$$

$$V = \{\varepsilon, 0\}$$

$$W = \{1\}$$

$$\begin{aligned} \mathcal{W}(A) &= V \cdot \left( \bigcup_{i=0}^{m-n+1} I^i \right) \cdot W \\ &= \{1, 0.1, 1.1, 0.0.1, 0.1.1, 1.0.1, 1.1.1, \\ &\quad 0.0.0.1, 0.0.1.1, 0.1.0.1, 0.1.1.1\} \end{aligned}$$

SUT error is uncovered by  
0.0.0.1

# Property-based Testing

- ▶ In property-based testing (or random testing), we generate **random** input values, and check the results against a given **executable** specification.
- ▶ Attention needs to be paid to the **distribution** values.
- ▶ Works better with **high-level languages**, where the datatypes represent more information on an abstract level and where the language is powerful enough to write comprehensive executable specifications (i.e. Boolean expressions).
  - ▶ Implementations for e.g. Haskell (QuickCheck), Scala (ScalaCheck), Java
- ▶ Example: consider list reversal in C, Java, Haskell
  - ▶ Executable spec: reversal is idempotent and distributes over concatenation.
  - ▶ Question: how to generate random lists?

# White-Box Tests

- ▶ In white-box tests, we derive test cases based on the structure of the program (**structural testing**)
  - ▶ To abstract from the source code (which is a purely syntactic artefact), we consider the **control flow graph** of the program.

## Def: Control Flow Graph (CFG)

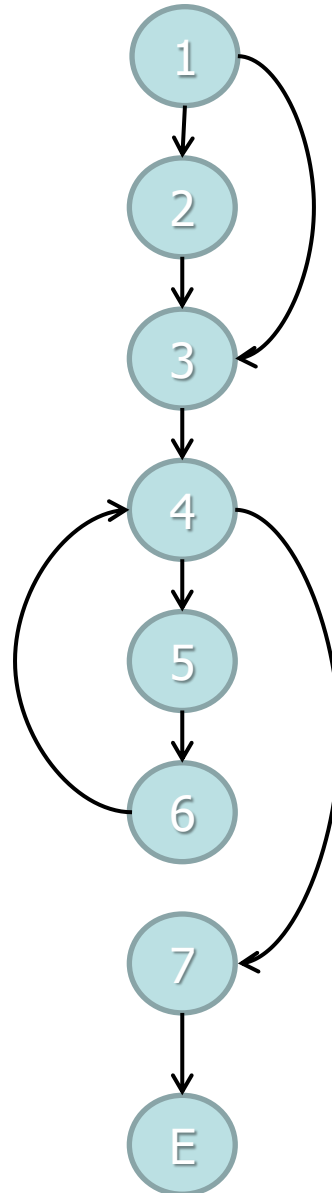
- nodes as elementary statements (e.g. assignments, **return**, **break**, . . . ), as well as control expressions (e.g. in conditionals and loops), and
- vertices from  $n$  to  $m$  if the control flow can reach a node  $m$  coming from a node  $n$ .

- ▶ Hence, **paths** in the CFG correspond to **runs** of the program.



# Example: Control-Flow Graph

```
if (x < 0) /*1*/ {  
    x = -x; /*2*/  
}  
z = 1; /*3*/  
  
while (x > 0) /*4*/ {  
    z = z * y; /*5*/  
    x = x - 1; /*6*/  
}  
return z; /*7*/
```



An execution path is a path through the CFG ending with an exit node.

Examples:

- [1,3,4,7, E]
- [1,2,3,4,7, E]
- [1,2,3,4,5,6,4,7, E]
- [1,3,4,5,6,4,5,6,4,7, E]
- ...

# Coverage

- ▶ **Statement coverage:**

Measures the percentage of statements that were covered by the tests.

100% statement coverage is reached if each **node** in the CFG has been visited at least once.

- ▶ **Branch coverage:**

Measures the percentage of **edges** (emanating from branching or non-branching nodes) covered by the tests.

100% branch coverage is reached if every edge of the CFG has been traversed at least once.

- ▶ **Path coverage:**

Measures the percentage of CFG **paths** that have been covered by the tests.

100% path coverage is achieved if every path of the CFG has been covered at least once.

# Decision Coverage

- ▶ **Decision coverage:**

Measures the coverage of conditional branches (i.e., edges emanating from conditional nodes). 100% decision coverage is reached if the tests cover all conditional branches.

- ▶ **Decision coverage vs. branch coverage:**

- ▶ If branch coverage is 100%, then decision coverage is 100% and vice versa.
- ▶ A lower percentage  $p < 100\%$  of branch coverage, however, has a different meaning than a decision coverage of  $p$ , because
- ▶ branch coverage considers all edges, whereas
- ▶ decision coverage considers edges emanating from decision nodes only

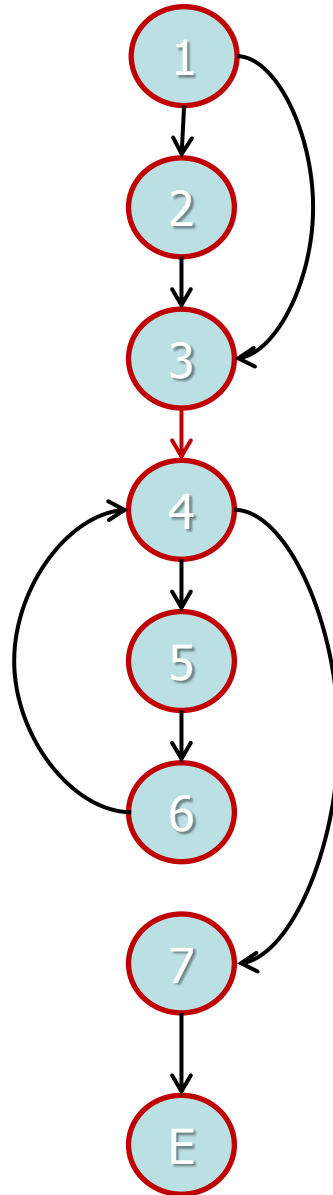
# Example: Statement Coverage

```
if (x < 0) /*1*/ {  
  x = -x; /*2*/  
}
```

```
z = 1; /*3*/
```

```
while (x > 0) /*4*/ {  
  z = z * y; /*5*/  
  x = x - 1; /*6*/  
}
```

```
return z; /*7*/
```



- ▶ Which (minimal) path covers all statements?

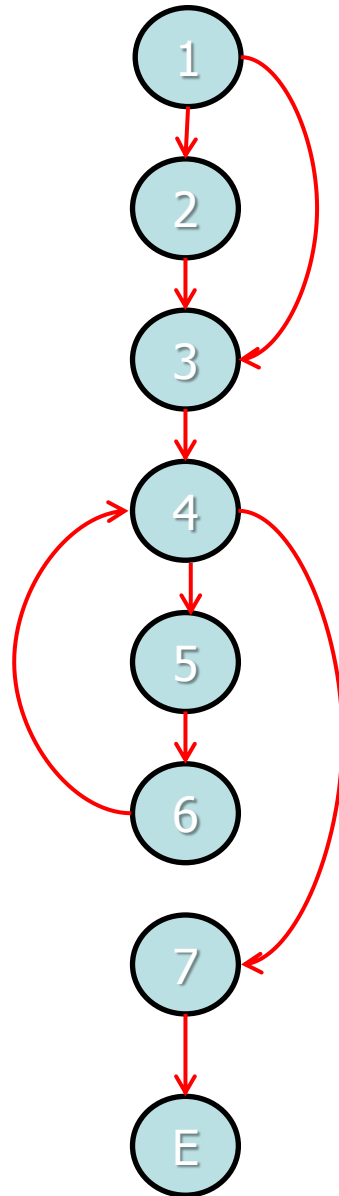
$p = [1, 2, 3, 4, 5, 6, 4, 7, E]$

- ▶ Which state generates  $p$ ?

$x = -1$   
 $y$  any  
 $z$  any

# Example: Branch Coverage

```
if (x < 0) /*1*/ {  
  x = -x; /*2*/  
}  
z = 1; /*3*/  
while (x > 0) /*4*/ {  
  z = z * y; /*5*/  
  x = x - 1; /*6*/  
}  
return z; /*7*/
```



- ▶ Which (minimal) paths cover all vertices?

$p_1 = [1, 2, 3, 4, 5, 6, 4, 7, E]$

$p_2 = [1, 3, 4, 7, E]$

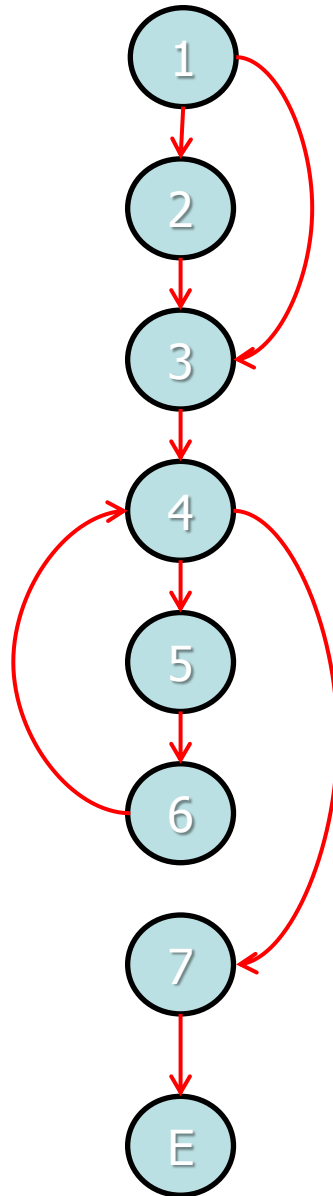
- ▶ Which states generate  $p_1, p_2$ ?

	$p_1$	$p_2$
x	-1	0
y	any	any
z	any	Any

- ▶ Note  $p_3$  with  $x = 1$  does not add coverage.

# Example: Path Coverage

```
if (x < 0) /*1*/ {  
    x = -x; /*2*/  
}  
z = 1; /*3*/  
while (x > 0) /*4*/ {  
    z = z * y; /*5*/  
    x = x - 1; /*6*/  
}  
return z; /*7*/
```



► How many paths are there?

► Let  $q_1 = [1,2,3]$   
 $q_2 = [1,3]$   
 $p = [4,5,6]$   
 $r = [4,7,E]$

then all paths are

$$P = (q_1|q_2) p^* r$$

► Number of possible paths:

$$|P| = 2 \cdot \text{MaxInt} - 1$$

# Statement, Branch and Path Coverage

## ▶ **Statement Coverage:**

- ▶ Necessary but not sufficient, not suitable as only test approach.
- ▶ Detects dead code (code which is never executed).
- ▶ About 18% of all defects are identified.

## ▶ **Branch coverage:**

- ▶ Least possible single approach.
- ▶ Needs to be achieved by (specification-based) tests for avionic software of DAL-C – does not suffice for DAL-B or DAL-A.
- ▶ Detects dead code, but also frequently executed program parts.
- ▶ About 34% of all defects are identified.

## ▶ **Path Coverage:**

- ▶ Most powerful structural approach;
- ▶ Highest defect identification rate (close to 100%);
- ▶ But no **practical** relevance.

# Decision Coverage Revisited

- ▶ Decision coverage requires that for all decisions in the program, each possible outcome is considered once.
- ▶ **Problem:** cannot sufficiently distinguish Boolean expressions.
  - ▶ Example: for  $A \parallel B$ , the following are sufficient:

A	B	Result
False	False	False
True	False	True

- ▶ But this does not distinguish  $A \parallel B$  from  $A$ ;  $B$  is effectively not tested.



# Decomposing Boolean Expressions

- ▶ The binary Boolean operators include conjunction  $x \wedge y$ , disjunction  $x \vee y$ , or anything expressible by these (e.g. exclusive disjunction, implication)

## Elementary Boolean Terms

An elementary Boolean term does not contain binary Boolean operators, and cannot be further decomposed.

- ▶ An elementary term is a variable, a Boolean-valued function, a relation (equality  $=$ , orders  $<$ ,  $\leq$ ,  $>$ ,  $\geq$ , etc.), or a negation of these.
- ▶ This is a fairly syntactic view, e.g.  $x \leq y$  is elementary, but  $x < y \vee x = y$  is not, even though they are equivalent.
- ▶ In formal logic, these are called **literals**.

# Simple Condition Coverage

- ▶ For each decision in the program, each elementary Boolean term (condition) evaluates to *True* and *False* at least once
- ▶ Note that this does not say much about the possible value of the condition
- ▶ Example:

if (temperature > 90 && pressure > 120) { ... }

<i>C1</i>	<i>C2</i>	<i>Result</i>
False	False	False
False	True	False
True	False	False
True	True	True

-- These two would be enough  
-- for condition coverage

# Modified Condition Coverage

- ▶ It is not always possible to generate all possible combinations of elementary terms, e.g.  $3 \leq x \ \&\& \ x < 5$ .
- ▶ In modified (or minimal) condition coverage, all possible combinations of those elementary terms the value of which determines the value of the whole condition need to be considered.
- ▶ Example:  $3 \leq x \ \&\& \ x < 5$

$3 \leq x$	$x < 5$	Result
False	False	False
False	True	False
True	False	True
True	True	True

- ▶ Another example:  $(x > 1 \ \&\& \ ! p) \ || \ p$

# Modified Condition/Decision Coverage

- ▶ Modified Condition/Decision Coverage (MC/DC) is required by the “aerospace norm” **DO-178B** for Level A software.
- ▶ It is a **combination** of the previous coverage criteria defined as follows:
  - ▶ Every point of entry and exit in the program has been invoked at least once;
  - ▶ Every decision in the program has taken all possible outcomes at least once;
  - ▶ Every condition (i.e. elementary Boolean terms earlier) in a decision in the program has taken all possible outcomes at least once;
  - ▶ Every condition in a decision has been shown to independently affect that decision’s outcome.

# How to achieve MC/DC

- ▶ **Not:** Here is the source code, what is the minimal set of test cases?
- ▶ **Rather:** From requirements we get test cases, do they achieve MC/DC?
- ▶ Example:
  - ▶ Test cases:

<b>Test case</b>	<b>1</b>	<b>2</b>	<b>3</b>	<b>4</b>	<b>5</b>
Input A	F	F	T	F	T
Input B	F	T	F	T	F
Input C	T	F	F	T	T
Input D	F	T	F	F	F
<i>Result Z</i>	<i>F</i>	<i>T</i>	<i>F</i>	<i>T</i>	<i>T</i>

Source Code:

$Z = (A \parallel B) \&\& (C \parallel D)$

**Question:** do test cases achieve MC/DC?

Source: Hayhurst *et al*, A Practical Tutorial on MC/DC. NASA/TM2001-210876

# Example: MC/DC

Determining MC/DC:

1. Are all decisions covered?
2. Eliminate masked inputs (recursively)
  - ▶ False for && masks other input
  - ▶ True for || masks other input
3. Remaining unmasked test cases must cover all conditions.

Here:

- ▶ Result is both F and T, so decisions covered.
- ▶ Masking:
  - ▶ In test case 1, C and D are masked
  - ▶ In test case 3, A and B are masked
  - ▶ Recursive masking as shown
- ▶ Remaining cases cover T, F for A, B, C, D
  - ▶ MC/DC achieved
  - ▶ In fact, test case 4 not even needed (?)

Source Code

$Z = (A \ || \ B) \ \&\& \ (C \ || \ D)$

<b>Test case</b>	<b>1</b>	<b>2</b>	<b>3</b>	<b>4</b>	<b>5</b>
Input A	F	F	T	F	T
Input B	F	T	F	T	F
Input C	T	F	F	T	T
Input D	F	T	F	F	F
<i>Result Z</i>	<i>F</i>	<i>T</i>	<i>F</i>	<i>T</i>	<i>T</i>

# Summary

- ▶ (Dynamic) Testing is the controlled execution of code, and comparing the result against an expected outcome.
- ▶ Testing is (traditionally) the main way for **verification**.
- ▶ Depending on how the test cases are derived, we distinguish **white-box** and **black-box** tests.
- ▶ In black-box tests, we can consider **limits** and **equivalence classes** for input values to obtain test cases.
- ▶ In white-box tests, we have different notions of **coverage**: statement coverage, path coverage, condition coverage, etc.
- ▶ Next week: **Static testing** aka. static **program analysis**