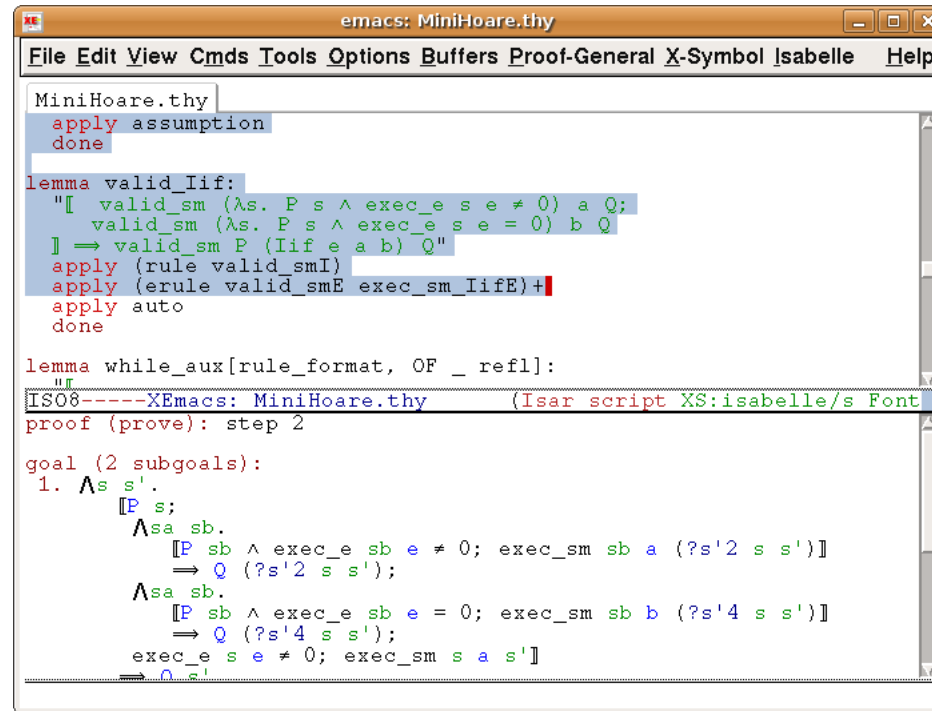

Managing Proof Documents for Asynchronous Processing

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Traditional User Interfaces



```
emac: MiniHoare.thy
File Edit View Cmds Tools Options Buffers Proof-General X-Symbol Isabelle Help
MiniHoare.thy
apply assumption
done

lemma valid_Iif:
  "[ valid_sm (λs. P s ^ exec_e s e ≠ 0) a Q;
    valid_sm (λs. P s ^ exec_e s e = 0) b Q
  ] ⇒ valid_sm P (Iif e a b) Q"
  apply (rule valid_smI)
  apply (erule valid_smE exec_sm_IifE)+
  apply auto
  done

lemma while_aux[rule_format, OF _ refl]:
  "!!
  ISO8-----XEmacs: MiniHoare.thy (Isar script XS:isabelle/s Font
proof (prove): step 2

goal (2 subgoals):
1. Λs s'.
  [P s;
   Λsa sb.
   [P sb ^ exec_e sb e ≠ 0; exec_sm sb a (?s'2 s s')]
   ⇒ Q (?s'2 s s');
   Λsa sb.
   [P sb ^ exec_e sb e = 0; exec_sm sb b (?s'4 s s')]
   ⇒ Q (?s'4 s s');
   exec_e s e ≠ 0; exec_sm s a s']
  ⇒ Q s'
```

- Idea: Script buffer [3, 1]
- UI keeps proof script for batch replay
- Linear processing, commands become “locked”
- Focus on mechanics of proving – user-friendly?

Document-Centered View

- Metaphor “proof document”
 - User edits human-readable a proof document
 - Prover checks the consistency
- Assisted authoring [2]
 - Isar as human-readable proof language [8]
 - Backflow: Assistance by prover for editing
 - Processing linear
- Plat Ω approach [6]
 - Near-natural, text-book style input language
 - Front-end parses structure & computes structural diff
 - Triggers necessary (re-)checking by Ω mega prover

Asynchronous Proof Processing [7]

- Observation
 - 95% of processing time spent in proofs
 - . . . but later commands do not depend on proofs at all
- Asynchronous proof processing (APP)
 - Process definitions, theorems, etc. immediately
 - Check syntax & explicit references immediately
 - Execute proofs when processor would be idle

Applications

- Parallel execution
- Traverse intermediate theory quickly between . . .
 - Definition & theorem
 - Theorem & reference
 - Tactic definition & application
- Strategy: Continuous proof processing
 - Isabelle processes the newest version of each command
 - Success/failure is presented as annotations

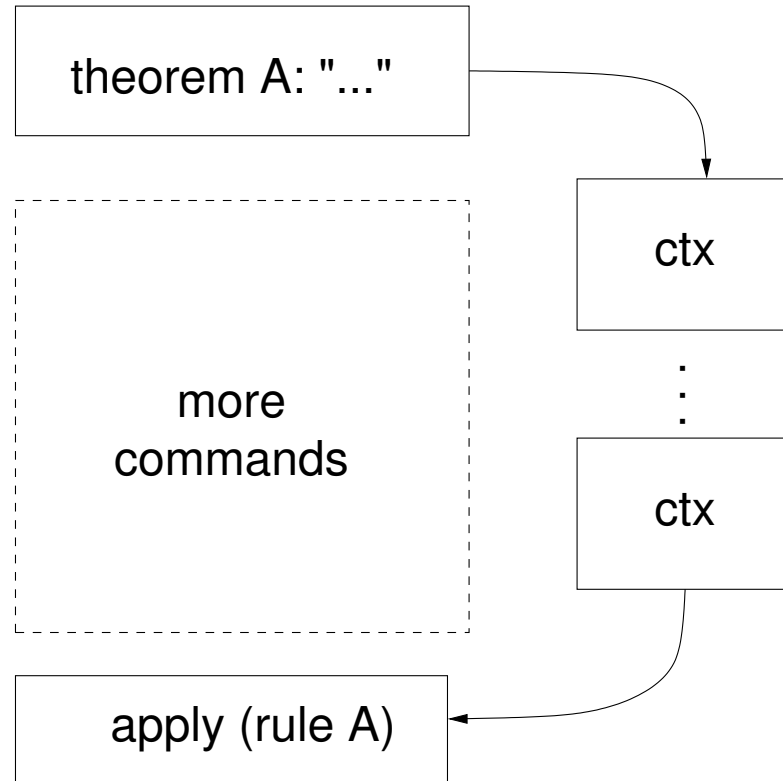
⇒ IDE-like theory development

Outline

- Design considerations for asynchronous proof checking
- Protocol for UI/prover communication
- Infrastructure for asynchronous proof checking
- Design principle: minimal assumptions = maximal re-use
- Applied to Isabelle/Isar – ongoing work with Makarius Wenzel

Design Considerations

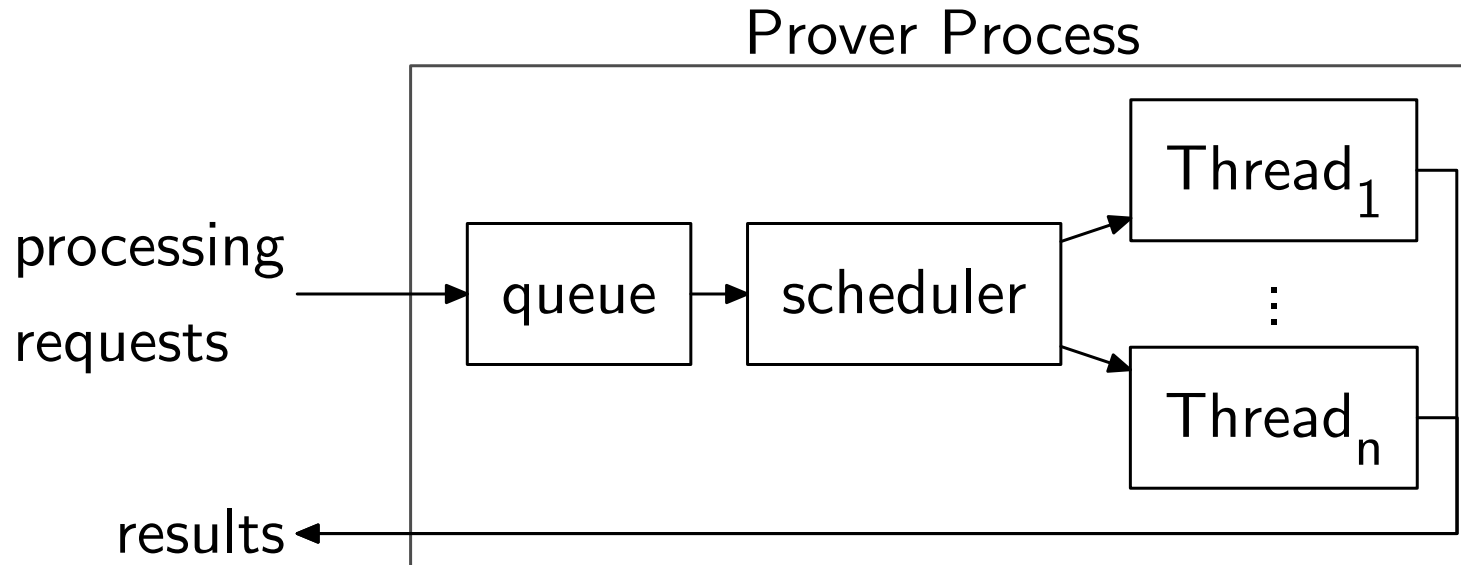
Dependencies in Isabelle



Implicit Dependencies

- Implicit rule sets: `intro`, `simp`, `cong`, ...
- Used by various methods: `blast`, `simp`, `auto`, ...
- Proofs may break because new rules are added
- Extensible: Isar framework is not aware of dependencies
- Conclusions
 - Dependencies are hard to track
 - Design APP to be independent of possible solution⇒ Don't prescribe particular dependency model

Order of Processing



- Active Object [5] / Lightweight Executable Framework [4]
- Asynchronous / parallel / concurrent / distributed processing
 - Results computed at the convenience of the service provider
 - Clients perform other computations until result arrives

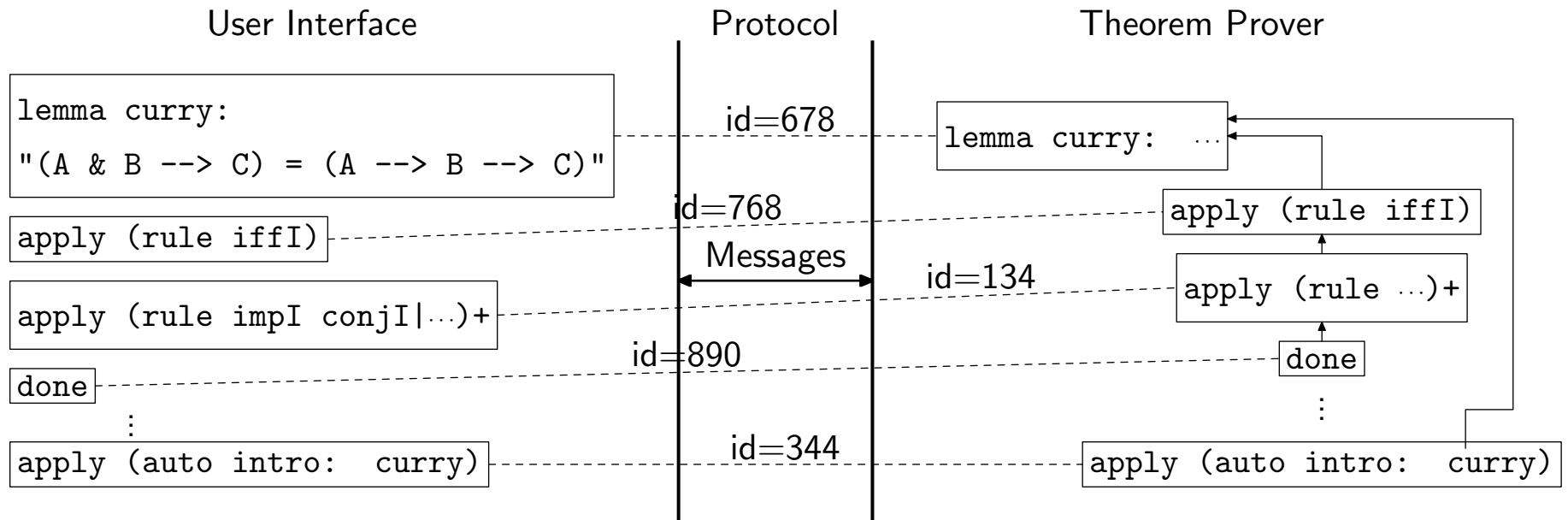
Minimal Assumptions — The Prover

- Principle: The UI makes no assumptions about
 - The processing order of individual commands
 - The dependencies between / independence of commands
- Benefit: Incremental implementation
 1. Start with linear processing
 2. Develop dependency tracking models
 3. Multi-threading & asynchronous processing
- Benefit: Software structure
 - Adapt implementation to existing mechanisms
 - Exploit internal knowledge in implementation
 - Changes possible without changing the UI

Division of Responsibilities

- Isabelle
 - Processing order
 - Dependency tracking
- User Interface
 - Identification commands in the proof document
 - Synchronization of access by user / Isabelle

Identifying Commands



- Splitting done by user interface
 - Isabelle works at command level \Rightarrow Facilitate re-work
 - Immediate reaction; auxiliary information (outline, . . .)
 - Service available to other provers

Demands on Identifying Commands

- Reliability
 - Entire protocol depends on identity of commands
 - ⇒ Must not fail in any case
- Stability
 - State information attached to commands
 - Creation/deletion may cause expensive (re-)proving
 - ⇒ Avoid spurious changes to document structure
- Interactivity
 - User keeps typing
 - Prover keeps proving
 - ⇒ Recognize & propagate changes immediately
- Synchronization: Negotiate creation/deletion with prover

Solution: Separating Keywords

- Conclusion: build small, specialized incremental parser
- Isar design: keywords separate commands
- Implementation
 - Maintain partitioning of proof document into commands
 - Split elements on creation of keywords
 - Join elements on deletion of keywords
 - Keep track of quotes (comments, inner syntax, etc.)
 - Delay split/join during negotiation with prover
- Generalization: Splitter
 - Informed about textual changes
 - Split/join components in arbitrary (Java) code

Protocol Part 1: Document Structure

- UI sends messages
 - `create(id, prevId)`
 - `destroy(id)`
- Isabelle
 - Maintains doubly-linked list of commands
 - And hash map $id \mapsto \textit{command object}$

Synchronization

- Two complementary models of the proof document
 - UI: Text, keywords, markups, outline, . . .
 - Isabelle: Processing state, dependencies, . . .
- Need synchronization for consistency
- Mutexes, semaphores, etc. need shared memory
- Approach “ownership”
 - Only one process owns a shared object at a given time
 - And only the owner may access the object

Ownership of Commands

- UI is Owner
 - User may edit the command text
 - Split/join with commands are possible
 - UI may send messages create / destroy
 - Isabelle maintains no references (dependencies etc.)
 - But: Command is still present in the document structure
- Isabelle is Owner
 - Isabelle may choose to execute the command
 - Isabelle may register dependencies on command
 - UI receives processing information from Isabelle
 - User may not edit the command

Protocol Part 2: Negotiating Ownership

- Transfer UI \rightarrow Isabelle
 - Send command text to Isabelle
 - Prevent user from editing
 - Quick — except for long sequences
- Transfer Isabelle \rightarrow UI
 - Undo any dependent commands & erase references
 - Kill thread currently processing command
 - May take time
- Protocol
 - UI \Rightarrow Prover
 - `send(id, cmdStr)`
 - `revoke(id)`
 - Prover \Rightarrow UI
 - `released(id)`
 - `request(id)`

That's it — basically

Protocol Part 3: Information Messages

- Isabelle decides on processing order
- But reports progress, success/failure
- Informational messages
 - `queued(id)`
 - `startProcessing(id)`
 - `success(id)`
 - `error(id)`
 - `result(id,result)`
- Note: no reaction by interface required

Sequences of Messages?

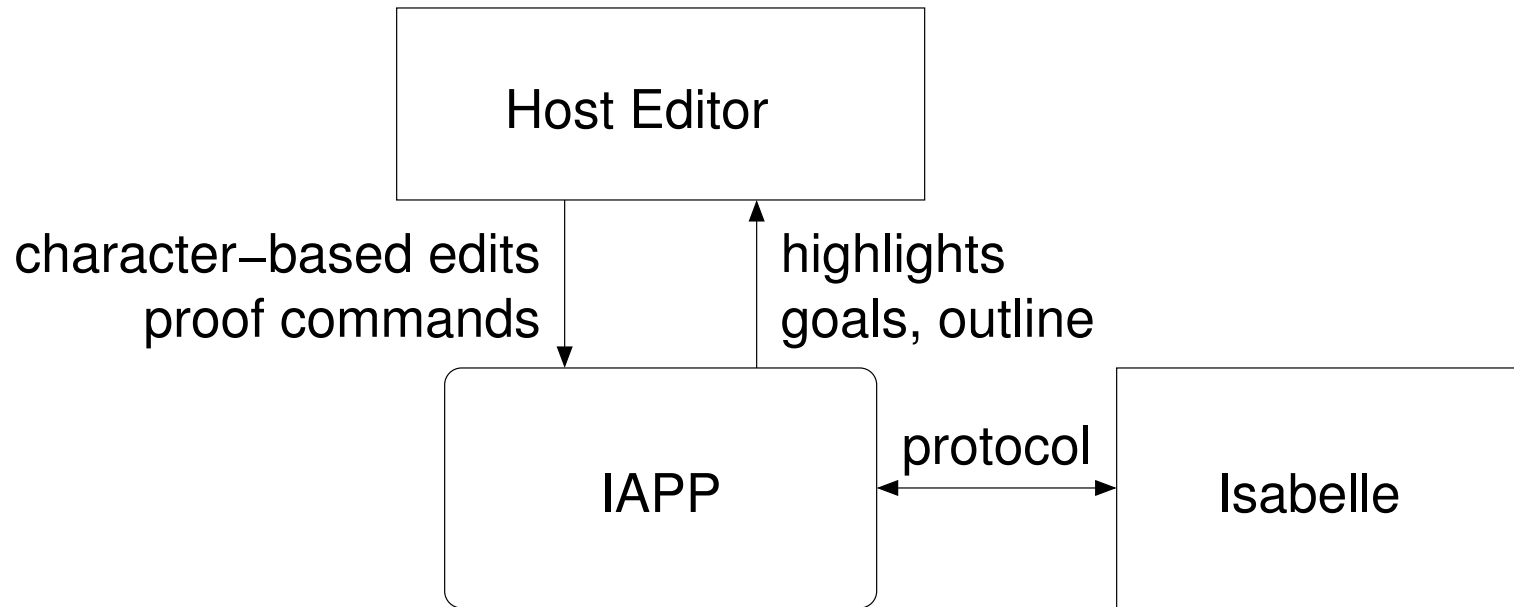
- Approach: Asynchronous processing
 - Prover & UI work independently
 - Events occur on commands (created, destroyed, . . .)
 - Messages are only notifications of events
- ⇒ No specific ordering of messages prescribed
- Instead: state model for commands
 - Events change the state of a command
 - The state determines which events may legally occur

Software Design

Points of Variation

- APP is current development
- Three points of variation
 - Editor: Emacs, Eclipse, Netbeans, JEdit, . . .
 - Interaction mode: Continuous, explicit requests, . . .
 - Prover: Extension of Isabelle; perhaps other provers
- Solution: Infrastructure for Asynchronous Proof Processing
 - Abstracts over concrete editor used
 - Abstracts over the prover used

Minimal Assumptions — The Editor



- Infrastructure for Asynchronous Proof Processing (IAPP)
- Uses APP protocol for Isabelle communication
- Displays results in editor
- Receives user actions from editor

Controlling Command Execution

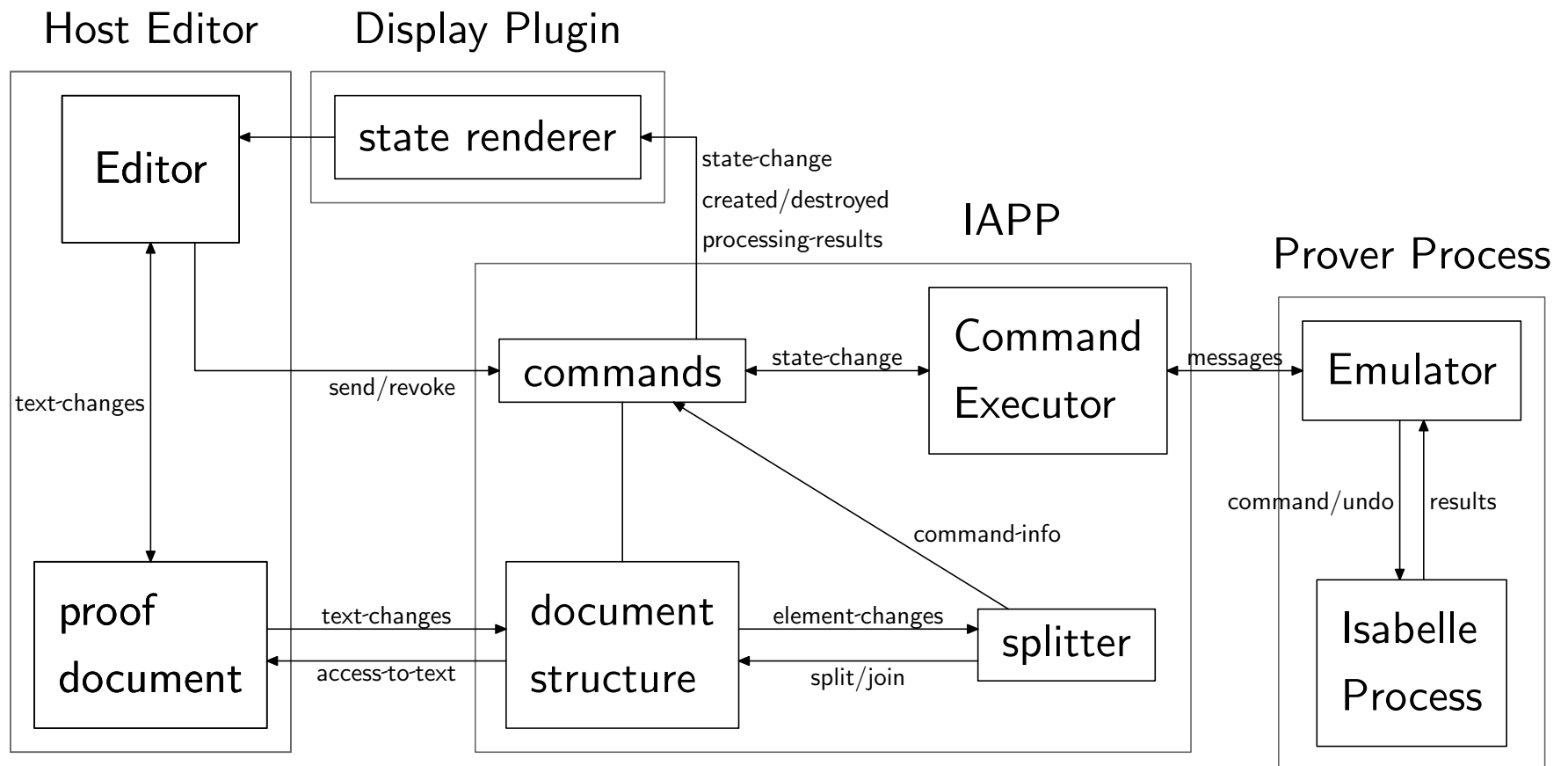
- The state of a command determines its execution
 - to be send \Rightarrow will be executed at some point
 - to be revoked \Rightarrow will be undone at some point
 - Nice: Program logic for sending/revoking is small
- \Rightarrow Facilitates experiments with interaction styles

```
CommandState state = elem.getData();  
if (state.isIdle())  
    state.send();  
else  
    state.revoke();
```

Synchronous Proving with the IAPP

- Provers may always choose to linear processing order
 - Possible: Emulator that translates IAPP \leftrightarrow TTY
- ⇒ Talk to existing provers in the background
- Gains
 - Experiment with styles of user interactions
 - Elements of APP desirable for traditional interface
 - Identification of commands for outline etc.
 - Event-based \Rightarrow non-stalling user interface
 - Deals gracefully with long traces

Overall Software Structure



Conclusion

- UI for asynchronous proof processing
- Design principle: minimal assumptions = maximal re-use
- Collaboration of UI + prover for APP
 - Maintenance of document structure
 - Protocol
 - State model for commands
- Characteristics
 - Design entirely event-based \Rightarrow asynchronous processing
 - Also TTY-based provers receive benefits

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