Query Conservative Extensions in Horn Description Logics with Inverse Roles

Jean Christoph Jung, Carsten Lutz, Mauricio Martel, and Thomas Schneider
Fachbereich Informatik, Universität Bremen, Germany

Abstract

We investigate the decidability and computational complexity of query conservative extensions in Horn description logics (DLs) with inverse roles. This is more challenging than without inverse roles because characterizations in terms of unbounded homomorphisms between universal models fail, blocking the standard approach to establishing decidability. We resort to a combination of automata and mosaic techniques, proving that the problem is \( \sharp \text{PSPACE} \)-complete in Horn-\( \text{ALCH} \bowtie \mathcal{F} \) and \( \mathcal{ECL} \). We obtain the same upper bound for deductive conservative extensions, for which we also prove a \( \text{ConEXPTIME} \) lower bound.

1 Introduction

In the past years, access of incomplete data mediated by description logic (DL) ontologies has gained increasing importance [Poggi et al., 2008; Bienvenu and Ortiz, 2015]. The main idea is to specify domain knowledge and semantics of the data in the ontology, resulting in more complete answers to queries. Significant research activity has led to efficient algorithms and tools for a wide range of DLs such as DL-Lite [Calvanese et al., 2007], more expressive Horn-DLs [Eiter et al., 2012; Trivela et al., 2015; Bienvenu et al., 2016], and “full Boolean” DLs such as \( \text{ALC} \) [Kolliou and Glimm, 2013; Zhou et al., 2015].

In contrast to query answering, which is by now well-understood, there is a need to develop reasoning services for ontology engineering that are tailored towards query-centric applications and support tasks such as ontology versioning and module extraction from ontologies. For example, if one wants to safely replace an ontology with a new version or with a smaller subset of itself (a module), then the new ontology should preserve the answers to all queries over all ABoxes (which store the data) [Kontchakov et al., 2010]. The same guarantee ensures that one can safely replace an ontology with another version in an application [Konev et al., 2012]. In both cases, ontologies need to be tested not for their logical equivalence, but for giving the same answers to relevant queries over relevant datasets.

This requirement can be formalized using conservative extensions. In the following, we use the DL term TBox instead of ontology. A TBox \( \mathcal{T}_2 \supseteq \mathcal{T}_1 \) is a \( (\Gamma, \Sigma) \)-query conservative extension of a TBox \( \mathcal{T}_1 \), where \( \Gamma \) and \( \Sigma \) are signatures of concept/role names relevant for data and queries, respectively, if all \( \Sigma \)-queries give the same answers w.r.t. \( \mathcal{T}_1 \) and \( \mathcal{T}_2 \), for every \( \Gamma \)-ABox. Note that the subset relationship \( \mathcal{T}_2 \supseteq \mathcal{T}_1 \) is natural when replacing a TBox with a module, but not in versioning, so we might not want to insist on it. In this more general case, \( \mathcal{T}_1 \) and \( \mathcal{T}_2 \) are called \( (\Gamma, \Sigma) \)-query inseparable. Conservativity and inseparability of TBoxes, as defined above, are useful when knowledge is considered static and data changes frequently. Variants of these notions for knowledge bases (KBs), which consist of a TBox and an ABox, can be used for applications with static data [Wang et al., 2014; Arenas et al., 2016].

We also consider the basic notion of query entailment: \( \mathcal{T}_1 \) \((\Gamma, \Sigma)\)-query entails \( \mathcal{T}_2 \) if all \( \Sigma \)-queries give at least the answers w.r.t. \( \mathcal{T}_1 \) that they give w.r.t. \( \mathcal{T}_2 \), on any \( \Gamma \)-ABox. Query inseparability and conservativity are special cases of entailment: inseparability is bidirectional entailment and conservativity is entailment with the assumption that \( \mathcal{T}_1 \subseteq \mathcal{T}_2 \). It thus suffices to prove upper bounds for query entailment and lower bounds for conservative extensions. As a query language, we concentrate on conjunctive queries (CQs); since we work with Horn-DLs and quantify over the queries, this is equivalent to using unions of CQs (UCQs) or positive existential queries (PEQs). CQ entailment has been studied for various DLs [Kontchakov et al., 2009; Lutz and Wolter, 2010; Konev et al., 2012; Botoeva et al., 2016a], also in the KB version [Botoeva et al., 2016b; Botoeva et al., 2016c] and for OBDA specifications [Bienvenu and Rosati, 2015], see also the survey [Botoeva et al., 2016a]. Nevertheless, there is still a notable gap in our understanding of this notion: query entailment between TBoxes is poorly understood in Horn DLs with inverse roles, often considered a crucial feature, for which there do not seem to be any available results. This is for a reason: it has been observed in [Botoeva et al., 2016a; Botoeva et al., 2016b] that standard techniques for Horn DLs without inverse roles fail when inverse roles are added.

In fact, for Horn-DLs without inverse roles query entailment can be characterized by the existence of homomorphisms between universal models [Lutz and Wolter, 2010; Botoeva et al., 2016a]. The resulting characterizations provide an important foundation for decision procedures, often based on tree automata [Botoeva et al., 2016a].
the presence of inverse roles, however, such characterizations are only correct if we require the existence of n-bound homomorphisms, for any n. [Boteva et al., 2016a; Boteva et al., 2016b]. It is not obvious how the existence of such infinite families of bounded homomorphisms can be verified using tree automata (or related techniques) and, consequently, decidability results for query conservative extensions in Horn-DLs with inverse roles are difficult to obtain. The only result we are aware of concerns inseparability of KBs, and it is proved using intricate game-theoretic techniques.

In this paper, we develop decision procedures for query entailment and related problems in Horn DLs with inverse roles. The main idea is to provide a more refined characterization, mixing unbounded and bounded homomorphisms and using bounded homomorphisms only in places where this is strictly necessary. We can then deal with the “unbounded part” using tree automata while the “bounded part” is addressed by precomputing relevant information using a mosaic technique. In this way, we establish decidability and a 2EXPSPACE upper bound for query entailment (and thus inseparability and conservativity) in Horn-ALCHIF. Together with lower bounds from [Boteva et al., 2016c], we get 2EXPSPACE-completeness for all fragments of Horn-ALCHIF that contain ELI or Horn-ACC.

We additionally study the case of deductive entailment between TBoxes, i.e., the question whether T1 entails at least the same concept and role inclusions as well as functionality assertions over Σ as T2. This problem too has not previously been studied for Horn DLs with inverse roles. We consider ECHIF -TBoxes and show that deductive entailment is equivalent to a restricted version of query entailment. We obtain a model theoretic characterization, a decision procedure, and a 2EXPSPACE upper complexity bound. We also give a CONEXPTIME lower bound.

2 Preliminaries

2.1 Horn-ALCHIF

We introduce Horn-ALCHIF, a member of the Horn-SHIF family of DLs whose reasoning problems have been widely studied [Hustadt et al., 2007; Kotzsch et al., 2007; Eiter et al., 2008; Kazakov, 2009; Lutz and Wolter, 2012; Bárány, García et al., 2014]. Let NC, NR, NI be sets of concept, role, and individual names. A role is either a role name or an inverse role r−. As usual, we identify (r−)− and r, allowing to switch between roles names and their inverses easily. A concept inclusion (CI) is of the form L ⊑ R, where L and R are concepts defined by the syntax rules

\[
R, R' ::= \top | \bot | A | \neg A | R \sqcap R' | \neg L \sqcup R | \exists r.R | \forall r.R
\]

L, L' ::= \top | \bot | A | L \sqcap L' | L \sqcup L' | \exists r.L

with A ranging over concept names and r over roles. A role inclusion (RI) is of the form r ⊑ s with r, s roles and a functionality assertion (FA) is of the form func(r) with r a role. ELI-concepts are expressions that are built according to the syntax rule for L above, but do not use “¬”.

A Horn-ALCHIF TBox T is a set of CIs, RIs, and FAs. An ECHIF TBox is a set of ELI-concepts, RIs, and FAs. To avoid dealing with rather messy technicalities that do neither seem to be very illuminating from a theoretical viewpoint nor too useful from a practical one, we generally assume that functional roles cannot have any subroles, that is, r ⊑ s in T implies func(s) ∉ T. We conjecture that our main results also hold without that restriction. An ABox A is a non-empty set of concept and role assertions of the form A(a) and r(a, b), where A ∈ NC, r ∈ NR and a, b ∈ NI. We write ind(A) for the set of individuals in A.

The semantics is defined as usual in terms of interpretations I = (ΔI, A) complying with the standard name assumption, i.e., a IV = a for all a ∈ NI [Baader et al., 2017]. An interpretation I is a model of a TBox if it satisfies all inclusions and assertions in it, and likewise for ABoxes. A is consistent with T if T and A have a common model.

A signature Σ is a set of concept and role names. A Σ-ABox is an ABox that uses only concept and role names from Σ, and likewise for Σ-ELI-concepts and other syntactic objects.

Generally and without further notice, we work with Horn-ALCHIF TBoxes that are in a certain nesting-free normal form, that is, they contain only CIs of the form

\[
T \subseteq A, A \subseteq \bot, A \sqcap A_1 \sqsubseteq B, A \sqsubseteq \exists r.B, A \sqsubseteq \forall r.B,
\]

where A, B, A1, A2 are concept names and r, s are roles. It is well-known that every Horn-ALCHIF TBox T can be converted into a TBox T′ in normal form (introducing additional concept names) such that T is a logical consequence of T′ and every model of T can be extended to one of T′ by interpreting the additional concept names, see e.g. [Bienvenu et al., 2016]. As a consequence, all results obtained in this paper for TBoxes in normal form lift to the general case.

2.2 Query Conservative Extensions and Entailment

A conjunctive query (CQ) is of the form q(x) = ∃y ϕ(x, y), where x and y are tuples of variables and ϕ(x, y) is a conjunction of atoms of the form A(v) or r(v, w) with A ∈ NC, r ∈ NR, and v, w ∈ x ∪ y. We call x answer variables and y quantified variables of q. A CQ q is tree-shaped if it does not contain atoms of the form r(x, x) and the undirected graph (x ∪ y, { (v, w) | r(v, w) is an atom in q }) is a tree; tree-shaped CQs are thus connected and may contain multiedges. A tree-shaped CQ q is strongly tree-shaped or an stCQ if the root is the one and only answer variable and q has no multi-edges, that is, for any distinct variables z, z′ in q, there is at most one role atom that contains both z and z′.

A match of q in an interpretation I is a function π : x ∪ y → ΔI such that π(v) ∈ ΔI for every atom A(v) of q and (π(v), π(w)) ∈ r for every atom r(v, w) of q. We write I |= q(a1, . . . , an) if there is a match of q in I with π(xi) = ai for all i < n. A tuple a of elements from NI is a certain answer to q over an ABox A given a TBox T, written T, A |= q(a), if I |= q(a) for all models of T and A.

1E.g., out of 439 available ontologies in BioPortal [Matentzoglu and Parsia, 2017], only 21 (≤ 4.8%) contain the described pattern. A significant fraction of the occurrences of the pattern appear to be due to modeling mistakes.
Definition 1 Let $\Gamma, \Sigma$ be signatures and $T_1, T_2$ Horn-\text{$\mathcal{ALCHIF}$} TBoxes. We say that $T_1 (\Gamma, \Sigma)$-CQ entails $T_2$, written $T_1 \models^{\text{CQ}}_{\Gamma, \Sigma} T_2$, if for all $\Gamma$-ABoxes $A$ consistent with $T_1$ and $T_2$, all $\Sigma$-CQs $q(x)$ and all tuples $a \subseteq \text{ind}(A)$, $T_2, A \models q(a)$ implies $T_1, A \models q(a)$. In addition $T_1 \subseteq T_2$, we say that $T_2$ is a $(\Gamma, \Sigma)$-CQ conservative extension of $T_1$. If $T_1 \models^{\text{CQ}}_{\Gamma, \Sigma} T_2$ and vice versa, then $T_1$ and $T_2$ are $(\Gamma, \Sigma)$-CQ inseparable.

We also consider $(\Gamma, \Sigma)$-stCQ entailment, denoted $\models^{\text{stCQ}}_{\Gamma, \Sigma}$ and defined in the obvious way by replacing CQs with stCQs.

If $T_1 \not\models^{\text{CQ}}_{\Gamma, \Sigma} T_2$ because $T_2, A \models q(a)$ but $T_1, A \not\models q(a)$ for some $\Gamma$-ABox $A$ consistent with both $T_1$, $\Sigma$-CQ $q(x)$ and $a$, we call the triple $(A, q, a)$ a witness to non-entailment.

Example 2 Let $T_1 = \{ \text{PhDSud} \subseteq \exists \text{AdvBy}, \text{Prof}, \text{adv} \subseteq \text{AdvBy} \}$ and $T_2 = T_1 \cup \{ \text{func}(\text{AdvBy}) \}$, $\Sigma = \{ \text{Prof} \}$ and $\Gamma = \{ \text{PhDSud}, \text{adv} \}$. Then $T_1 \not\models^{\text{CQ}}_{\Gamma, \Sigma} T_2$ because of the witness $(\text{PhDSud}(\text{john}), \text{adv}(\text{mary}, \text{john})).$

If we drop from Definition 1 the condition that $A$ must be consistent with both $T_1$ and $T_2$, then we obtain an alternative notion of CQ entailment that we call CQ entailment with inconsistent ABoxes. While this notion trivially implies CQ entailment in the original concept, the converse fails.

Example 3 Let $T_1 = \emptyset$, $T_2 = \{ A_1 \cap A_2 \subseteq \bot \}$ and $\Gamma = \{ A_1, A_2 \}$. Then $T_1 \not\models^{\text{CQ}}_{\Gamma, \Sigma} T_2$ but $T_1$ does not $(\Gamma, \Sigma)$-CQ entail $T_2$ with inconsistent ABoxes.

The following lemma relates the two notions of CQ entailment.

Lemma 4 CQ entailment with inconsistent ABoxes can be decided in polynomial time given access to oracles deciding CQ entailment and CQ evaluation.

Consequently and since CQ evaluation is in EXPTime in Horn-\text{$\mathcal{ALCHIF}$} [Lenter et al., 2008], all complexity results obtained in this paper also apply to CQ entailment with inconsistent ABoxes.

It is easy to see that $T_1 \not\models^{\text{CQ}}_{\Gamma, \Sigma} T_2$ if there is a $\Gamma$-role $r$ and a $\Sigma$-role $s$ with $T_2 \models r \subseteq s$ but $T_1 \not\models r \subseteq s$. We write $T_1 \models^{\text{RI}}_{\Gamma, \Sigma} T_2$ if there are no such $r$ and $s$. Clearly, $T_1 \models^{\text{RI}}_{\Gamma, \Sigma} T_2$ can be decided via $|\Gamma| \cdot |\Sigma|$ many Horn-\text{$\mathcal{ALCHIF}$} subsumption tests, thus in EXPTime [Tobies, 2001]. It is thus safe to assume $T_1 \models_{\Gamma, \Sigma}^{\text{RI}} T_2$ when deciding CQ entailment, which we will generally do from now on to avoid dealing with special cases.

2.3 Deductive Conservative Extensions

Another natural notion of entailment is deductive entailment, which generalizes the notion of deductive conservative extensions [Ghilardi et al., 2006, Lutz et al., 2007, Konev et al., 2009, Lutz and Wolter, 2010], and which separates two $\Box$Box in terms of concept and role inclusions and functionality assertions, instead of ABoxes and queries.

Definition 5 Let $\Sigma$ be a signature and let $T_1$ and $T_2$ be $\mathcal{ELHI}, \mathcal{FL}_T$ TBoxes. We say that $T_1$ $\Sigma$-deductively entails $T_2$, written $T_1 \models^{\text{ELHI}, \mathcal{FL}_T}_{\Sigma} T_2$, if for all $\Sigma$-$\mathcal{ELI}$-concept inclusions $\alpha$ and all $\Sigma$-$\mathcal{RI}$s and $\Sigma$-$\mathcal{FA}$s $\alpha$: $T_2 \models \alpha$ implies $T_1 \models \alpha$. If additionally $T_1 \subseteq T_2$, then we say that $T_2$ is a $\Sigma$-deductive conservative extension of $T_1$. If $T_1 \models^{\text{ELHI}, \mathcal{FL}_T}_{\Sigma} T_2$ and vice versa, then $T_1$ and $T_2$ are $\Sigma$-deductively inseparable.

Although closely related, it is not difficult to see that deductive and query entailment are orthogonal.

Example 6 (1) Let $T_1$, $T_2$ be as in Example 2 and $\Sigma = \{ A_1, A_2, B \}$. Then $T_1 \models^{\text{CQ}}_{\Sigma} T_2$, but $T_1 \not\models^{\text{ELHI}, \mathcal{FL}_T}_{\Sigma} T_2$.

(2) Let $T_1 = \emptyset$ and $T_2 = \{ A \subseteq \exists \text{r.B} \}$, and $\Sigma = \{ A, B \}$. Then $T_1 \not\models^{\text{CQ}}_{\Sigma} T_2$, but $T_1 \not\models^{\text{ELHI}, \mathcal{FL}_T}_{\Sigma} T_2$ as witnessed by $\{ (A(a)) \}, \exists \text{r.B}(x, a)$. However, $T_1 \models^{\text{ELHI}, \mathcal{FL}_T}_{\Sigma} T_2$.

Nevertheless, the two notions are sufficiently closely related so that we have the following.

Lemma 7 In $\mathcal{ELHI}, \mathcal{FL}_T$, deductive entailment can be decided in polynomial time given access to oracles for stCQ entailment and stCQ evaluation.

2.4 Homomorphisms and the Universal Model

For interpretations $I_1, I_2$ and a signature $\Sigma$, a $\Sigma$-homomorphism from $I_1$ to $I_2$ is a total function $h : \Delta I_1 \rightarrow \Delta I_2$ such that (1) $h(a) = a$ for all $a \in \text{at}_1$, (2) $h(d)(d')$ is in $\Delta I_2$ for all $d, d'$ in $\Delta I_1$, $A \subseteq \Delta I_\Sigma \cap \Sigma$, and (3) $h(d)(d')$ is in $\Delta I_2$ for all $(d, d')$ in $\Delta I_2$, $r \in \Delta I_\Sigma \cap \Sigma$. If there is a $\Sigma$-homomorphism from $I_1$ to $I_2$, we write $I_1 \rightarrow_{\Sigma} I_2$.

Let $T$ be a Horn-\text{$\mathcal{ALCHIF}$} TBox in normal form and $A$ an ABox consistent with $T$. A type for $T$ is a set $t \subseteq \text{sub}(T) \cap \Delta I_\Sigma$ such that $T \models \bigcap t \subseteq A$ implies $A \models t$ for all concept names $A$. For $a \in \text{ind}(A)$, let $\text{tp}_T(a)$ be the type of a relative to $T$. When $a \in \text{ind}(A)$, $t, t'$ are types for $T$, and $r$ is a role, we write

- $a \sim_{r, T} t$ if $T, A \models \exists r. \bigcap t(a)$ and $t$ is maximal with this condition, and
- $t \sim_{r, T} t'$ if $T \models \bigcap t \subseteq \exists r. \bigcap t'$ and $t'$ is maximal with this condition.

A path for $A$ and $T$ is a finite sequence $\pi = ar_1 \cdots ar_n$ such that $T \models \bigcap t \subseteq A$ implies $a \in \text{ind}(A)$ for all concept names $A$, and $t_1, \ldots, t_n$ types for $T$ such that

(i) $a \sim_{r_1} t_1$ and, if $\text{func}(r_0) \in T$, then there is no $b \in \text{ind}(A)$ such that $T, A \models r_0(a, b)$;

(ii) for every $1 \leq i < n$, we have $t_i \sim_{r_i} t_{i+1}$ and, if $\text{func}(r) \in T$, then $r_{i-1} \neq r_i$.

When $n > 0$, we use $\text{tail}(\pi)$ to denote $t_n$. Let Paths be the set of all paths for $T$ and $A$. The universal model $I_{\Sigma, T, A}$ of $T$ and $A$ is defined as follows:

$\Delta I_{\Sigma, T, A} = \text{Paths}$

$\text{A}^{I_{\Sigma, T, A}} = \{ a \in \text{ind}(A) \mid T, A \models A(a) \}$

$\{ \pi \in \Delta I \setminus \text{ind}(A) \mid T \models \bigcap \text{tail}(\pi) \subseteq A \}$

$\text{r}^{I_{\Sigma, T, A}} = \{ (a, b) \in \text{ind}(A)^2 \mid s(a, b) \in A, T \models s \subseteq r \} \cup \{ (\pi, \text{st}) \mid \text{st} \in \text{Paths} \text{ and } T \models s \subseteq r \} \cup \{ (\pi, \text{st}) \mid \text{st} \in \text{Paths} \text{ and } T \models s \subseteq r \}$

It is standard to prove that $I_{\Sigma, T, A}$ is indeed a model of $T$ and $A$ and that it is universal in the sense that for every model $I$
of $T$ and $A$, we have $I \rightarrow I_{T,A}$. Consequently, $T, A \models q(a)$ iff $I_{T,A} \models q(a)$, for all CQs $q(x)$ and tuples $a$ of individuals.

We also need universal models of a TBox $T$ and a type $t$, instead of an ABox. More precisely, we define $I_{T,t} = I_{T,A}$, where $A_t = \{ A(a) \mid A \in t \}$ for a fixed $a \in \mathbb{N}$.

3 Model-theoretic Characterization

We aim to provide a model-theoretic characterization of query entailment that will be the basis for our decision procedure later on. The first step towards this characterization consists in showing that non-entailment is always witnessed by tree-shaped ABoxes and tree-shaped CQs with at most one answer variable. Here, an ABox $A$ is tree-shaped if it does not contain an assertion of the form $r(a, a)$, the undirected graph $G_A = (\text{ind}(A), \{ \{a, b\} \mid r(a, b) \in A\})$ is a tree, and for any $a, b \in \text{ind}(A)$, $A$ contains at most one role assertion that involves both $a$ and $b$.

**Lemma 8** Let $T_1$ and $T_2$ be Horn-$\mathcal{ALCHIF}$ TBoxes with $T_1 \models_{I_{T_1}} T_2$. If $T_1 \not\models_{\Sigma, T_2}$, then there is a witness $(A, q, a)$ where $A$ and $q$ are tree-shaped and $|A| \leq 1$, i.e., $q$ has at most one answer variable. If $T_1 \models_{I_{T_2}} T_2$, then there is such a witness where additionally $q$ is an stCQ.

Our goal is to characterize query entailment in terms of homomorphisms between (universal) models. Homomorphisms are natural because answers to CQs are preserved under homomorphisms between (universal) models. Homomorphisms are partly preserved even under bounded homomorphisms if the bound is not smaller than the number of variables in the CQ.

Let $I_1, I_2$ be interpretations, $d \in \Delta_{T_1}^2$, and $n \geq 0$. We say that there is an $n$-bounded $\Sigma$-homomorphism from $I_1$ to $I_2$, written $I_1 \rightarrow_{\Sigma, n} I_2$, if for any subinterpretation $I'_1$ of $I_1$ with $|\Delta_{\Sigma}^2| \leq n$, we have $I'_1 \rightarrow_{\Sigma} I_2$. Moreover, we write $I_1 \rightarrow_{\Sigma, n} I_2$ if $I_1 \rightarrow_{\Sigma} I_2$ for any $n$. The following characterization follows from the definition of CQ entailment, Lemmas 8 and the connection between CQs and suitably bounded homomorphisms.

**Lemma 9** Let $T_1$ and $T_2$ be Horn-$\mathcal{ALCHIF}$ TBoxes with $T_1 \models_{I_{T_1}} T_2$. Then $T_1 \models_{I_{T_2}} T_2$ iff for all tree-shaped $\Gamma$-ABoxes $A$ consistent with $T_1$ and $T_2$, $\mathcal{I}_{T_1,A} \rightarrow_{\Sigma, \mathcal{I}_{T_2,A}} ^{\text{fin}} \mathcal{I}_{T_1,A}$. Ideally, we would like to use Lemma 9 as a basis for a decision procedure based on tree automata. To this end, it is useful that the ABox $A$ and models $I_{T_1,A}$ and $I_{T_2,A}$ in the lemma are tree-shaped. What is problematic is that Lemma 9 speaks about bounded homomorphisms, for any bound (corresponding to the bounded size of CQs), since it does not seem possible to verify such a condition using automata. We would thus like to replace bounded homomorphisms with unbounded ones, which does not compromise the characterization in the case of Horn-DLs without inverse roles [Lutz and Wolter, 2010; Botoeva et al., 2016c]. However, this is not true already for ELT TBoxes [Botoeva et al., 2016a].

**Example 10** Let $T_1 = \{ A \subseteq \exists s.B, B \subseteq \exists r.B \}$, $T_2 = \{ A \subseteq \exists s.B, B \subseteq \exists r.B \}$, $\Gamma = \{ A \}$, and $\Sigma = \{ r \}$. Then both $\mathcal{I}_{T_1,A}$ and $\mathcal{I}_{T_2,A}$ contain an infinite $r$-path; the $r$-path in $\mathcal{I}_{T_2,A}$ has a final element while the one in $\mathcal{I}_{T_1,A}$ does not. Hence $\mathcal{I}_{T_2,A} \not\rightarrow_{\Sigma} \mathcal{I}_{T_1,A}$, but $T_1 \models_{I_{T_2}} T_2$ (see Thm. 11 below).

We now show that it is possible to refine Lemma 9 so that it makes a much more careful statement in which bounded homomorphisms are partly replaced by unbounded ones. It is then possible to check the unbounded homomorphism part of the characterization using tree automata as desired, and to deal with bounded homomorphisms using a mosaic technique that “precompiles” relevant information about unbounded homomorphisms to be used in the automaton construction.

We start with introducing relevant notation. For a signature $\Sigma$, we use $\mathcal{I}_{\Sigma}^\text{con}$ to denote the restriction of the interpretation $I$ to those elements that can be reached from an ABox individual by traveling along $\Sigma$-roles (forwards or backwards). Tree-shaped interpretations are defined analogously to tree-shaped CQs (thus multi-edges are allowed). For a TBox $T$, an ABox $A$, and $a \in \text{ind}(A)$, we use $\mathcal{I}_{T,A}|_a$ to denote the subtree interpretation in the universal model $\mathcal{I}_{T,A}$ rooted at $a$. $\Sigma$-subtree in $\mathcal{I}_{T,A}$ is a maximal tree-shaped, $\Sigma$-connected sub-interpretation $I$ of $\mathcal{I}_{T,A}$ that does not comprise any ABox individuals. The root of $I$ is the (unique) element of $\Delta^2$ that can be reached from an ABox individual on a shortest path among all elements of $\Delta^2$. The refined characterization uses simulations instead of homomorphisms for the stCQ case because they are insensitive to multi-edges. Given a signature $\Sigma$ and two interpretations $I, J$, a $\Sigma$-simulation of $I$ in $J$ is a relation $s \subseteq \Delta^2 \times \Delta^2$ such that: (1) $(a, a) \in s$ for all $a \in \mathbb{N}$, (2) if $d \in A^2$ with $A \in \Sigma$ and $(d, e) \in s$, then $e \in A^2$, and (3) if $(d, d') \in s^r$ with $\Sigma$-role and $(d, e) \in \Sigma$, then there is some $e'$ with $(e, e') \in s^r$ and $(d', e') \in \Sigma$. We write $I \leq_{\Sigma} J$ if there is a $\Sigma$-simulation of $I$ in $J$.

**Theorem 11** Let $T_1$ and $T_2$ be Horn-$\mathcal{ALCHIF}$ TBoxes with $T_1 \models_{I_{T_1}} T_2$. Then $T_1 \models_{I_{T_2}} T_2$ iff for all tree-shaped $\Gamma$-ABoxes $A$ consistent with $T_1$ and $T_2$, and for all tree-shaped, finitely branching models $I_1$ and $I_2$, the following hold:

(1) $T_1, A \models_{\Sigma} \mathcal{I}_{T_1,A}^\text{con} \rightarrow_{\Sigma} I_1$;

(2) for all $\Sigma$-subtrees $I$ in $\mathcal{I}_{T_2,A}$, one of the following holds:

(a) $I \rightarrow_{\Sigma} I_1$;

(b) $I \rightarrow_{\Sigma, a} \mathcal{I}_{T_1, \mathcal{I}_{T_2,A}}$ for some $a \in \text{ind}(A)$.

Furthermore, $T_1 \models_{\Sigma} T_2$ iff $\mathcal{I}_{T_2,A}^\text{con} \not\rightarrow_{\Sigma} \mathcal{I}_{T_1,A}$.

4 Decidability and Complexity

We prove that, in Horn-$\mathcal{ALCHIF}$, CQ entailment can be decided in 2EXPTIME. By existing lower bounds, the former is thus 2EXPTIME-complete in all fragments of Horn-$\mathcal{ALCHIF}$ that contain $\mathcal{ELT}$ or Horn-$\mathcal{ALC}$. Moreover, stCQ entailment in Horn-$\mathcal{ALCHIF}$ and deductive entailment in $\mathcal{ELHIF}$ can also be decided in 2EXPTIME. We establish a ConNEXPTIME lower bound and leave the precise complexity open.

To obtain the upper bounds, we use a combination of tree automata and mosaics to implement the characterization in Theorem 11. We start with a mosaic-based decision procedure for Condition (2b). Note that a $\Sigma$-subtree $I$ in $\mathcal{I}_{T_2,A}$ can be uniquely identified by the type $t_2$ of its root. It therefore suffices to show the following.
Theorem 12. Given two Horn-ALCHIF TBoxes $T_1$ and $T_2$ and types $t_i$ for $T_i$, $i \in \{1, 2\}$, it can be decided in time $2^{2^{|T_2|+|T_1|}}$ whether $\mathcal{I}_{T_2,t_2}^{\cong} \rightarrow_{\Sigma} \mathcal{I}_{T_1,t_1}$, $p$ a polynomial.

Although we cannot get rid of bounded homomorphisms in Theorem 11, a central idea for applying a mosaic approach to prove Theorem 11 is to first replace homomorphisms with unbounded ones. To make this possible, we replace $\mathcal{I}_{T_1,t_1}$ with a suitable class of interpretations used as targets for the unbounded homomorphisms.

To illustrate, consider Example 10 and let $t_1 = t_2 = \{B\}$. The difference between $\mathcal{I}_{T_2,t_2} \rightarrow_{\Sigma} \mathcal{I}_{T_1,t_1}$ and $\mathcal{I}_{T_2,t_2} \rightarrow_{\Sigma} \mathcal{I}_{T_1,t_1}$ is that unbounded homomorphisms fail once they reach the root of $\mathcal{I}_{T_1,t_1}$, while bounded homomorphisms can, depending on the bound, map the root of $\mathcal{I}_{T_2,t_2}$ deeper and deeper into $\mathcal{I}_{T_1,t_1}$, thus never reaching its root. The latter is possible because $\mathcal{I}_{T_1,t_1}$ is regular in the sense that any two elements which have the same type root isomorphic subtrees. This is of course not only true in this example, but by construction in any universal model. To transition back from bounded to unbounded homomorphisms, we replace $\mathcal{I}_{T_1,t_1}$ with a class of (finite and infinite) interpretations that can be seen as a “backwards regularization” of $\mathcal{I}_{T_1,t_1}$. In our concrete example, we would include an interpretation $\mathcal{I}$ where that predecessor has a predecessor, and $\mathcal{I}$ is a tuple $(t_1, t_2)$, the universal model.

To unbounded homomorphisms, we replace $\mathcal{I}_{T_1,t_1}$ to be homomorphically embedded into the neighborhood. Formally, a mosaic for $t_1$ is a tuple $M = (t_-, t, t, S, t)$ such that $(t_-, t, t, S) = n_1^{\Delta^r_{t,t_1}}(d)$ for some $d \in \Delta^r_{T_1,t_1}$ and $t: \{\hat{t}_1, t\} \cup S \rightarrow 2^{\mathcal{I}_{T_1,t_1}}$ satisfies the following condition: (M) For all $\hat{t} \in \ell(t)$ we have $\hat{t} \in \Sigma \subseteq t$ and, for all $\hat{t} \in \ell(t_2)$, $r \in \mathcal{R}(t_2)$ where $\hat{t} \in \ell(t_2)$, one of the following holds for $s = \mathcal{R}(t_2) \mid t_2 = r \subseteq s.$

To ease notation, we use $t_M$ to denote $t$, $\rho_M$ to denote $\rho$, and likewise for the other components of a mosaic $M$. Let $M$ be the set of all mosaics for $t_1$ and $M' \subseteq M$. An $M'$ is good in $M'$ if the following conditions are satisfied:

1. For each $(\rho, t) \in S_M$, there is an $N \in M'$ such that $(M, \rho, t) = (N, \rho_M, t_N)$, $\ell_M(\rho, t) = \ell_N(t_N)$, and $\ell_M(t_M) = \ell_N(t_N)$.
2. If $\rho = \rho_M$ and $M' \subseteq M$.

Let $M_0, M_1, \ldots$ be the sequence obtained by starting with $M_0 = M$ and defining $M_{i+1}$ to be $M_i$ when all mosaics that are not good in $M_i$ have been removed. Assume that $M_p$ is where the sequence stabilizes.

Lemma 14. Let $t_i \in \mathcal{I}(T_1)$ for $i \in \{1, 2\}$. Then there is a $J \in \mathcal{I}_{\Sigma}^{\cong} \rightarrow_{\Sigma} J'$ if $M_p$ contains a mosaic $M$ with $t_2 \in \ell_M(t_M)$.

Since there are at most $2^{T_2}$ mosaics for $t_1$, we obtain the desired Theorem 12.

We now develop the decision procedure for CQ and stCQ entailment in Horn-ALCHIF, based on Theorems 11 and 12. Our main tool are alternating two-way tree automata with counting (2ATA$_c$), an extension of alternating tree automata over unranked trees [Gradel and Walukiewicz, 1999] with the ability to count. A tree is a non-empty (potentially infinite)
set of words $T \subseteq \{N \setminus \{0\}\}^*$ closed under prefixes. We assume that trees are finitely branching, i.e., for every $w \in T$, the set $\{ i \mid w \cdot i \in T \}$ is finite. For any $w \in \{N \setminus \{0\}\}$, we set $w \cdot 0 := w$. If $w = n_0n_1 \cdots n_k, k \geq 0$, we set $w \cdot -1 := n_0 \cdots n_{k-1}$. For an alphabet $\Theta$, a $\Theta$-labeled tree is a pair $(T, L)$ with $T$ a tree and $L : T \rightarrow \Theta$ a node labeling function.

A 2ATA $A$ is a tuple $\mathfrak{A} = (Q, \Theta, q_0, \delta, \Omega)$ where $Q$ is a finite set of states, $\Theta$ is the input alphabet, $q_0 \in Q$ is the initial state, $\delta$ is a transition function, and $\Omega : Q \rightarrow N$ is a priority function. The transition function $\delta$ maps every state $q$ and input letter $a \in \Theta$ to a positive Boolean formula $\delta(q, a)$ over the truth constants true and false and transition atoms of the form $q, \langle - \rangle q, \langle q \rangle \Theta, \bot$, $\top$, and $\tau$. Informally, a transition $\delta(q)$ expresses that a copy of $\mathfrak{A}$ is sent to the current node in state $q$; $\langle - \rangle q$ means that a copy is sent in state $q$ to the predecessor node, which is required to exist; $\top$ means that the predecessor node is not required to exist; $\langle q \rangle \Theta$ (resp., $\tau$) means that a copy of $\mathfrak{A}$ is sent to $n$ (resp., to all but $n$) successors. The semantics of 2ATA$\mathfrak{A}$ is given in terms of runs as usual, please see the appendix. We use $L(\mathfrak{A})$ to denote the set of trees accepted by $\mathfrak{A}$. It is standard to verify closure of 2ATA$\mathfrak{A}$ under intersection. The following is obtained via reduction to standard alternating parity tree automata [Vardi 1998].

**Theorem 15.** The emptiness problem for 2ATA$\mathfrak{A}$ can be solved in time exponential in the number of states.

Let $T_1, T_2$ be Horn-ALCH\(\mathcal{I}, \mathcal{F}\) TBoxes and $\Gamma, \Sigma$ signatures. We aim to show that one can construct a 2ATA$\mathfrak{A}$ such that $L(\mathfrak{A}) = \emptyset$ iff $T_1 \not\models^CQ T_2$. In fact, $\mathfrak{A}$ is the intersection of four 2ATA$\mathfrak{A}$, $\mathfrak{A}_1, \mathfrak{A}_2, \mathfrak{A}_3, \mathfrak{A}_4$. They run over $\Theta$-labeled trees with $\Theta = 2\Theta_0 \times 2\Theta_1 \times 2\Theta_2$, where $\Theta_0 = \Gamma \cup \{ r \mid r \in \Gamma \}$ and $\Theta_i = \Sigma \cup \{ r \mid r \in \Sigma_i \}$ for $i = 0, 1, 2$. For a $\Theta$-labeled tree $(T, L)$, we use $L_i, i \in \{0, 1, 2\}$ to refer to the $i$-th component of $L$, that is, $L(n) = (L_0(n), L_1(n), L_2(n))$, for all $n \in T$. The component $L_0$ represents a (possibly infinite) ABox $A = \{ A(n) \mid A \in L_0(n) \} \cup \{ r(n, \cdot, -1, n) \mid n \neq \varepsilon, r \in L_0(n) \}$, where $r(a,b)$ is identified with $r(b,a)$. The 2ATA$\mathfrak{A}_1$ accepts a $\Theta$-labeled tree $(T, L)$ if $\mathfrak{A}$ is finite, tree-shaped (and thus connected) and includes the root of $T$, and it is straightforward to construct.

Components $L_1, L_2$ give rise to interpretations $I_1 = (T, I_1)$ and $I_2 = (ind(\mathfrak{A}), I_2)$, where for $i \in \{1, 2\}$,

$A^{I_1} = \{ a \in L_i(n) \}$

$A^{I_2} = \{ (n, n \cdot \varepsilon, -1) \mid r \in L_i(n) \} \cup \{ (n \cdot \varepsilon, -1, n) \mid r \in L_i(n) \}$

$\mathfrak{A}_2$ verifies that $I_1$ is a model of $\mathfrak{A}$ and $\mathfrak{T}_1$, which is standard, too. $\mathfrak{A}_3$ verifies that $A$ is consistent with $T_2$, and $\mathfrak{A}_2$ is $T_2, A$ restricted to $\text{ind}(\mathfrak{A})$. This involves computing the type of an ABox element without having access to the anonymous (that is: non-\text{\texttt{ABox}}) part of $I_2, A$, using a characterization of ABox entailments [Bienvenu et al., 2013] in terms of derivation trees. Finally, $\mathfrak{A}_4$ verifies that either (1) or (2) from Theorem [17] is not satisfied. For (1), $\mathfrak{A}_4$ sends a copy of itself to every tree $T$ starting at an ABox element in $I_2, A$, and attempts to show that $T$ cannot be homomorphically embedded into a corresponding tree in $I_1$. This attempt is successful if either incompatible types are found in the root or, recursively, there is some successor of the current type in $T, A$ that cannot be mapped to any neighbor in $I_1$. Since the anonymous part of $T, A$ is not explicit in the input, the current type is stored in the states, and the generating relation $t \leadsto t'$ is “hard-coded” into the transition function. For Condition (2a), $\mathfrak{A}_4$ non-deterministically guesses a $\Sigma$-subtree $I$ and proceeds as in (1); Condition (2b) is verified based on Theorem [12] by pre-computing $\delta^{b_\perp}$. Thus the number of states of $\mathfrak{A}_4$ is exponential in $T_2$ (because of the types) but only polynomial in $\mathfrak{T}_1$. Automata $\mathfrak{A}_1, \mathfrak{A}_2, \mathfrak{A}_3$ have polynomially many states.

In the special case of stCQ entailment, we simply replace $\mathfrak{A}_4$ with a 2ATA$\mathfrak{A}_4$ that refutes the simulation condition of Theorem [11] analogously to how $\mathfrak{A}_4$ refutes Condition (1).

To obtain the desired upper complexity bounds for query and deductive entailment, we observe that, in both cases, $\mathfrak{A}$ can be constructed in time polynomial in $\mathfrak{T}_1$ and exponential in $\mathfrak{T}_2$, and the emptiness check adds an exponential blowup (Theorem [15]). For deductive entailment, we use the reduction to stCQ entailment (Lemma [7]).

**Theorem 16.** In Horn-ALCH\(\mathcal{I}, \mathcal{F}\), the following problems can be decided in time $2^{2^{2^{\log_2|T_2|}+2^{2^{\log_2|T_1|}}}}$, $p$ a polynomial: $\langle \Gamma, \Sigma \rangle$-CQ entailment, $\langle \Gamma, \Sigma \rangle$-CQ inseparability, and $\langle \Gamma, \Sigma \rangle$-CQ conservative extensions. The same holds for $\Sigma$-deductive entailment, $\Sigma$-deductive inseparability, and $\Sigma$-deductive conservative extensions in $\mathcal{ELH} \mathcal{I}, \mathcal{F}_1$.

Matching lower bounds for all problems except deductive entailment are provided by [Botoeva et al., 2016c]. They hold even in the case where $\Gamma = \Sigma$.

**Corollary 17.** In any fragment of Horn-ALCH\(\mathcal{I}, \mathcal{F}\) that contains $\mathcal{EL}$ or Horn-\text{\texttt{AC}}, the following problems are $\text{\texttt{EXPTIME}}$-complete: $\langle \Gamma, \Sigma \rangle$-CQ entailment, $\langle \Gamma, \Sigma \rangle$-CQ inseparability, and $\langle \Gamma, \Sigma \rangle$-CQ conservative extensions.

In the description logic $\mathcal{EL}$, which is $\mathcal{EL}$ without inverse roles, deductive conservative extensions and deductive $\Sigma$-entailment are $\text{\texttt{EXPTIME}}$-complete [Lutz and Wolter, 2010]. This raises the question whether the upper bound for deductive entailment reported in Theorem [15] is tight. While we leave this question open, we observe that the transition from $\mathcal{EL}$ to $\mathcal{EL}$ does increase the complexity of deductive conservative extensions and related problems to at least $\text{\texttt{EXPTIME}}$. We consider this a surprising result since in reasoning problems that are not defined in terms of conjunctive queries, adding inverse roles does typically not result in an increase of complexity. The following is established by a non-trivial reduction of a tiling problem.

**Theorem 18.** In any DL between $\mathcal{EL}$ and $\mathcal{ELH} \mathcal{I}, \mathcal{F}_1$, deductive conservative extensions, deductive $\Sigma$-entailment, and deductive $\Sigma$-inseparability are $\text{\texttt{EXPTIME}}$-hard.

5 Conclusion

As future work, it would be interesting to close the gap in complexity between $\text{\texttt{CONEXPTIME}}$ and $\text{\texttt{2EXPTIME}}$ for deductive entailment in $\mathcal{EL}$ and $\mathcal{ELH} \mathcal{I}, \mathcal{F}_1$. Furthermore, it would be interesting to extend the results to ontology languages from the family of DataLog+- (aka existential rules), in particular to frontier-guarded TGDs.
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References


A Proofs for Section 2

A.1 Additional Definitions and Properties

If an interpretation $I$ is a common model of a TBox $T$ and ABox $A$, then we also write $I \models (T, A)$ and call $I$ a model of $(T, A)$.

The following is standard to prove.

Lemma 19 For every Horn-$\text{ALCHIF}$ TBox $T$ in normal form and ABox $A$ consistent with $T$, the following hold:

1. $I_{T, A} \models (T, A)$.
2. For all models $I$ of $(T, A)$, we have $I_{T,A} \rightarrow I$.
3. For all types $t, t'$ for $T$ with $t \subseteq t'$, we have $I_{T,t} \rightarrow I_{T,t'}$.
4. $T, A \models q(a)$ iff $I_{T, A} \models q(a)$, for all CQs $q(x)$ and tuples $a$ of individuals.

In the following subsections, we need to deal with the question whether any $\Gamma$-ABox that is inconsistent with some TBox $T_2$ is inconsistent with another TBox $T_1$. We say that $T_1$ $\Gamma$-inconsistency entails $T_2$, written $T_1 \models \beta T_2$, if for all $\Gamma$-ABoxes $A$: if $A$ is inconsistent with $T_2$, then $A$ is inconsistent with $T_1$.

A.2 Proof of Lemma 4

Lemma 4 CQ entailment with inconsistent ABoxes can be decided in polynomial time given access to oracles deciding CQ entailment and CQ evaluation.

To prove Lemma 4 we proceed in two steps: first, we show how to deal with inconsistency entailment as defined in Section A.1; second we show how to use this type of entailment to deal with CQ entailment with inconsistent ABoxes.

For the first step, we can reduce inconsistency entailment to CQ entailment because, if $A$ is inconsistent with $T$, then either (a) $I_{T, A}$ contains a $B$-instance for some $B$ with $B \subseteq \bot \subseteq T$, or (b) $A$ contains a “fork” $\{r(a, b), r(a, c)\}$ that is prohibited by $\text{func}(r) \subseteq T$. Write $T_1 \models \beta T_2$ if for all $\Gamma$-ABoxes $A = \{r(a, b), r(a, c)\}$: if $A$ is inconsistent with $T_2$, then also with $T_1$. Thus, if we have a witness ABox $A$ for $T_1 \not\models \beta T_2$, then $A$ is inconsistent with $T_2$ either by Case (a) – we which we can detect via CQ entailment if we allow a fresh concept name in the CQ and slightly modify the TBoxes – or by Case (b), which implies $T_1 \not\models \beta T_2$.

Lemma 20 Let $\Gamma$ be a signature and let $T_1$ and $T_2$ be Horn-$\text{ALCHIF}$ TBoxes. Furthermore, let $A$ be a fresh concept name and $T_1^A$ be obtained from $T_1$ by replacing each occurrence of $\bot$ with $A$ and adding the axioms $\exists s.A \subseteq A$ and $\exists s-A \subseteq A$ for every role $s$ occurring in $T_r$, for $i = 1, 2$. Then the following are equivalent.

1. $T_1 \models \beta T_2$
2. $T_1^A \models \text{CQ}_{\Gamma, A} T_2^A$ and $T_1 \not\models \text{fork} T_2$
3. $T_2^A \models \text{stCQ}_{\Gamma, A} T_1^A$ and $T_1 \not\models \text{fork} T_2$

Note that we only need the equivalence between (1) and (2) to prove Lemma 4. However, we will need (3) later to prove Lemma 4. Indeed, (2) and (3) are obviously equivalent given the primitive query signature $\{A\}$ and the propagation of $A$ throughout the $T_i^A$.

Proof.

“1 $\Rightarrow$ 2”. We prove the contrapositive. Assume $T_1^A \not\models \text{CQ}_{\Gamma, A} T_2^A$ or $T_1 \not\models \text{fork} T_2$. In case $T_1 \not\models \text{fork} T_2$, every witness ABox is a witness for $T_1 \not\models \text{fork} T_2$ too.

In case $T_1^A \not\models \text{CQ}_{\Gamma, A} T_2^A$ is violated, consider a witness $(A, q, a)$. Since $A$ is the only symbol allowed in $q$, all atoms of $q$ have the form $A(z)$ for arbitrary variables $z$. If $q$ consists of several atoms, then it is disconnected and we can omit all but one atom from $q$ and still have a witness (see also proof of Lemma 3 Property d). Hence we can assume w.l.o.g. that $q$ is of the form (i) $q(x) = A(x)$ or (ii) $q() = \exists y A(y)$ and, furthermore, that $A$ and thus the universal models $I_{T_i, A}$ are connected. (Due to the “propagation” of $A$ in the $T_i$, we can even assume that $q$ is of the form (i) only, but that does not matter in the following argumentation.) We now have:

- $A$ is inconsistent with $T_2$.
- Assume to the contrary that $A$ is consistent with $T_2$ and consider the universal model $I_{T_2, A}$ for $T_2$ and $A$ (Section 2.3). Clearly, for all domain elements $d$ of $I_{T_2, A}$, we have $T_2 \not\models \bigcap \{\text{tp}_{T_2, A}(d) \subseteq A\}$. Since $A$ is fresh and by the definition of $T_2^A$ we get $T_2^A \not\models \bigcap \{\text{tp}_{T_2, A}(d) \subseteq A\}$.

Now Lemma 19(1) for $T_2^A$ implies that $I_{T_2, A} \not\models (T_2, A)$; hence $I_{T_2, A}$ satisfies all axioms in $T_2^A$ that have been taken over from $T_2$ without modification, i.e., all axioms that are not of the form $B \subseteq A$. But axioms of the latter form are also satisfied because $T_2^A \not\models \bigcap \{\text{tp}_{T_2, A}(d) \subseteq A\}$ for every domain element $d$. Hence $I_{T_2, A} \models (A, q, a)$. Now, since $I_{T_2, A}$ has no $A$-instance, we cannot have $T_2^A, A \models q(a)$ for any $\{A\}$-query $q$, contradicting the assumption that $(A, q, a)$ is a witness.

- $A$ is consistent with $T_1$.
- Since $(A, q, a)$ is a witness, we have $I_{T_1, A} \not\models q(a)$ by Lemma 19(4). Due to the additional axioms in the definition of $T_1^A$, which “propagate” $A$ into every domain element of the connected (see above) universal model $I_{T_1^A, A}$, we have $T_1^A \not\models \bigcap \{\text{tp}_{T_1^A, A}(d) \subseteq A\}$ for all domain elements $d$. Since $A$ is fresh, we have $T_1 \not\models \bigcap \{\text{tp}_{T_1^A, A}(d) \subseteq A\}$.

With the same reasoning as above, we get $I_{T_1^A, A} \models (T_1, A)$; hence $A$ is consistent with $T_1$.

Consequently $T_1 \not\models \beta T_2$, as desired.

“2 $\Rightarrow$ 3”. This is immediate because $T_1^A \models \text{CQ}_{\Gamma, A} T_2^A$ implies $T_1^A \models \text{stCQ}_{\Gamma, A} T_2^A$.

“$3 \Rightarrow 1$”. We prove the contrapositive. Assume $T_1 \not\models \beta T_2$, i.e., there is a $\Gamma$-ABox $A$ that is inconsistent with $T_2$ but consistent with $T_1$. We need to show that $T_1^A \not\models \text{CQ}_{\Gamma, A} T_2^A$ or $T_1 \not\models \text{fork} T_2$.
From \(A\) being inconsistent with \(T_2\), we first conclude that one of the following two properties must hold.

(i) There is some \(d \in B^{T_2-A}\) with \(B \subseteq \bot \in T_2\).

(ii) \(A\) contains a “fork” \(A' = \{r(a, b), r(a, c)\}\) such that \(A'\) is inconsistent with \(T_2\).

Indeed, if neither (i) nor (ii) holds, then we have \(\mathcal{I}_{T_2-A} \models (T_2, A)\), contradicting the inconsistency of \(A\) with \(T_2\): First, \(\mathcal{I}_{T_2-A} \models A\) follows directly from the construction of \(\mathcal{I}_{T_2-A}\). Second, \(\mathcal{I}_{T_2-A} \models T_2\) can be shown analogously to the (omitted) standard proof of Lemma 19, via a case distinction over the axioms in \(T_2\), using “not (i)” and “not (ii)” instead of the assumption that \(A\) is consistent with \(T_2\).

Now first assume that (ii) holds. Since \(A\) is consistent with \(T_1\), so is \(A'\). Hence \(\mathcal{I}_1 \not\models \top\).

In case (ii) does not hold, (i) must hold. To show that \(\mathcal{I}_{T_2-A} \not\models \top\), consider the stCQ \(q = A(x)\) and some \(a \in \text{ind}(A)\) to which the element \(d\) from (i) is connected in \(\mathcal{I}_{T_2-A}\), i.e., if \(d \in \text{ind}(A)\), then choose \(a = d\); otherwise choose \(a\) such that \(d\) is in the subtree \(\mathcal{I}_{T_2-A}|_a\). We then have:

- \(A\) is consistent with \(T_2^A\):

Since \(T_2^A\) does not contain \(\bot\) and \(A\) does not contain forks as in (ii), \(A\) is consistent with \(T_2^A\), and \(\mathcal{I}_{T_2-A} \not\models \top\).

- \(A\) is consistent with \(T_1^A\):

It is not difficult to see that \(\mathcal{I}_{T_1-A} \models (T_1^A, A)\): From \(\mathcal{I}_{T_1-A} \models (T_1, A)\), it follows that \(\mathcal{I}_{T_1-A} \models A\) and \(\mathcal{I}_{T_1-A} \models (T_1^A, A)\) shares with \(T_1\). The modified axioms \(B \subseteq A\) with \(B \subseteq \bot \in T_1\) are satisfied, too, because \(\mathcal{I}_{T_1-A}\) cannot have any \(B\)-instances. Finally, the additional propagation axioms are satisfied because \(\mathcal{I}_{T_1-A}\) has no \(A\)-instance as \(A\) is fresh.

- \(T_2^A, A \models q(a)\):

Due to (i), we have \(\mathcal{I}_{T_2-A} \models \exists y B(y)\) for some \(B \subseteq \bot \in T_2\). Hence \(\mathcal{I}_{T_2-A} \models \exists y B(y)\), which follows from the construction of both universal models (in fact the only difference between \(\mathcal{I}_{T_2^A-A}\) and \(\mathcal{I}_{T_2-A}\) is that some domain elements of \(\mathcal{I}_{T_2^A-A}\) may be \(A\)-instances). Hence \(\mathcal{I}_{T_2^A-A}\) has a \(B\)-instance in the subtree \(\mathcal{I}_{T_2-A}|_a\) and thus, by construction, an \(A\)-instance. By the “propagation” of \(A\) in \(T_2^A\), we have that \(a\) is an instance of \(A\) in \(\mathcal{I}_{T_2^A-A}\); hence \(\mathcal{I}_{T_2^A-A} \models A(a) = q\).

- \(T_1^A, A \not\models q(a)\):

Follows from \(\mathcal{I}_{T_1-A} \models (T_1^A, A)\) (as shown above) and \(\mathcal{I}_{T_1-A} \not\models q(a)\) (given the lack of \(A\)-instances).

**Proposition 21** Fork entailment \(T_1 \models \text{fork} T_2\) can be (Turing) reduced in polynomial time to stCQ evaluation.

**Proof.** Perform \(2|\Gamma|\) many ABox consistency checks by evaluating the stCQ \(A(a)\) on both \(T_1\), where \(A\) is a concept name that does not occur in any of the \(T_i\).

For the second step, we can now reduce CQ entailment with inconsistent ABoxes to the disjunction of our original notion of CQ entailment and inconsistency entailment. We need an additional notion: Given a TBox \(T\) and signatures \(\Gamma, \Sigma\), we say that \(T\) is \((\Gamma, \Sigma)\)-universal if

\[(\ast)\] for all \(\Gamma\)-ABoxes \(A\) and \(\Sigma\)-CQs \(q(x)\) and all tuples \(a \subseteq \text{ind}(A)\) with \(|a| = |x|\), we have \(T, A \models q(a)\).

**Lemma 22** Let \(\Gamma, \Sigma\) be signatures and let \(T_1\) and \(T_2\) be Horn-ALCHIF TBoxes. Then \(T_1 \models (\Gamma, \Sigma)\)-CQ entails \(T_2\) with inconsistent ABoxes if one of the two following conditions holds.

1. \(T_1 \models \Gamma \bowtie \Sigma T_2\) and \(T_1 \models \top T_2\)
2. \(T_1\) is \((\Gamma, \Sigma)\)-universal.

**Proof.** We prove both implications via contraposition.

\(\Rightarrow\). Assume (1) and (2) are both false, i.e., \(T_1\) is not \((\Gamma, \Sigma)\)-universal and either (a) \(T_1 \not\models \bowtie \Sigma T_2\) or (b) \(T_1 \not\models \top T_2\). In case (a), \(T_1\) trivially does not \((\Gamma, \Sigma)\)-CQ entail \(T_2\) with inconsistent ABoxes. In case (b), consider a witness \(\Gamma\)-ABox \(A\). Since \(T_1\) is not \((\Gamma, \Sigma)\)-universal, there is a \(\Gamma\)-ABox \(A'\), a \(\Sigma\)-CQ \(q(x)\) and a tuple \(a \subseteq \text{ind}(A)\) with \(|a| = |x|\) such that \(T_1, A' \not\models q(a)\). We assume w.l.o.g. that \(A\) and \(A'\) use distinct sets of individuals. We set \(A'' = A \cup A'\) and have:

- \(T_2, A'' \models q(a)\) because \(A\) is inconsistent with \(T_2\) and so is \(A''\).
- \(T_1, A'' \not\models q(a)\): let \(J\) be the disjoint union of the universal model \(\mathcal{I}_{T_1-A}\) and the model \(J\) witnessing \(T_1, A' \not\models q(a)\). Clearly \(J \not\models (T_1, A'')\) but \(J \not\models q(a)\).

Hence \(T_1\) does not \((\Gamma, \Sigma)\)-CQ entail \(T_2\) with inconsistent ABoxes, as desired.

\(\Leftarrow\). Assume \(T_1\) does not \((\Gamma, \Sigma)\)-CQ entail \(T_2\) with inconsistent ABoxes and consider a witness \((A, q, a)\). Then it is immediate that (2) does not hold. Furthermore, if \(A\) is consistent with both \(T_1\) and \(T_2\), then \(T_1 \not\models \bowtie \Sigma T_2\). Otherwise \(A\) must be inconsistent with \(T_2\) but consistent with \(T_1\); hence \(T_1 \not\models \top T_2\). Therefore (1) does not hold either.
Lemma 4 is now a direct consequence of Lemmas 22 and 20 and Propositions 21 and 23.

A.3 Proof of Lemma 7

Lemma 7 In ELHIF, deductive entailment can be decided in polynomial time given access to oracles for stCQ entailment and stCQ evaluation.

Lemma 7 is an immediate consequence of the following lemma because the additional $T_1 \not\models_{\Sigma;\Sigma} T_2$ can be reduced to stCQ entailment and stCQ evaluation via Lemma 20.

Lemma 24 Let $\Sigma$ be a signature and $T_1, T_2 \in \mathbb{ELHIF}_\Sigma$ TBoxes such that $T_1 \not\models_{\Sigma;\Sigma} T_2$. Then

$$T_1 \not\models_{\Sigma;\Sigma} T_2 \iff T_1 \not\models_{\Sigma;\Sigma} T_2 \text{ and } T_1 \not\models T_2.$$

Proof. We prove both implications via contraposition.

\(\Leftarrow\). We assume that $T_1 \not\models_{\Sigma;\Sigma} T_2$. In case this is witnessed by a $\Sigma$-FA $\text{func}(r)$, we immediately get a witness $\Sigma$-ABox $= \{r(a, b), r(a, c)\}$ for $T_1 \not\models\not\models T_2$ and are done.

Otherwise, $T_1$ contains all $\Sigma$-FAs from $T_2$, and there is a witness $\Sigma$-CI $C \subseteq D$ (witness RI s are excluded by the assumption $T_1 \not\models_{\Sigma;\Sigma} T_2$). Since $\mathbb{ELI}_1$-concepts that contain $\bot$ are equivalent to $\bot$, the left-hand side $C$ cannot contain $\bot$ (i.e., an $\mathbb{ELI}$ concept) and, if $D$ does, then $C \subseteq \bot$ is a witness. We show that such witnesses give rise to either a witness $A_C$ for $T_1 \not\models\not\models T_2$ or a witness $(A_C, q_d, a)$ for $T_1 \not\models_{\Sigma;\Sigma} T_2$ with $q_d(x)$ an stCQ.

We first consider the case that there is a witness $C \subseteq \bot$ with $C$ an $\mathbb{ELI}$ concept. We can construct from $C$ in the obvious way a tree-shaped $\Sigma$-ABox $A_C$ and root $a$: $A$ reflects the tree structure of $C$; however, to respect the $\Sigma$-FAs in $T_1$ (and thus those in $T_2$), we need to merge subtrees of all nodes that are $r$-neighbors of the same node, whenever $\text{func}(r) \in T_1$. Consider the universal model $I_{T_2, A_C}$ and observe that $a \in C^{T_2, A_C}$ from the construction of $I_{T_2, A_C}$. Since $T_2 \models D \subseteq \bot$, we have that $I_{T_1, A_C}$ is not a model of $T_2$. Hence, by the contraposition of Lemma 19, $A_C$ is inconsistent with $T_2$. On the other hand, since $T_1 \not\models C \subseteq \bot$, there is a model $I \models T_1$ and an instance $d \in C^2$. We can turn $I$ into a model of $A_C$ by interpreting the ABox individuals accordingly ("partial" unraveling might be necessary to ensure that the standard name assumption is respected), witnessing the consistency of $A$ with $T_1$. We thus have $T_1 \not\models\not\models T_2$ and are done.

In the second case, all witnesses $C \subseteq D$ consist solely of $\mathbb{ELI}$ concepts $C$. We construct the same ABox $A_C$ with root $a$ from $C$ and transform $D$ into a $\Sigma$-stCQ $q_d(x)$ with a single answer variable that represents the tree shape of $D$. Now $(A_C, q_d, a)$ is a witness to $T_1 \not\models_{\Sigma;\Sigma} T_2$ for the following reasons.

- $A_C$ is consistent with $T_2$: since $C \subseteq \bot$ is not a witness to $T_1 \not\models_{\Sigma;\Sigma} T_2$, there must be a model $I \models T_2$ with $d \in C^2$. We claim that we can turn $I$ into a model of $A_C$ by interpreting the ABox individuals without violating the standard name assumption. If we assume to the contrary that this is not possible, then there are subconcepts $C_1, \ldots, C_n$ of $C$ corresponding to subtrees that have been merged in the construction of $A_C$, such that $T_2 \models C_1 \land \cdots \land C_n \subseteq \bot$. However, $T_1 \not\models C_1 \land \cdots \land C_n \subseteq \bot$ because $A_C$ is consistent with $T_1$, as shown previously. Hence $C_1 \land \cdots \land C_n \subseteq \bot$ would be a witness to $T_1 \not\models_{\Sigma;\Sigma} T_2$, which we have ruled out – a contradiction.

- $T_2, A_C \models q_d(a)$, witnessed by $I_{T_2, A_C}$, together with $a \in C^{T_2, A_C}$ and $T_2 \models D$.

- $T_1, A_C \not\models q_d(a)$: take a model $I$ witnessing $T_1 \not\models C \subseteq D$ and an element $d \in C^2 \setminus D^2$. As in the previous case, we can turn $I$ into a model $J$ of $A_C$ by interpreting the ABox individuals (again involving unraveling if necessary), obtaining $J \not\models q_d(a)$.

\(\Rightarrow\). Assume $T_1 \not\models_{\Sigma;\Sigma} T_2$ or $T_1 \not\models T_2$. In case $T_1 \not\models T_2$, consider a witness $\Sigma$-Box $A$ and assume w.l.o.g. that $A$ is tree-shaped. Let $a \in \text{ind}(A)$ be its root. We can assume that $T_1$ contains all $\Sigma$-FAs from $T_2$ (otherwise $T_1 \not\models_{\Sigma;\Sigma} T_2$ and we are done). We turn $I$ into a $\Sigma$-ELHIF concept $C_A$ in the obvious way. Then $C_A \subseteq \bot$ is a witness to $T_1 \not\models_{\Sigma;\Sigma} T_2$.

- $T_2 \models C_A \subseteq \bot$ because, if there were a model $I$ of $T_2$ with $d \in C_A$, we could turn it into a model of $(T_2, A)$ by interpreting the ABox individuals accordingly (possibly involving partial unraveling as above), which would contradict the assumption that $A$ is a witness to $T_1 \not\models\not\models T_2$.

- $T_1 \not\models C_A \subseteq \bot$, witnessed by $I_{T_1, A}$.

In case $T_1 \not\models_{\Sigma;\Sigma} T_2$, by Lemma 6 there is a witness $(A, q, a)$ with $A$ tree-shaped and $q$ a $\Sigma$-stCQ with exactly one answer variable. We construct $C_A$ as above and another $\Sigma$-ELHIF concept $D_q$ from $q$ in the obvious way. It can be shown analogously to the previous case that $C_A \subseteq D_q$ is a witness to $T_1 \not\models_{\Sigma;\Sigma} T_2$.

B.1 Unraveling ABoxes

To obtain tree-shaped ABoxes or CQs, we use unraveling, which needs to be more cautious in the presence of inverse roles and functionality. In particular, we need to ensure that, whenever a role is functional in an ABox, then so it is in its unraveling. We define an unraveling for Horn-ALCHIF similar to the one for Horn-ALCIF in [Lutz and Wolter 2012].

Let $A$ be an ABox. The unraveling $U_A^\Delta$ of $A$ at an individual $a \in \text{ind}(A)$ is the following ABox:

- $\text{ind}(U_A^\Delta)$ is the set of sequences $b_0 r_0 b_1 \cdots r_n-1 b_n$ with $n \geq 0$, where $b_0 = a$, $b_i \in \text{ind}(A)$ for all $0 \leq i \leq n$, $r_i(b_i, b_{i+1}) \in A$ for all $0 \leq i < n$, and $(b_{i-1}, r_{i-1}) \neq (b_{i+1}, r_i)$ for all $0 \leq i < n$.

The assumption that $A$ is consistent with $T$ is not needed for the construction of $I_{T, A}$, only for the proof of Lemma 19.
(b_{i+1}, r_i)$ (the latter inequality is needed to ensure preservation of functionalit).  

- The concept assertions in $U^a_\mathcal{A}$ are all assertions of the shape $C(\alpha)$ such that $\alpha = b_0 \cdots b_{n-1}r_{n-1}b_n \in \text{ind}(\mathcal{A})$ and $C(b_n) \in \mathcal{A}$. The role assertions in $U^r_\mathcal{A}$ are all assertions of the shape $r(b_0 \cdots b_{n-1}, \alpha)$ such that $\alpha = b_0 \cdots b_{n-1}r_{n-1}b_n \in \text{ind}(\mathcal{A})$.

The following is standard to prove:

**Proposition 25** Let $\mathcal{T}$ be a Horn-\textsc{ACHTIF} TBox, $\mathcal{A}$ an ABox, and $a \in \text{ind}(\mathcal{A})$. If $\mathcal{A}$ is consistent with $\mathcal{T}$, then so is $U^a_\mathcal{A}$.

**B.2 Proof of Lemma 8**

We reformulate the lemma to make its statement more explicit.

**Lemma 8 reformulated equivalently.** Let $T_1, T_2$ be Horn-\textsc{ACHTIF} TBoxes with $T_1 \models_{\Sigma, \tau} T_2$. If $T_1 \not\models_{\Sigma, \tau} T_2$, then there is a tree-shaped $\Gamma$-ABox $\mathcal{A}$ consistent with $T_1$ and $T_2$, and a tree-shaped $\Sigma$-CQ $q$ such that one of the following holds:

1. $q$ has a single answer variable and there is an $a \in \text{ind}(\mathcal{A})$ such that $T_2, \mathcal{A} \models q(a)$ but $T_1, \mathcal{A} \not\models q(a)$;
2. $q$ is Boolean and $T_2, \mathcal{A} \models q$ but $T_1, \mathcal{A} \not\models q$.

If $T_1 \not\models_{\Sigma, \tau} T_2$, then there is a tree-shaped $\Gamma$-ABox $\mathcal{A}$ and a tree-shaped $\Sigma$-CQ $q$ with (1).

**Proof. Unrestricted CQs.** Assume $T_1 \not\models_{\Sigma, \tau} T_2$, i.e., $T_2, \mathcal{A} \models q(\alpha)$ and $T_1, \mathcal{A} \not\models q(\alpha)$, for some $\Gamma$-ABox $\mathcal{A}$ consistent with both $T_1$, some $\Sigma$-CQ $q$ and some tuple $\alpha$. Lemma 19 yields $I_{T_2, \mathcal{A}} \models q(\alpha)$ and $I_{T_1, \mathcal{A}} \not\models q(\alpha)$. We first show that the following properties of $q$ and $\alpha$ are without loss of generality:

- (a) Any match of $q(x)$ into $I_{T_2, \mathcal{A}}$ maps every quantified variable into the anonymous part.
- (b) $q(x)$ does not contain atoms of the form $r(x_1, x_2)$ with $x_1, x_2$ answer variables.
- (c) If $x = (x_1, \ldots, x_n)$ and $\alpha = (a_1, \ldots, a_n)$, then $a_i \neq a_j$ for all $i, j$ with $1 \leq i < j \leq n$.
- (d) $q(x)$ is connected.

For (a), take a match $\pi$ of $q$ in $I_{T_2, \mathcal{A}}$ and a quantified variable $y$ such that $\pi(y) = b \in \text{ind}(\mathcal{A})$. Obtain $q'(x, y)$ from $q(x)$ by removing the quantification over $y$, thus making $y$ an answer variable. Clearly, we have $I_{T_2, \mathcal{A}} \models q'(a, b)$ and $I_{T_1, \mathcal{A}} \not\models q'(a, b)$, and thus $T_2, \mathcal{A} \models q'(a, b)$ and $T_1, \mathcal{A} \not\models q'(a, b)$.

For (b), observe that such atoms can always be dropped, since they cannot be inferred via $T_1$ or $T_2$: Let $q(x) = \exists y(y(r(x_1, x_2) \land \varphi(x', y))$ with $x_1, x_2 \in x$, and let $I_{T_2, \mathcal{A}} \models q(a)$ be witnessed by the match $\pi$ with $\pi(x_1) = a_i, i = 1, 2$. Construct the CQ $q'(x', y) = \exists y(r(x_1, x_2) \land \varphi(x', y))$ by dropping the atom $r(x_1, x_2)$ (and thus possibly removing $x_1$ and/or $x_2$ from the free variables). It is clear that $I_{T_2, \mathcal{A}} \models q'(a')$ for the corresponding restriction $a'$ of the tuple $a$; thus it suffices to show that $I_{T_1, \mathcal{A}} \not\models q'(a')$.

From $I_{T_2, \mathcal{A}} \models q(a)$ we can conclude that $(a_1, a_2) \in r^+_{T_2, \mathcal{A}}$. By construction of $I_{T_2, \mathcal{A}}$ there is some $\Gamma$-role $r'$ with $r'(a_1, a_2) \in A$ and $T_2 \models r' \sqsubseteq r$ (which includes the possibility $r' = r$, i.e., $r(a_1, a_2) \in A$). Due to $T_1 \models_{\Sigma, \tau} T_2$, we also have $T_1 \models r' \sqsubseteq r$ and hence $(a_1, a_2) \in r^+_{T_1, \mathcal{A}}$. This implies the desired $I_{T_1, \mathcal{A}} \not\models q'(a')$ because, otherwise, any match $q'$ of $q$ in $I_{T_1, \mathcal{A}}$ with $\pi(x_1) = a_i, i = 1, 2$, could be extended to a match of $q$.

This construction does not introduce any violations of (a).

For (c), observe that, whenever $a_i = a_j$ for some $i, j$ with $1 \leq i < j \leq n$, we can always drop $x_i$ and $x_j$: Let $x'$ and $x'$ be $x$ and $x$ with $x_i$ and $x_j$ removed, and transform $q(x)$ into $q'(x')$ by replacing every occurrence of $x_i$ with $x_j$. Now $T_2, A \models q'(x')$ and $T_1, A \not\models q'(x')$. This construction does not introduce any violations of (a) or (b).

For (d), observe that $T_2, \mathcal{A} \models q(\alpha)$ and $T_1, \mathcal{A} \not\models q(\alpha)$ implies $T_2, \mathcal{A} \models q'(\alpha)$ and $T_1, \mathcal{A} \not\models q'(\alpha)$ for some connected component $q'(\alpha)$ of $q$. This construction does not introduce any violations of (a) or (b). While this is easy to see for (a) and (b), Property (a) requires a closer look: If the possibly disconnected CQ $q$ satisfies (a) and has at least one match $\pi$ in $I_{T_1, \mathcal{A}}$, then every match of any connected component $q'$ in $I_{T_1, \mathcal{A}}$ can be extended to a match of $q$ in $I_{T_2, \mathcal{A}}$ via $\pi$ restricted to the remaining connected components. Since the match of $q$ satisfies (a), so does the match of $q'$.

Thus, as long as $q$ violates any of the above properties, we apply the corresponding modification as described and in the order given. From now on, we assume that $q$ satisfies properties (a) to (d). Furthermore, they imply:

- (e) $q(x)$ does not contain a proper path between any two answer variables, which is a non-empty sequence of atoms $r_1(z_1, z_2), r_2(z_2, z_3), \ldots, r_n(z_n, z_{n+1})$ with variables $z_1, z_{n+1} \in x$ and $z_i \notin y$ for $1 \leq i \leq n$, and with roles $r_i$ such that $z_{i+1} \neq z_{i-1}$ for every $1 \leq i \leq n$.

To show this, assume the opposite, i.e., $q(x)$ contains a proper path as above between two answer variables $x, x'$. By (f), we have $n > 1$. By (h) and (i), $\pi$ maps all $z_i$ with $1 < i \leq n$ to the anonymous part of $I_{T_2, \mathcal{A}}$. However, there is no corresponding proper path between any two ABox individuals in $I_{T_2, \mathcal{A}}$: a contradiction.

Assume now that $q(x) = \exists y \varphi(x, y)$ is not tree-shaped, i.e., there is a cycle $r_1(z_1, z_2), r_2(z_2, z_3), \ldots, r_n(z_n, z_{n+1})$ with variables $z_i \in x \cup y$, $z_1 = z_{n+1}$, and roles $r_i$ such that $z_{i+1} \neq z_{i-1}$ for every $1 \leq i \leq n$ and $z_2 \neq z_n$. By (c), we have $z_i \in y$ for all $1 \leq i \leq n + 1$. Let $\pi$ be a match of $q(x)$ in $I_{T_2, \mathcal{A}}$. By (c), $\pi$ maps all variables to the anonymous part of $I_{T_2, \mathcal{A}}$ which, by construction, is acyclic. Hence $\pi$ cannot satisfy the properties of a match; contradiction.

Assume now that $x$ in $q(x)$ contains more than one answer variable, say $x \neq x'$, matched by $a$ and $a'$ in $I_{T_2, \mathcal{A}}$, with $a \neq a'$ due to (c). By (d), $q$ is connected, and thus, there is a path from $x$ to $x'$ in $q$. Since $x \neq x'$, there is even a proper path; contradicting (c).

Thus, now we have that $q$ is tree-shaped and behaves as required by (1) or (2). It remains to transform $\mathcal{A}$ into a tree-shaped ABox: In case $q$ is Boolean, we get from (c) and (d) that every match of $q$ in $I_{T_2, \mathcal{A}}$ is into the anonymous subtree rooted
at some ABox individual a; in case q has one answer variable, let a = a. Consider the unraveling $U_\alpha^1$ of A at a. Clearly, $T_2, U_\alpha^1 \models q(a)$ and $T_1, U_\alpha^1 \not\models q(a)$, which is still consistent with both $T_2$. By compactness, there is a finite subset $B \subseteq U_\alpha^1$ with $T_1, B \models q(a)$ and $T_2, B \not\models q(a)$. Clearly, we can also assume that $B$ is connected.

**stCQs.** Since stCQs are already tree-shaped and have exactly one answer variable, the previous argument for unrestricted CQs reduces to observing Properties (a) and (c) and unraveling the witness ABox as described. \qed

### B.3 Proof of Lemma 9

**Lemma 9** Let $T_1$ and $T_2$ be Horn-ALCHIFF TBoxes with $T_1 \models_{T_1} T_2$. Then $T_1 \models_{T_1} T_2$ iff for all tree-shaped $\Gamma$-ABoxes $A$ consistent with $T_1$ and $T_2$, $I_{T_2, A} \models_{T_1} I_{T_2, A}$.

**Proof.** We prove both implications via contraposition.

("$\Rightarrow$". Assume $T_1 \not\models_{T_1} T_2$ and consider a witness $(A, q, a)$. By Lemma 8 we can assume that $A$ is tree-shaped. From Lemma 9 (4) we get $I_{T_2, A} \models q(a)$ and $I_{T_1, A} \not\models q(a)$. If we take the finite subinterpretation $I$ of $I_{T_2, A}$ given by a match of $q$ in $I_{T_2, A}$, then we must have $I \not\models q(a)$ because of $I_{T_1, A} \not\models q(a)$. Hence $I_{T_2, A} \not\models_{T_1} I_{T_1, A}$.

("$\Leftarrow$". Assume $I_{T_2, A} \not\models_{T_1} I_{T_1, A}$, i.e., there is a finite subinterpretation $I$ of $I_{T_2, A}$ with $I \not\models_{T_1} I_{T_1, A}$. Let $a$ be the ABox individuals in $I$ and let $q_I$ be $I$ viewed as a CQ whose variables correspond to the domain elements of $I$ and the ABox individuals are represented by answer variables. Then it can be verified that $I_{T_2, A} \models q_I(a)$ and $I_{T_1, A} \not\models q_I(a)$. \qed

### B.4 Proof of Theorem 11

To prove the second part of Theorem 11 (the stCQ case), we need a bounded variant of simulations, analogous to bounded homomorphisms. We write $I_1 \subseteq^* I_2$ if for any subinterpretation $I'$ of $I_1$ with $|I'| \leq n$, we have $I'_1 \subseteq^* I_2$. Moreover, we write $I_1 \subseteq^* I_2$ if $I_1 \subseteq^* I_2$ for any $n$.

We begin with two useful insights about bounded homomorphisms (and simulations) and their connection to unbounded ones. We use $I_1^{\|d}$ to denote the restriction of $I_1$ to elements that can be reached by starting at $d$ and traveling along at most $n$ role edges (forwards or backwards).

**Fact 26** Let $\Sigma$ be a signature and $I_1, I_2$ be interpretations such that $I_2$ is finitely branching.

1. The following are equivalent.
   a. $I_1 \rightarrow_{\Sigma} I_2$
   b. For every $d \in \Delta I_1$ and every $i \geq 0$ : $I_1[d]^i \rightarrow_{\Sigma} I_2$

2. The following are equivalent.
   a. $I_1 \subseteq^* I_2$
   b. For every $d \in \Delta I_1$ and every $i \geq 0$ : $I_1[d]^i \subseteq I_2$

We will thus use Conditions (1b) and (2b) as alternative characterizations of bounded homomorphisms and simulations.

The second insight shows that, under additional conditions, we can extract an unbounded homomorphism from a suitable family of bounded ones.

**Lemma 27** Let $I_1, I_2$ be finitely branching interpretations and let $I_1$ be $\Sigma$-connected.

1. If there are $d_0 \in \Delta I_1$ and $e_0 \in \Delta I_2$ such that for each $i \geq 0$ there is a $\Sigma$-homomorphism $h_i$ from $I_1[d_0]^i$ to $I_2$ with $h_i(d_0) = e_0$, then $I_1 \rightarrow_{\Sigma} I_2$.

2. If there are $d_0 \in \Delta I_1$ and $e_0 \in \Delta I_2$ such that for each $i \geq 0$ there is a $\Sigma$-simulation $\rho_i$ of $I_1[d_0]^i$ in $I_2$ with $(d_0, e_0) \in \rho_i$, then $I_1 \subseteq I_2$.

**Proof.** We only show (1); Part (2) is analogous. We are going to construct a $\Sigma$-homomorphism $h$ from $I_1$ to $I_2$ step by step, obtaining the desired homomorphism in the limit. We will take care that, at all times, the domain of $h$ is finite and

(*) there is a sequence $h_0, h_1, \ldots$ with $h_i : \Sigma$-homomorphism from $I_1[d_0]^i$ to $I_2$ such that whenever $h(d)$ is already defined, then $h_i(d) = h(d)$ for all $i \geq 0$.

Start with setting $h(d_0) = e_0$. The original sequence $h_0, h_1$ from the lemma witnesses (*). Now consider the set $\Lambda$ that consists of all elements $d \in \Delta I_1$ such that $h(d)$ is undefined and there is an $e \in \Delta I_2$ with $h(e)$ defined and such that $d$ is reachable from $e$ along a $\Sigma$-role edge. Since the domain of $h$ is finite and $I_1$ is finitely branching, $\Lambda$ is finite. By (*), since every $d \in \Lambda$ is reachable in one step from an element $e$ such that $h(e)$ is defined, and since $I_2$ is finitely branching, for each $d \in \Lambda$ there are only finitely many $e$ such that $h_i(d) = e'$ for some $i$. Thus there must be a function $\delta : \Lambda \rightarrow \Delta I_2$ such that, for infinitely many $i$, we have $h_i(d) = \delta(d)$ for all $d \in \Delta I_2$. Extend $h$ accordingly, that is, set $h(d) = \delta(d)$ for all $d \in \Lambda$. Clearly, the sequence $h_0, h_1, \ldots$ from (*), after the extension is no longer sufficient to witness (*). We fix this by skipping homomorphisms that do not respect $\delta$, that is, define a new sequence $h_0', h_1', \ldots$ by using as $h_i'$ the restriction of $h_i$ to the domain of $I_1[d_0]^i$, where $j \geq i$ is smallest such that $h_i(d) = \delta(d)$ for all $d \in \Delta I_2$. This finishes the construction. Note that we will automatically have $h(a) = a$ for all individual names $a$ (as required), no matter whether $d_0$ is an individual name or not.

We are now ready to prove Theorem 11.

**Theorem 11** Let $T_1$ and $T_2$ be Horn-ALCHIFF TBoxes with $T_1 \models_{T_1} T_2$. Then $T_1 \models_{T_1} T_2$ iff for all tree-shaped $\Gamma$-ABoxes $A$ consistent with $T_1$ and $T_2$, and for all tree-shaped, finitely branching models $I_1$ of $(T_1, A)$, the following conditions are satisfied:

1. $I_{T_2, A} \models_{\Sigma} I_{T_1, A}$.

2. For all $\Sigma$-subtrees $I$ in $I_{T_2, A}$, one of the following holds:
   a. $I \rightarrow_{\Sigma} I_{T_1, A}$.
   b. $I \rightarrow_{\Sigma} I_{T_1, \{d_1\}(a)}$ for some $a \in \text{ind}(A)$.

Furthermore, $T_1 \models_{\text{stCQ}} T_2$ iff $I_{T_2, A} \models_{\Sigma} I_{T_1, A}$ for all $A$ and $I_1$ as above iff $I_{T_2, A} \models_{\Sigma} I_{T_1, A}$.
Proof. Unrestricted CQs, “if”. We show the contrapositive. Thus, first assume that $\Gamma \not\models_{\Sigma}^\tilde{T}_1 T_2$. By Lemma 8 there is a tree-shaped $\Gamma$-ABox $A$ consistent with both $T_1$, and a tree-shaped $\Sigma$-CQ $q$ such that either

(1') $q$ has a single answer variable and there is an element $a \in \text{ind}(A)$ such that $T_2, A \models q(a)$ but $T_1, A \not\models q(a)$

or

(2') $q$ is Boolean and $T_2, A \models q$ but $T_1, A \not\models q$.

In case (1') holds, $q$ is connected. Let $h$ be a match of $q$ in $I_{T_1, A}$; in particular $h(x) = a$. Since $q$ contains an answer variable, we must have $I_{T_2, A} \models^\con_\Sigma^I I_{T_1, A}$ as otherwise the composition of $h$ and the witnessing homomorphism shows $I_{T_1, A} \models q(a)$, which is not the case. Thus Condition (1) is violated for $I_1 = I_{T_1, A}$.

In case (2') holds, consider again a match $h$ of $q$ in $I_{T_1, A}$. Let $I_{T_2, A}$ be the restriction of $I_{T_2, A}$ to the elements in the range of $h$. Clearly, we have $I_{T_2, A} \models^\con_\Sigma I_{T_1, A}$. Consequently, $I_{T_2, A} \models^\con_\Sigma_0 I_{T_1, A}$ where $n$ is the number of variables in $q$, implying that Conditions (2) and (2b) are both false.

Unrestricted CQs, “only if”. Assume that $T_1 \models_{\Gamma, \Sigma}^\con T_2$ and let $A$ be a tree-shaped $\Gamma$-ABox consistent with both $T_1$. We first show the following:

Claim. For all models $I_1$ of $(T_1, A)$, we have $I_{T_2, A} \not\models_{\Sigma}^f I_1$.

Proof of claim: Assume to the contrary that $I_{T_2, A} \models^f_\Sigma I_1$. Then $I_{T_2, A} \models^\con_\Sigma I_1$ for some $n$, that is, there is a subinterpretation of $I$ of $I_{T_2, A}$ with $|\Delta I| \leq n$ such that $I \models^\con_\Sigma I_1$. Let $a$ be the ABox individuals in $I$ and let $q_I$ be $I$ viewed as a CQ whose variables correspond to the domain elements of $I$ and the ABox individuals are represented by answer variables. Then it can be verified that $I_{T_2, A} \models q_I(a)$ and $I_1 \not\models q_I(a)$.

Condition 1 is a consequence of Lemma 27. Fix a tree-shaped, finitely branching model $I_1 = (T_1, A)$ and let $I_{T_2, A}$ be the disjoint union of the connected interpretations $I_1, \ldots, I_k$. In each $I_i$, we find at least one individual $a_i$, $i \not\models^\con_\Sigma I_1$. We can assume w.l.o.g. that $I$ is connected and contains at least one ABox individual (otherwise we just extend $I$ and increase $n$ accordingly). Let $a$ be the ABox individuals in $I$ and let $q_I$ be $I$ viewed as a tree-shaped CQ whose variables correspond to the domain elements of $I$ and the ABox individuals are represented by answer variables. Clearly $I \models q_I(a)$ and thus $I_{T_2, A} \models q_I(a)$.

Now for Condition 2. Let $I$ be a $\Sigma$-subtree in $I_{T_2, A}$ with root $d_0$. By the claim above and Fact 26 there is a sequence $h_0, h_1, \ldots$ such that $h_i$ is a $\Sigma$-homomorphism from $I_{T_2, A}^d_{d_0}$ to $T_1, A$.

First assume that there is an $e_0 \in \Delta I_{T_2, A}$ such that $h_i(d_0) = e_0$ for infinitely many $i$. Construct a new sequence $h'_0, h'_1, \ldots$ with $h'_i$ a $\Sigma$-homomorphism from $I_{T_2, A}^d_{d_0}$ to $T_1, A$ by skipping homomorphisms that do not map $d_0$ to $e_0$, that is, $h'_i$ is the restriction of $h_i$ to the domain of $I_{T_2, A}^d_{d_0}$ where $j \geq i$ is smallest such that $h_j(d_0) = e_0$. Clearly, $h'_i(d_0) = e_0$ for all $i$. Thus, Lemma 27 yields $I \models \Sigma I_{T_2, A}$ and thus, by Lemma 19 (2) $I \models \Sigma I_1$ for every tree-shaped, finitely branching model $I_1 = (T_1, A)$.

It remains to deal with the case that there is no $e_0 \in \Delta I_{T_2, A}$ such that $h_i(d_0) = e_0$ for infinitely many $i$. We can assume that there is an $a_0 \in \text{ind}(A)$ such that $h_i(d_0) \in \Delta I_{T_2, A}$ for all $i$; in fact, there must be an $a_0$ such that $h_i(d_0) \in \Delta I_{T_2, A}$ for infinitely many $i$ and we can again skip homomorphisms to achieve this for all $i$. It is important to note that the remaining homomorphisms do not necessarily map all ancestors of $d_0$ in $I$ to elements in $I_{T_2, A}$, due to the presence of inverse roles. Now, since $I_{T_2, A}$ is finitely branching, for all $i, n \geq 0$ we must find a $j \geq i$ such that $h_j(d_0)$ is a domain element whose distance from $a_0$ exceeds $n$ (otherwise the previous case would apply). We can use this fact to construct a sequence $h'_0, h'_1, \ldots$ with $h'_i$ a $\Sigma$-homomorphism from $I_{T_2, A}^d_{d_0}$ to $I_{T_2, A}^d_{a_0}$, it is easy to verify that this implies $I \models \Sigma I_{T_2, A}$, in fact, $h'_0, h'_1, \ldots$ can again be found by skipping homomorphisms.

If we now fix an arbitrary (tree-shaped, finitely branching) model $I_1 = (T_1, A)$, by Lemma 19 (2) and (3) we have $tp_{I_{T_1, A}}(a_0) \subseteq tp_{I_1}(a_0)$ and thus $I_{T_2, A} \models^\con_\Sigma I_{T_2, tp_{I_1}(a_0)}$ as required.

stCQs. We need to show that the following three conditions are equivalent.

(i) $T_1 \models_{\Gamma, \Sigma}^\con T_2$

(ii) $I_{T_2, A} \models^\con_\Sigma I_1$ for all tree-shaped $\Gamma$-ABoxes $A$ consistent with $T_1$ and $T_2$, and for all tree-shaped, finitely branching models $I_1$ of $(T_1, A)$.

(iii) $I_{T_2, A} \models^\con_\Sigma I_{T_2, A}$ for all $A$ as above.

Proof of claim: Assume to the contrary that $I_{T_2, A} \models^\con_\Sigma I_1$. Then $I_{T_2, A} \models I_1$ for some $n$, that is, there is a subinterpretation of $I$ of $I_{T_2, A}$ such that $I \models I_1$. Let $a$ be the ABox individuals in $I$ and let $q_I$ be $I$ viewed as a CQ whose variables correspond to the domain elements of $I$ and the ABox individuals are represented by answer variables. Then it can be verified that $I_{T_2, A} \models q_I(a)$ and $I_1 \not\models q_I(a)$.

(iii) $I_{T_2, A} \models^\con_\Sigma I_{T_2, A}$ for all $A$ as above.

Proof of claim: Assume to the contrary that $I_{T_2, A} \models^\con_\Sigma I_1$. Then $I_{T_2, A} \models I_1$ for some $n$, that is, there is a subinterpretation of $I$ of $I_{T_2, A}$ such that $I \models I_1$. We can assume w.l.o.g. that $I$ is connected and contains at least one ABox individual (otherwise we just extend $I$ and increase $n$ accordingly). Let $a$ be the ABox individuals in $I$ and let $q_I$ be $I$ viewed as a tree-shaped CQ whose variables correspond to the domain elements of $I$ and the ABox individuals are represented by answer variables. Clearly $I \models q_I(a)$ and thus $I_{T_2, A} \models q_I(a)$.

We now show that the following three conditions are equivalent.

(iii) $I_{T_2, A} \models^\con_\Sigma I_{T_2, A}$ for all $A$ as above.

Proof of claim: Assume to the contrary that $I_{T_2, A} \models^\con_\Sigma I_1$. Then $I_{T_2, A} \models I_1$ for some $n$, that is, there is a subinterpretation of $I$ of $I_{T_2, A}$ such that $I \models I_1$. We can assume w.l.o.g. that $I$ is connected and contains at least one ABox individual (otherwise we just extend $I$ and increase $n$ accordingly). Let $a$ be the ABox individuals in $I$ and let $q_I$ be $I$ viewed as a tree-shaped CQ whose variables correspond to the domain elements of $I$ and the ABox individuals are represented by answer variables. Clearly $I \models q_I(a)$ and thus $I_{T_2, A} \models q_I(a)$.

We now show that the following three conditions are equivalent.

Proof of claim: Assume to the contrary that $I_{T_2, A} \models^\con_\Sigma I_1$. Then $I_{T_2, A} \models I_1$ for some $n$, that is, there is a subinterpretation of $I$ of $I_{T_2, A}$ such that $I \models I_1$. We can assume w.l.o.g. that $I$ is connected and contains at least one ABox individual (otherwise we just extend $I$ and increase $n$ accordingly). Let $a$ be the ABox individuals in $I$ and let $q_I$ be $I$ viewed as a tree-shaped CQ whose variables correspond to the domain elements of $I$ and the ABox individuals are represented by answer variables. Clearly $I \models q_I(a)$ and thus $I_{T_2, A} \models q_I(a)$.
and its subtree, and redirect each of the n original atoms to its corresponding copy. Apply this step exhaustively.

The result of this transformation is an stCQ q', which still satisfies $I_{\alpha,A} \models q'(a)$. On the other hand, $I_1 \not\models q'(a)$ because, otherwise, a match $\pi'$ of $q'(x)$ in $I_1$ would give rise to a simulation of $\mathcal{I}$ in $I_1$.

Having established the claim, we proceed as follows: Let $a$ be an ABox individual in $I_{\alpha,A}^{\Sigma,\alpha}$. By the claim and Fact 26 there is a sequence $h_0, h_1, \ldots$ such that $h_i$ is a $\Sigma$-homomorphism from $I_{\alpha,A}^{\Sigma,\alpha}$ to $I_1$. Obviously $h_i(a) = a$ for all $i$. From Lemma 27 we obtain $I_{\alpha,A}^{\Sigma,\alpha} \cong \Sigma I_1$ as desired.

\[ \square \]

C Proofs for Section 4

C.1 Proof of Theorem 12

Theorem 12 Given two Horn-ALCHIIF TBoxes $\mathcal{T}_1$ and $\mathcal{T}_2$ and types $t_i$ for $\mathcal{T}_i$, $i \in \{1, 2\}$, it can be decided in time $2^{2^{|T_1|}2^{|T_2|}}$ whether $I_{\alpha,A}^{\Sigma,\alpha} \rightarrow_{\Sigma} I_{\alpha,A}^{\Sigma,\alpha} P \alpha$.

Proof. By Lemma 13 we can decide $I_{\alpha,A}^{\Sigma,\alpha} \rightarrow_{\Sigma} I_{\alpha,A}^{\Sigma,\alpha}$ by checking whether a $\mathcal{J} \in \text{can}_c(I_{\alpha,A}^{\Sigma,\alpha})$ with $I_{\alpha,A}^{\Sigma,\alpha} \rightarrow_{\Sigma} I_{\alpha,A}^{\Sigma,\alpha}$ $\mathcal{J}$. By Lemma 14 this can be done by constructing the corresponding set $\mathcal{M}$ of mosaics for $t_1$, removing all mosaics that are not good, and checking whether the remaining set $\mathcal{M}_p$ contains a mosaic $M$ with $t_2 \in T_M$.

The desired upper bound is now a consequence of the following observations:

- The size of each 1-neighborhood in $I_{\alpha,A}^{\Sigma,\alpha}$ is bounded by $q(|T_1|)$, for a polynomial $q$.
- The number of mosaics for $t_1$ is bounded by $2^{q(|T_1|)2^{|T_2|}}$ for a polynomial $q'$; there are at most $2^{|T_1|^2}$ many 1-neighborhoods in $I_{\alpha,A}^{\Sigma,\alpha}$, and each such neighborhood admits at most $2^{|T_1|^2}2^{|T_2|}$ many decorations with sets of types.
- Given a tuple $(t', \rho, t, s, \ell)$, one can decide in time $2^q(|T_1|)$, $q'$ a polynomial, whether $(t', \rho, t, s)$ is a 1-neighborhood. Moreover, we can decide in time $2^q(|T_1|)$, $q'$ a polynomial, whether $\mathcal{M}$ is satisfied.
- Conditions 1 and 2 of a mosaic being good can be checked in the desired time.

\[ \square \]

C.2 Proof of Lemma 13

Lemma 13 Let $\mathcal{T}$ be a Horn-ALCHIIF TBox, $t_0 \in \text{tp}(\mathcal{T})$, and $\mathcal{I}$ a tree-shaped interpretation. Then $\mathcal{I} \rightarrow_{\Sigma} I_{\alpha,A}^{\Sigma,\alpha}$ if there is a $\mathcal{J} \in \text{can}_c(I_{\alpha,A}^{\Sigma,\alpha})$ with $\mathcal{I} \rightarrow_{\Sigma} \mathcal{J}$.

Proof. \textit{\(\Rightarrow\)} Let $d_0$ be the root of $\mathcal{I}$. By Fact 26 there is a sequence $h_0, h_1, \ldots$ such that $h_i$ is a $\Sigma$-homomorphism from $I_{\alpha,A}^{\Sigma,\alpha}$ to $I_{\alpha,A}^{\Sigma,\alpha}$. Note that the set $\text{tp}(\mathcal{T})$ is finite, and that $I_{\alpha,A}^{\Sigma,\alpha}$ is finitely branching. By skipping homomorphisms, we can thus construct a new sequence $h'_0, h'_1, \ldots$ such that $h'_i$ is a $\Sigma$-homomorphism from $I_{\alpha,A}^{\Sigma,\alpha}$ to $I_{\alpha,A}^{\Sigma,\alpha}$, and, additionally, for every $0 \leq i \leq j$ and $d \in \Delta I_{\alpha,A}^{\Sigma,\alpha}$ the following properties hold:

\[
\begin{align*}
(i) & \quad n_{I_{\alpha,A}^{\Sigma,\alpha}}(h'_i(d)) = n_{I_{\alpha,A}^{\Sigma,\alpha}}(h'_j(d)), \\
(ii) & \quad \text{if } e \text{ is a successor of } d \text{ in } \mathcal{I}, \text{ then } h'_i(e) \text{ is a successor of } h'_j(d) \text{ in } I_{\alpha,A}^{\Sigma,\alpha}.
\end{align*}
\]

Guided by $h'_i$, we construct a sequence of interpretations $\mathcal{J}_0, \mathcal{J}_1, \ldots$ and a sequence $g_0, g_1, \ldots$ with $g_i$ a $\Sigma$-homomorphism from $I_{\alpha,A}^{\Sigma,\alpha}$ to $\mathcal{J}_i$ such that for all $0 \leq i \leq j$ and $d \in \Delta I_{\alpha,A}^{\Sigma,\alpha}$, we have $g_i(d) = g_j(d)$. Throughout the construction, we maintain the invariant

\[
n_{I_{\alpha,A}^{\Sigma,\alpha}}(g_i(d)) \subseteq n_{I_{\alpha,A}^{\Sigma,\alpha}}(h'_i(d))
\]

for all $i, d$ such that $g_i(d)$ is defined.

We start with $\mathcal{J}_0 = \{(e_0), \cdot \}$ such that $\text{tp}_{\mathcal{J}_0}(e_0) = \text{tp}_{I_{\alpha,A}^{\Sigma,\alpha}}(h'_0(d_0))$ and $g_0(d_0) = e_0$. Clearly (*) is satisfied. Assuming that $\mathcal{J}_i$ and $g_i$ are already defined, we extend them to $\mathcal{J}_{i+1}$ and $g_{i+1}$ by doing the following for every $(d, d') \in \rho^2$ with $d \in \Delta I_{\alpha,A}^{\Sigma,\alpha}$ and $d' \notin \Delta I_{\alpha,A}^{\Sigma,\alpha}$.

- By the claim and Item (i), we have $n_{I_{\alpha,A}^{\Sigma,\alpha}}(g_i(d)) \subseteq n_{I_{\alpha,A}^{\Sigma,\alpha}}(h'_i(d))$ for all $j \geq i$; thus, we can apply (R) to $g_i(d)$ and $h'_i(d)$. More precisely, we obtain $\mathcal{J}_{i+1}$ by adding a predecessor and/or successors to achieve $n_{I_{\alpha,A}^{\Sigma,\alpha}}(g_i(d)) = n_{I_{\alpha,A}^{\Sigma,\alpha}}(h'_i(d))$. 

To define $g_i(d')$, we distinguish two cases according to Item (ii):

- $h'_i(d')$ is a successor of $h'_i(d)$ for all $j \geq i$. Then there is some $(\rho', t')$ in component $S$ of $n_{I_{\alpha,A}^{\Sigma,\alpha}}(h'_i(d))$ such that $(h'_i(d), h'_i(d')) \subseteq \rho^2 \rho \rho \rho$ and $\text{tp}_{\mathcal{J}_{i+1}}(h'_i(d')) = t'$. By (**) that pair is also in component $S$ of $n_{I_{\alpha,A}^{\Sigma,\alpha}}(g_i(d))$. Take a corresponding $\rho'$-successor $e'$ of $e$ in $\mathcal{J}_{i+1}$ and set $g_{i+1}(d') = e'$. Clearly (*) is satisfied.

- $h'_i(d)$ is a successor of $h'_i(d')$ for all $j \geq i$. Then $t' = \text{tp}_{\mathcal{J}_{i+1}}(h'_i(d))$ and $\rho$ is maximal with $(h'_i(d), h'_i(d')) \subseteq \rho^2 \rho \rho \rho$. By (**), the $t'$- and $\rho$-component in $n_{I_{\alpha,A}^{\Sigma,\alpha}}(g_i(d))$ are identical. Take a corresponding $\rho$-predecessor $e'$ of $e'$ in $\mathcal{J}_{i+1}$ and set $g_{i+1}(d') = e'$. Clearly (*) is satisfied.

The construction of $\mathcal{J}$ and $h$ is finished by setting $h = \bigcup_{i \geq 0} g_i$ and $\mathcal{J}' = \bigcup_{i \geq 0} \mathcal{J}_i$, and defining $\mathcal{J}$ as the result of exhaustive application of rule (R) to $\mathcal{J}'$.

\[ \leftarrow \rightarrow \]

It suffices to show $\mathcal{I} \rightarrow_{\Sigma} I_{\alpha,A}^{\Sigma,\alpha}$ To this end, denote with $\mathcal{J}_i, i \geq 0$, the finite submodel of $\mathcal{J}$ obtained after i rule applications, and with $d_i$ the root of $\mathcal{J}_i$. We verify the following claim, which implies $\mathcal{I} \rightarrow_{\Sigma} I_{\alpha,A}^{\Sigma,\alpha}$.

Claim. For all $i \geq 0$, we have:

- (i) there is an $e_0 \in \Delta I_{\alpha,A}^{\Sigma,\alpha}$ with $\text{tp}_{\mathcal{J}_{i+1}}(e_0) = \text{tp}_{\mathcal{J}_i}(d_i)$;
- (ii) for all $e_0 \in \Delta I_{\alpha,A}^{\Sigma,\alpha}$ with $\text{tp}_{\mathcal{J}_{i+1}}(e_0) \subseteq \text{tp}_{\mathcal{J}_i}(d_i)$, we have $(\mathcal{J}_i, d_i) \rightarrow (\mathcal{J}_{i+1}, e_0)$.

\footnote{We write $\mathcal{I} \rightarrow_{\Sigma} \mathcal{J}$ to denote that, for every $n \geq 0$, there are $n$-bounded homomorphisms from $\mathcal{I}$ to $\mathcal{J}$, without restricting the signature.}
We prove the claim by induction on $i$. For $i = 0$, Points (i) and (ii) are clear by definition of $J_0$. For the inductive step, consider $J_{i+1}$ and suppose $(R)$ has been applied to some $d \in \Delta^{J_i}$ and $e \in \Delta^{\Sigma_{T_{i+1}}}$.

Observe that Point (i) is trivially preserved when $d$ is not the root of $J_i$. In case $d = d_i$, it is preserved by the condition on the choice of $e$ in $(R)$: $e$ has the same type as $d_i$ and, by construction, the predecessor $e'$ of $e$ (if it exists) has the same type as $d_{i+1}$.

For Point (ii), we distinguish two cases:

- The extension of $J_i$ to $J_{i+1}$ has added any predecessor to $d$. In particular, we then have $d_{i+1} = d_i$. Let $e_0$ be as in (ii), i.e., $t_{J_i} = t_{J_i} \ni e_0$.
- By induction hypothesis, there is a homomorphism $h : (J_i, d_{i+1}) \rightarrow (T_{i+1}, e_0)$. We extend $h$ to the domain of $J_{i+1}$ by doing the following for each newly added successor $d'$ of $d$.

Let $t_{J_{i+1}}(d) = t$ and $t_{J_{i+1}}(d') = t'$ and $\rho$ maximal with $(d, d') \in \rho^{T_{i+1}}$. By the choice of $e$ in $(R)$, $e$ is of type $\rho$ and has a $\rho$-successor of type $t'$. By construction of the universal model, there is $r \in \rho$ with $t \sim_{\rho} t'$ and $\rho = \{s \mid T \models r \subseteq s\}$. Denote with $\hat{t} = t_{J_{i+1}}(h(d))$. The definition of a homomorphism yields $t \subseteq \hat{t}$. Thus, there is $\hat{t}' \supseteq t'$ such that $t \sim_{\rho} \hat{t}'$. By definition of the universal model, $h(d)$ has a $\rho$-successor of type $t'$ or a $\rho$-predecessor of type $\hat{t}'$, for $\hat{t} \supseteq \hat{t}'$. We extend $h$ by setting $h(d')$ to that predecessor or successor, respectively.

- The extension of $J_i$ to $J_{i+1}$ has added a $\rho$-predecessor $d'$ to $d$. Then $d = d_i$ and $d' = d_{i+1}$. Let $t_{J_{i+1}}(d) = t$ and $t_{J_{i+1}}(d') = t'$. By construction of the universal model, there is $r \in \rho$ with $t \sim_{\rho} t'$ and $\rho = \{s \mid T \models r \subseteq s\}$. Let $e_0$ be as in (ii), that is, $\hat{t}' := t_{J_i}(e_0) \ni t'$. We then have that $\hat{t}' \sim_{\rho} t$ for some $i \in \rho$. By definition of the universal model, $e_0$ has a $\rho$-successor of type $t$ or a $\rho$-predecessor of type $t_i$. Let this element be $\overline{\tau}_0$. By induction hypothesis, there is a homomorphism $h : (J_i, d) \rightarrow (T_{i+1}, \overline{\tau}_0)$. We extend $h$ by first setting $h(d) = e_0$ and then extending $h$ to all successors of $d$ as in the previous case.

It should be clear that $h$, updated as above, witnesses $(J_{i+1}, d_{i+1}) \rightarrow (T_{i+1}, e_0)$.

C.3 Proof of Lemma [14]

**Lemma** [14] Let $t_i \in t_{\mu}(T_i)$ for $i \in \{1, 2\}$. Then there is a $J \in \text{can}_{\omega}(T_1, t_1)$ such that $\mathbb{T}_{\Sigma, T_2}^{\text{can}} \rightarrow \Sigma J$ iff $M_p$ contains a mosaic $M$ with $t_2 \in \epsilon_M(t_M)$.

**Proof.** “⇒”. Let $h$ be a $\Sigma$-homomorphism from $\mathbb{T}_{\Sigma, T_2}^{\text{can}}$ to some $J \in \text{can}_{\omega}(T_1, t_1)$. For every $d \in \Delta^{\Sigma J}$, denote with $T_{h}(d)$ the set of all types mapped to $d$ by $h$, that is,

$$T_{h}(d) = \{ \text{tp}_{\Sigma, T_2}^{\text{can}}(e) \mid h(e) = d, e \in \Delta^{\Sigma_{T_2}, e^{\text{can}}} \}.$$

For every element $d \in \Delta^{\Sigma J}$, we define a tuple $M(d) = (t^-, \rho, t, S, \ell)$ by taking:

- $(t^-, \rho, t, S) = n_1^J(d)$;
- $\ell(t) = T_h(d)$;
- If there is a predecessor $d'$ of $d$, then $\ell(t^-) = T_h(d')$; otherwise, set $\ell(t^-) = \emptyset$ (not important);
- For every successor $d'$ of $d$ with $tp_{\Sigma, T_2}(d') = t'$ and $\rho = \{r \mid (d, d') \in r^s\}$ add $(\rho', t') \in S$ and set $\ell(\rho', t') = T_h(d')$;

It is easy to verify that every $M(d) = (t^-, \rho, t, S, \ell)$ obtained in this way is actually a mosaic: By definition of $J$, we know that $(t^-, \rho, t, S) = n_1^J(d')$ for some $d' \in \Delta^{\Sigma_{T_2}, e^{\text{can}}}$. Moreover, by definition of the universal model $\mathbb{T}_{\Sigma, T_2}$ and the fact that $h$ is a homomorphism, Condition (M) is satisfied.

Let $\mathcal{M}(J) = \{ M(d) \mid d \in \Delta^{\Sigma J} \}$.

“⇐”. Assume $M_p$ contains a mosaic $M$ with $t_2 \in \epsilon_M(t_M)$. We define the interpretation $J$ as the limit of the following process. We maintain a partial function $q : \Delta^{\Sigma} \rightarrow M_p$, intuitively mapping each domain element of $J$ to the mosaic that gave rise to it. Throughout the construction, the following invariant is preserved:

If $q(d) = (t^-, \rho, t, S, \ell)$, then $n_1^J(d) = (t^-, \rho, t, S, \ell)$. (*)

We start with defining $J$ as the interpretation corresponding to the 1-neighborhood represented by $M$, and define $q(e_0) = M$, where $e_0$ is the “center” of that 1-neighborhood. By definition, the invariant (*) is satisfied. Then extend $J$ by applying the following step exhaustively in a fair way: Choose some $d \in J$ such that $q(d)$ is undefined, and:

- If $d$ has a predecessor $d'$ such that $q(d') = M'$ then, due to (*), there is $(\rho, t) \in S_{M'}$ such that $(d', d) \in \rho$ and $t_{\mu}(d) = t$. Let $N \subseteq M_p$ be the mosaic that exists according to Condition 1 of being good for $(\rho, t) \in S_{M'}$. Then extend $J$ such that $n_1^J(d) = (t_N, \rho_N, t_N, S_N)$ and set $q(d) = N$.
- If $d$ has a successor $d'$ such that $q(d') = M'$ then, due to (*), we know that $\overline{T}_M \ni T_{\mu}(d') \not\subseteq \Delta$. Let $N \subseteq M_p$ be the mosaic that exists according to Condition 2 of being good. Then extend $J$ such that $n_1^J(d) = (t_N, \rho_N, t_N, S_N)$ and set $q(d) = N$.

It is immediate from the construction that these steps preserve (*), and that always one of the cases applies. Moreover, by construction, any interpretation $J$ obtained in the limit of such a process is an element of $\text{can}_{\omega}(T_1, t_1)$. It thus remains to construct a $\Sigma$-homomorphism $h$ witnessing $\mathbb{T}_{\Sigma, T_2}^{\text{can}} \rightarrow \Sigma J$. We proceed again inductively, maintaining the invariant:

If $h(d)$ is defined, then $tp_{\Sigma, T_2}(d) \in \ell_{q(h(d))}(t_{q(h(d)))}$.

Let $d_0$ be the root of $\mathbb{T}_{\Sigma, T_2}$. We start with setting $h(d_0) = e_0$, where $e_0$ is as above. By the assumption that $t_2 \in \epsilon_M(t_M)$, invariant (*) is satisfied. Now, exhaustively apply the following step. Choose $d \in \Delta^{\Sigma_{T_2}, e^{\text{can}}}$ such that $h(d)$ is not defined but $h(d') = e$ is defined for the predecessor $d'$ of $d$. Let
\(t = t_{p_1, p_2}(d), t' = t_{p_1, p_2}(d'), M' = q(d').\) By definition of \(T_{r_1, r_2}\), we know that \(t' \sim_{r_2} T_{r_1, r_2}\) for some \(r \in \text{rol}(T_2)\). Let \(\sigma = \{s \mid T \models s \subseteq s\}\). By invariant \(\{\}\), we know that \(t' \in \mathcal{M}(\{\})\). Thus, one of (a)–(c) in Condition (M) applies. Since \(d, d' \in \Delta^2 r_1 \rightarrow r_2\), we know that \(\sigma \not\subseteq \emptyset\), thus only (b) or (c) are possible. In case of (b), we extend \(h\) by setting \(h(d)\) to the predecessor of \(h(d')\). In case of (c), we extend \(h\) by setting \(h(d)\) to the according successor of \(h(d')\). Note that \(h\) extended like this satisfies the homomorphism conditions and preserves \(\mathfrak{A}\) due to the conditions in (b) and (c).

\[\Box\]

C.4 Proof of Theorem 15

We first make precise the semantics of 2ATA\(_\infty\). Let \((T, L)\) be a \(\Theta\)-labeled tree and \(\mathfrak{A} \equiv (Q, \Theta, q_0, \delta, \Omega)\) a 2ATA\(_\infty\). A run of \(\mathfrak{A}\) over \((T, L)\) is a \(T \times Q\)-labeled tree \((T_r, r)\) such that \(\varepsilon \in T_r\), \(r(\varepsilon) = (\varepsilon, q_0)\), and for all \(y \in T_r\), with \(r(y) = (x, q)\) and \(\delta(q, V(x)) = \emptyset\), there is an assignment \(v\) of truth values to the transition atoms in \(\Theta\) such that \(v\) satisfies \(\Theta\) and:

- if \(v(q') = 1\), then \(r(y') = (x, q')\) for some successor \(y'\) of \(y\) in \(T_r\);
- if \(v(\neg q') = 1\), then \(x \not\varepsilon\) and \(r(y') = (x, q')\) for some successor \(y'\) of \(y\) in \(T_r\);
- if \(v(q') = 1\), then \(x = \varepsilon\) or \(r(y') = (x, q')\) for some successor \(y'\) of \(y\) in \(T_r\);
- if \(v(\neg q') = 1\), then there are pairwise different
\(i_1, \ldots, i_n\), such that, for each \(y'\) of \(y\) in \(T_r\), with \(r(y') = (x, i_j, q')\);
- if \(v(\square q') = 1\), then for all but \(n\) successors \(x'\) of \(x\), there is a successor \(y'\) of \(y\) in \(T_r\), with \(r(y') = (x', q')\).

Let \(\gamma = i_0 i_1 \cdots\) be an infinite path in \(T_r\), and denote, for all \(j \geq 0\), with \(q_j\) the state such that \(r(i_j) = (x, q_j)\). The path \(\gamma\) is accepting if the largest number \(m\) such that \(\Omega(q_i) = m\) for infinitely many \(i\) is even. A run \((T_r, r)\) is accepting, if all infinite paths in \(T_r\) are accepting. \(\mathfrak{A}\) accepts a tree if \(\mathfrak{A}\) has an accepting run over it.

Theorem 15 The emptiness problem for 2ATA\(_\infty\) can be solved in time exponential in the number of states.

The proof is by reduction to the emptiness problem of standard two-way alternating tree automata on trees of fixed outdegree [Vardi, 1998]. We need to introduce strategy trees similar to [Vardi, 1998] Section 4. A strategy tree for \(\mathfrak{A}\) is a tree \((T, \tau)\) where \(\tau\) labels every node in \(T\) with a subset \(\tau(x) \subseteq Q \times Q \cup \{-1\} \times Q\), that is, with a graph with nodes from \(Q\) and edges labeled with natural numbers or \(-1\). Intuitively, \((q, i, p) \in \tau(x)\) expresses that, if we reached node \(x\) in state \(q\), then we should send a copy of the automaton in state \(p\) to \(x, i\). For each label \(\zeta\), we define \(\text{state}(\zeta) = \{q \mid (q, i, q') \in \zeta\}\), that is, the set of sources in the graph \(\zeta\). A strategy tree is on an input tree \((T', L')\) if \(T = T', q_0 \in \text{state}(\tau(\varepsilon))\), and for every \(x \in T\), the following conditions are satisfied:

1. if \((q, i, p) \in \tau(x)\), then \(x, i \in T\);
2. if \((q, i, p) \in \tau(x)\), then \(p \in \text{state}(\tau(x, i))\);
3. if \(q \in \text{state}(\tau(x))\), then the truth assignment \(v_{q, x}\) defined below satisfies \(\delta(q, L(x))\):

- \(v_{q, x}(p) = 1\) iff \((q, 0, p) \in \tau(x)\);
- \(v_{q, x}((-p)) = 1\) iff \((q, -1, p) \in \tau(x)\);
- \(v_{q, x}((-p)) = 1\) iff \((x, q, -1, p) \in \tau(x)\);
- \(v_{q, x}(\diamond_n p) = 1\) iff \((q, i, p) \in \tau(x)\) for \(n\) pairwise distinct \(i \geq 1\);
- \(v(\square_n p) = 1\) iff for all but at most \(n\) values \(i \geq 1\) with \(x, i \in T\), we have \((q, i, p) \in \tau(x)\).

A path \(\beta\) in a strategy tree \((T, \tau)\) is a sequence \(\beta = (u_1, \varepsilon_q)(u_2, q_2)\cdots\) of pairs from \(T \times Q\) such that for all \(i > 0\), there is some \(c_i\) such that \((q_i, c_i, q_{i+1}) \in \tau(u_i)\) and \(u_{i+1} = u_i \cdot c_i\). Thus, \(\beta\) is obtained by moves prescribed in the strategy tree. We say that \(\beta\) is accepting if the largest number \(m\) such that \(\Omega(q_i) = m\), for infinitely many \(i\), is even. A strategy tree \((T, \tau)\) is accepting if all infinite paths in \((T, \tau)\) are accepting.

Lemma 28 A 2ATA\(_\infty\), accepts an input tree iff there is an accepting strategy tree on the input tree.

Proof. The “if”-direction is immediate: just read off an accepting run from the accepting strategy tree.

For the “only if”-direction, we observe that acceptance of an input tree can be defined in terms of a parity game between Player 1 (trying to show that the input is accepted) and Player 2 (trying to challenge that). The initial configuration is \((\varepsilon, q_0)\) and Player 1 begins. Consider a configuration \((x, q)\). Player 1 chooses a satisfying true assignment \(v\) of \(\delta(q, L(x))\). Player 2 chooses an atom \(\alpha\) with \(v_{q, x}(\alpha) = 1\) and determines the next configuration as follows:

- if \(\alpha = p\), then the next configuration is \((x, p)\);
- if \(\alpha = \neg(p)\), then the next configuration is \((x, -1, p)\) unless \(x = \varepsilon\); in this case, Player 1 loses immediately;
- if \(\alpha = \neg(p)\), then the next configuration is \((x, -1, p)\) unless \(x = \varepsilon\); in this case, Player 2 loses immediately;
- if \(\alpha = \diamond_n p\), then Player 1 selects pairwise distinct \(i_1, \ldots, i_n\) with \(x, i_j \in T\), for all \(j\) (and loses if she cannot); Player 2 then chooses some \(i_j\) and the next configuration is \((x, i_j)\);
- if \(\alpha = \square_n p\), then Player 1 selects \(n\) values \(i_1, \ldots, i_n\); Player 2 then chooses some \(\ell \neq \{i_1, \ldots, i_n\}\) such that \(x, \ell \in T\) (and loses if he cannot) and the next configuration is \((x, \ell)\).

Player 1 wins an infinite play \((x_0, q_0)(x_1, q_1)\cdots\) if the largest number \(m\) such that \(\Omega(q_i) = m\), for infinitely many \(i\), is even. It is not difficult to see that Player 1 has a winning strategy on an input tree iff \(\mathfrak{A}\) accepts the input tree.

Observe now that the defined game is a parity game and thus Player 1 has a winning strategy iff she has a memoryless winning strategy [Emerson and Jutla, 1991]. It remains to observe that a memoryless winning strategy is nothing else than an accepting strategy tree.

\[\Box\]

Lemma 29 If \(L(\mathfrak{A}) \neq \emptyset\), then there is some \((T, L) \in L(\mathfrak{A})\) such that \(T\) has outdegree at most \(n\cdot C\), where \(n\) is the number of states in \(\mathfrak{A}\) and \(C\) is the largest number in \((\text{some transition} \diamond_{n} p \text{ or } \square_{n} p \text{ in } \delta)\).
We inductively construct a tree $(T', L')$ with $T' \subseteq T$ and $L'$ the restriction of $L$ to $T'$ and an accepting strategy tree $\tau'$ on $(T', L')$. For the induction base, we start with $T' = \{ \varepsilon \}$ and $\tau'$ the empty mapping. For the inductive step, assume that $\tau'(x)$ is still undefined for some $x \in T'$, and proceed as follows:

1. For every $(q, i, p) \in \tau(x)$ with $i \in \{-1, 0\}$, add $(q, i, p) \in \tau'(x)$.
2. For every $p \in \mathcal{Q}$, define $N_p = \{ i \geq 1 \mid (q, i, p) \in \tau(x), x \cdot i \in T \}$ and let $N_p' \subseteq N_p$ be a subset of $N_p$ with precisely $\min(C|N_p|)$ elements. Then:
   - (a) for all $i \in N_p'$, add $x \cdot i \in T'$;
   - (b) for all $(q, i, p) \in \tau(x)$ with $i \in N_p'$, add $(q, i, p) \in \tau'(x)$;
   - (c) for all $q \in \text{state}(x)$ and $i \in N_p'$, add $(q, i, p) \in \tau'(x)$.

By Step 2 above, $T'$ has outdegree bounded by $|\mathcal{Q}| \cdot C$. It remains to show that $\tau'$ is an accepting strategy tree on $T'$.

We verify Conditions (i)–(iii) of a strategy tree being on an input tree. Conditions 1 follows directly from the construction. For (ii), assume that $(q, i, p) \in \tau'(x)$. By construction, there is some $q' \in \{ q \cup (q', i, p) \in \tau(x) \}$, and, by Condition (ii) $(p, j, p') \in \text{state}(x \cdot i)$. Hence, there is some $(p, j, p') \in \text{state}(x \cdot i)$. By construction, there is also some $(p, j, p') \in \text{state}(x \cdot i)$, thus $p \in \text{state}(x \cdot i)$. For Condition (iii), take any $x \in T'$ and $q \in \text{state}(x \cdot i)$. As $q \in \text{state}(x \cdot i)$, we know that the truth assignment $v_{q,x}$ defined for $\tau$ in Condition (iii) satisfies $\delta(q, L, x)$. We show that for all transitions $\alpha$ with $v_{q,x}(\alpha) = 1$, we also have $v'_{q,x}(\alpha) = 1$, where $v'_{q,x}$ is the truth assignment defined for $\tau'$. By Step 1 of the construction, this is true for all $\alpha$ of the shape $p$, $(\neg p)$, and $[\neg p]$. Let now be $\alpha = \bigwedge k p$, that is, there are $k$ pairwise distinct $i \geq 1$ such that $(q, i, p) \in \tau(x)$. By the choice of $C$, we have $|N_p'| \geq k$. We have $2C(2c), b(k, i, q', p)$, that is, for all but at most $k$ values $i \geq 1$ with $x \cdot i \in T$, we have $(q, i, p) \in \tau(x)$. By Step 2(b), this remains true for $\tau'$, hence $v'_{q,x}(\alpha) = 1$.

We finally argue that $\tau'$ is also accepting. Let $\beta = (u_1, q_1)(u_2, q_2) \cdots$ be an infinite path in $(T', \tau')$. We construct an infinite path $\beta' = (u'_1, q_1)(u'_2, q_2)(u'_3, q_3) \cdots$ in $(T, \tau)$ as follows:

- $u'_1 = u_1$;
- Let $u_{i+1} = u_i \cdot \ell$ for some $\ell$ with $(q_i, \ell, q_{i+1}) \in \tau'(x)$. If $\ell \in \{0, 1\}$, we have $(q_i, \ell, q_{i+1}) \in \tau(x)$, by Step 1. We set $u'_{i+1} = u'_i \cdot \ell$. If $\ell \geq 0$ then, by Step 2(c), there is some $\ell'$ with $(q_i, \ell', q_{i+1}) \in \tau(x)$ and $x \cdot \ell' \in T'$. Set $u'_{i+1} = u'_i \cdot \ell'$.

Since every infinite path in $(T, \tau)$ is accepting, so is $\beta'$, and thus $\beta$.

We now reduce to reduce the emptiness problem of $2\mathcal{A}_{2\mathcal{A}}$ to the emptiness of alternating automata running on trees of fixed outdegree [Vardi, 1998], which we recall here. A tree $T$ is $k$-ary if every node has exactly $k$. A two-way alternating tree automaton over $k$-ary trees $(2\mathcal{A}_{2\mathcal{A}}^k)$ that are $\Theta$-labeled is a tuple $\mathcal{A} = (Q, \Theta, q_0, \delta, \Omega)$ where $Q$ is a finite set of states, $\Theta$ is the input alphabet, $q_0 \in Q$ is an initial state, $\delta$ is the transition function, and $\Omega : Q \rightarrow \mathbb{N}$ is a priority function. The transition function maps a state $q$ and some input letter $\theta$ to a transition condition $\delta(q, \theta)$, which is a positive Boolean formula over the truth constants true, false, and transitions of the form $(i, q) \in [k] \times Q$ such that $S$ satisfies $\varphi$ and for $1 \leq i \leq n$, we have $x \cdot i \in T_r$. $\omega \cdot i$ is defined, and $\tau_r(x \cdot i) = (w \cdot i, q_0)$. Accepting runs and accepted trees are defined as for $2\mathcal{A}_{2\mathcal{A}}$.

The emptiness problem for $2\mathcal{A}_{2\mathcal{A}}$'s can be solved in time exponential in the number of states [Vardi, 1998].
C.5 Proof of Theorem 16

Theorem 16 In Horn-ALCHIFY, the following problems can be decided in time $2^{o(|T|)}$, p a polynomial: (T, $\Sigma$)-CQ entailment, inseparability, and conservative extensions. The same holds for $\Sigma$-deductive entailment in ECHIFY.

We show the following lemma which, together with Theorem 15 and Lemma 7, implies Theorem 16.

Lemma 30 There are 2ATA, $\mathfrak{A}_1, \mathfrak{A}_2, \mathfrak{A}_3, \mathfrak{A}_4, \mathfrak{A}_4'$ such that:
- $\mathfrak{A}_1$ accepts $(T, L)$ iff $A$ is finite, tree-shaped, and contains $\varepsilon$;
- $\mathfrak{A}_2$ accepts $(T, L)$ iff $T_1$ is a model of $A$ and $T_2$;
- $\mathfrak{A}_3$ accepts $(T, L)$ iff $A$ is consistent with $T_2$, and $T_3$ is $T_2.A$ restricted to $ind(A)$;
- $\mathfrak{A}_4$ accepts $(T, L)$ iff either (1) or (2) from Theorem 12 is not satisfied, when $T_2.A$ is replaced with $T_2$;
- $\mathfrak{A}_4'$ accepts $(T, L)$ iff $T_2|\varepsilon \models \not\exists_T \exists \varepsilon$.

The number of states of $\mathfrak{A}_1$ and $\mathfrak{A}_2$ is polynomial in $|T_1|$ (and independent of $T_2$); the number of states of $\mathfrak{A}_3$ is polynomial in $|T_2|$ (and independent of $T_1$), and the number of states of $\mathfrak{A}_4, \mathfrak{A}_4'$ is exponential in $|T_2|$ (and independent of $T_1$). An automata can be constructed in time polynomial in $|T_1|$ and double-exponential in $|T_2|$.

The construction of the automaton $\mathfrak{A}_1$ is straightforward, so we concentrate on $\mathfrak{A}_2$, $\mathfrak{A}_3$, and $\mathfrak{A}_4$.

In what follows, we use $\Diamond q$ and $\Box q$ to abbreviate $\Diamond_1 q$ and $\Box_0 q$, respectively. We define $\mathfrak{A}_2 = \langle Q_2, \Theta, q_0, \delta_2, \Omega_2 \rangle$ where

$$Q_2 = \{q_0, q_A\} \cup \{q_\alpha \mid \alpha \in T_1\} \cup \{q_\rho, \eta_\rho \mid \rho \in \Theta_1\} \cup \{q_r, q_{r,B}, q_{r, B}^+, q_{r, B}^- | \exists r. B \in cl(T_1)\},$$

and $\Omega_2$ assigns 0 to all states. The idea of $\mathfrak{A}_2$ is to check that the ABox is satisfied, realized in state $q_A$, and that every axiom TBox axiom in $T_1$ is satisfied everywhere, realized using states $q_\alpha$ below. Formally, the transition function $\delta_2$ is given as follows, for $\sigma = (L_0, L_1, L_2)$:

$$\delta_2(q_0, \sigma) = \Box q_0 \land q_A \land \bigwedge_{\alpha \in T_1} q_\alpha$$

$$\delta_2(q_A, \sigma) = \bigwedge_{\rho \in \Lambda_0} q_\rho$$

$$\delta_2(q_{\text{hme}}(r), \sigma) = (q_{r^-} \land \Box q_r) \lor (\Box q_{r^-} \land \Diamond q_r)$$

$$\delta_2(q_{r, s}, \sigma) = \eta_r \lor q_s$$

$$\delta_2(q_{A,N A, B}, \sigma) = \eta_A \lor q_{A, B}$$

$$\delta_2(q_{A, A, B}, \sigma) = \eta_A$$

$$\delta_2(q_{T, A}, \sigma) = q_A$$

$$\delta_2(q_{A, r, B}, \sigma) = \eta_A \lor q_r \lor q_B$$

$$\delta_2(q_{r, A, B}, \sigma) = \eta_r \lor q_A \lor q_B$$

$$\delta_2(q_{r, B}, \sigma) = \Box q_r \lor (q_{r^-} \land \neg q_B)$$

$$\delta_2(q_{r, B}^+, \sigma) = \bigwedge q_{r, B}^+ \lor (q_{r^-} \land \neg q_B)$$

$$\delta_2(q_{r, B}^-, \sigma) = q_r \land q_B$$

$$\delta_2(q_{r, B}^-, \sigma) = q_r \land q_B$$

Finally, we set for all $\rho \in \Theta_1$:

$$\delta_2(q_\rho, \sigma) = \begin{cases} \text{true} & \text{if } \rho \in L_1 \\ \text{false} & \text{if } \rho \notin L_1 \end{cases}$$

$$\delta_2(q_\rho, \sigma) = \begin{cases} \text{true} & \text{if } \rho \notin L_1 \\ \text{false} & \text{if } \rho \in L_1 \end{cases}$$

Automaton $\mathfrak{A}_4$ relies on a syntactic characterization of ABox entailment [Bienvenu et al., 2013], which we introduce first.

Let $\mathcal{T}$ be a Horn-ALCHIFY TBox and $A$ a tree-shaped ABox. A derivation tree for an assertion $A_0(a_0)$ in $A$ w.r.t. $\mathcal{T}$ with $A_0 \in \mathcal{N}_C$ is a finite $\text{ind}(A) \times \mathcal{N}_C$-labeled tree $(T, V)$ that satisfies the following conditions:

1. $V(\varepsilon) = (a_0, A_0)$;
2. if $V(x) = (a, A)$ and neither $A(a) \notin A$ nor $\top \subseteq A$, then one of the following holds:
   (i) $x$ has successors $y_1, \ldots, y_k$, $k \geq 1$ with $V(y_i) = (a, A_i)$ for $1 \leq i \leq k$ and $T \models A_1 \cap \cdots \cap A_k \subseteq A$;
   (ii) $x$ has a single successor $y$ with $V(y) = (b, B)$ and there is an $\exists r. B \subseteq A \in T$ and an $(s, a, b) \in A$ such that $T \models s \subseteq r$;
   (iii) $x$ has a single successor $y$ with $V(y) = (b, B)$ and there is a $B \subseteq \exists r. A \in T$ such that $r(b, a) \in A$ and $\text{fun}(r) \in T$.

Note that the first item of Point 2 above requires $T \models A_1 \cap \cdots \cap A_n \subseteq A$ instead of $A_1 \cap \cdots \cap A_n \in T$ to ‘shortcut’ anonymous parts of the universal model. In fact, the derivation of $A$ from $A_1 \cap \cdots \cap A_n$ by $T$ can involve the introduction of anonymous elements.

The main property of derivation trees is the following.

Lemma 31 Let $\mathcal{T}$ be a Horn-ALCHIFY TBox and $A$ an ABox consistent with $\mathcal{T}$. Then for all assertions $A(a)$ with $A \in \mathcal{N}_C$ and $a \in \text{ind}(A)$ we have $T, A \models A(a)$ iff there is a derivation tree for $A(a)$ in $\mathcal{A}$ w.r.t. $\mathcal{T}$.

Proof. The “if”-direction is immediate, so we concentrate on the “only if”-direction. We construct a sequence of interpretations $\mathcal{I}_0, \mathcal{I}_1, \ldots$ by the following procedure. We start with setting:

$$\Delta_{\mathcal{I}_0} = \text{ind}(A)$$

$$\mathcal{A}_{\mathcal{I}_0} = \{a \mid A(a) \in A\}$$

$$r_{\mathcal{I}_0} = \{(a, b) \mid a \in A \text{ and } \text{fun}(r) \in T\}$$

For every $i \geq 0$, we obtain $\mathcal{I}_{i+1}$ from $\mathcal{I}_i$ by setting $\mathcal{I}_{i+1} = \mathcal{I}_i$ and applying the following rules to all $d, e \in \Delta_{\mathcal{I}_i}$:

1. If $d \in (A_1 \cap A_2)^+, \text{ but } d \notin A^{\mathcal{I}_i}$ for some $A_1 \cap A_2 \subseteq A \in T$, then add $d \in A^{\mathcal{I}_{i+1}}$;
2. If $d \in (\exists r. B)^{\mathcal{I}_i}$, but $d \notin A^{\mathcal{I}_i}$ for some $\exists r. B \subseteq A \in T$, then add $d \in A^{\mathcal{I}_{i+1}}$;
3. If $(d, e) \in r^{\mathcal{I}_i}$, but $(d, e) \notin s^{\mathcal{I}_i}$, for some $s$ with $T \models r \subseteq s$, then add $(d, e) \in s^{\mathcal{I}_{i+1}}$;
4. If $d \in A^{\mathcal{I}_i}$, but $d \notin (\exists r. B)^{\mathcal{I}_i}$ for some $A \subseteq \exists r. B \in T$, then...
(a) if there is e with \((d, e) \in r^T\) and func(r) ∈ T then add e ∈ B^(T)+
(b) otherwise add a fresh domain element e with \((d, e) \in r^{T+1}\) and e ∈ B^(T+1).

Let I be defined as \(A^I = \bigcup_{n \geq 0} I_n, A^2 = \bigcup_{n \geq 0} A^2\), and \(r^I = \bigcup_{n \geq 0} r^I\). It is standard to verify the following:

Claim 1. I → J for all models J of T and A.

By definition of \(I_0\), we have \(I \models A\). Moreover, we have \(I \models T'\) where \(T' \subseteq T\) is obtained from \(T\) by dropping all CLs of the form \(A \sqsubseteq \perp\) and all FAs. Since \(A\) is consistent with \(T\), there is a model \(J\) of \(A\) and \(T\); in particular, \(A^J = \emptyset\) for all \(A \sqsubseteq \perp \subseteq T\). By Claim 1, we have \(I \models J\), and thus \(A^J = \emptyset\). For the FAs func(s), observe that they are obeyed by \(A\) (because of consistency with \(T\)) and that they are preserved, by rule 4(a). Thus, I is a model of T.

Claim 2. For all i ≥ 0, we have:

(a) For all \(a \in \text{ind}(A)\): if \(a \in A^i\), then there is a derivation tree for \(a(A)\) in \(A\) w.r.t. T.

(b) If \(e\) was created because of \(d\) in Rule 4(b), then \(T \models \bigcap \{A \mid d \in A^i\} \subseteq \exists r. \{A \mid e \in A^i\}\) for all \(r\) with \((d, e) \in r^i\).

Proof of Claim 2. It is standard to show Part (b) of the Claim. We show Part (a) by induction on \(i\). By construction of \(I_0\), it is true for \(i = 0\). Consider \(I_{i+1}\). If \(a \in A^{i+1}\) because of Rule 1, construct a derivation tree of type (i) from the derivation trees for \(A^i\) and \(A^i\) which exist due to the induction hypothesis. If \(a \in A^{i+1}\) because of Rule 2, there is some \(d \in B^{i+1}\) with \((a, d) \in r^i\) and \(\exists r. B \subseteq A \subseteq T\). If \(d \in \text{ind}(A)\), then there is some \(s(a, d) \in A\) with \(T \models s \subseteq r\), by Rule 3. We can thus construct a derivation of type (ii) from the derivation tree of \(B(d)\), which exists due to induction hypothesis. If \(d \notin \text{ind}(A)\), then \(d\) was created because of \(a\) in Rule 4(b). By Part (b) of the Claim, we have \(T \models \bigcap \{A' \mid a \in A^i\} \subseteq \exists r. B\). Hence, \(T \models \bigcap \{A' \mid a \in A^i\} \subseteq A\), and we can construct a derivation tree of type (i) for \(a(A)\). If \(a \in A^{i+1}\) because of Rule 4(a), there is \((d, a) \in r^i\) and \(d \in B^{i+1}\), and \(B \subseteq \exists r. A\), func(r) ∈ T. If \(d \in \text{ind}(A)\), we can construct a derivation tree of type (ii) for \(A\) from the derivation tree of B(d) which exists by induction. If \(d \notin \text{ind}(A)\), then \(d\) was created because of \(a\) in Rule 4(b). By Part (b) of the Claim, we have \(T \models \bigcap \{A' \mid a \in A^i\} \subseteq \exists r. B\). Hence, \(T \models \bigcap \{A' \mid a \in A^i\} \subseteq A\), and construct a derivation tree of type (i) for \(a(A)\) based on this. This finishes the proof of Claim 2 and the Lemma.

In the following Lemma, we characterize consistency of ABoxes with TBoxes.

Lemma 32 Let T be a Horn-ALCH(L,F) TBox and A an ABox. Then A is consistent with T if the following points are satisfied for all \(a \in \text{ind}(A)\):

1. the following ABox \(A_a\) is consistent with T:

\[
A_a = \{ B(a) \mid B(a) \text{ has a derivation tree in } A \text{ w.r.t. } T \}.
\]

2. for all \(\text{func}(s) \in T\), there is at most one \(b \in \text{ind}(A)\) with \(s(a, b) \in A\).

Proof. The “only if”-direction is immediate, so we concentrate on the “if”-direction. Assume that all \(a \in \text{ind}(A)\) satisfy both items above. By the first item, there is a model \(I_0\) of \(A_0\) and \(T\). Since we are considering Horn-ALCH(L,F), there is also a tree-model \(A_0\) with root \(d_a \in \Delta^0\) satisfying, for all concept names \(B \in N_c\):

\[\star\] \(d_a \in B^0\) iff \(T, A_0 \models B(a)\).

We construct an interpretation \(I\) as follows. Start with \(I_0\) by taking

\[
\Delta^0 = \text{ind}(A)
\]

\[
A^0 = \{ a \mid A(a) \text{ has a derivation tree in } A \text{ w.r.t. } T \}
\]

\[
r^0 = \{ (a, b) \mid s(a, b) \in A, T \models s \subseteq r \}
\]

Now, obtain \(I\) from \(I_0\) by performing the following operation for every \(a \in \text{ind}(A)\) and \(b \in \Delta^0\) such that \((d_a, b) \in r^0\) for some set of roles \(\rho\) which contains no role \(r\) such that there is \(a'\) with \(r(a, a') \in A\). Extend \(I\) by adding the sub-interpretation of \(I_0\) rooted at \(b\) as a \(r\)-successor of \(a\).

Based on \((\star)\) and the assumptions, it is straightforward to show that \(I\) is a model of \(A\) and \(T\).

We are now ready to give the automaton \(A_3\). We take \(A_3 =\) \((Q_3, \Theta, q_0, \delta_3, \Omega_3)\) where

\[
Q_3 = \{ q_0, q_0, \} \cup \{ q_A, \bar{q}_A \mid A \in \Omega_3 \cap N_c \} \cup \{ q_r, q_r, q_r, q_r, q_r, q_r \mid r \in \Omega_3 \cap N_c \} \cup \{ q_a, q_a, q_a, b \mid r \in \Omega_3 \}
\]

and \(\Omega_3\) assigns zero to all states, except for states of the form \(q_A\), to which it assigns 1. The automaton \(A_3\) ensures that, for all \(n \in \text{ind}(A)\) we have:

(i) \(A \in L_2(n)\) iff there is a derivation tree for \(A(n)\) in \(\mathcal{A}\),

(ii) for all \(n \neq r, r \in R_2(n)\) iff there is some \(s\) such that \(s(n \cdot -1, n) \notin \mathcal{T}_2\) and \(\mathcal{T}_2 \models r \subseteq s\).

Intuitively, these points ensure that the represented interpretation \(I_2\) is the universal model of \(T_2\) and \(\mathcal{A}\), in case \(\mathcal{A}\) is consistent with \(T_2\). Having (i) and (ii), we can check inconsistency of \(\mathcal{A}\) with \(T_2\) based on Lemma 32 that is, we verify the following conditions for all \(n \in \text{ind}(\mathcal{A})\):

(iii) the set \(L_2(n) \cap N_c\) is consistent with \(T_2\);

(iv) for each \(s\) with \(\text{func}(s) \in T\), there are no \(n_1 \neq n_2\) such that both \(s(n, n_1) \in \mathcal{A}\) and \(s(n, n_2) \in \mathcal{A}\).

For Point (i), we use states \(q_A\) for the “if” part, and states \(\bar{q}_A\) for the “only if” part; for Point (ii), we use states \(q_r\) and \(\bar{q}_r\), respectively. Intuitively, a state \(q_A\) assigned to some node \(n\) is an obligation to verify the existence of a derivation tree for \(A(n)\). Conversely, \(\bar{q}_A\) is the obligation that there is no such derivation tree. Similar obligations hold for \(q_r\) and \(\bar{q}_r\). For Point (iii), we precompute the set of consistent types and check (iii) while visiting all \(n \in \text{ind}(\mathcal{A})\). Point (iv) can be checked directly on \(\mathcal{A}\), that is, independent from \(T_2\). The automaton starts with the following transitions, where we assume \(\sigma = (L_0, L_1, L_2)\):

\[
- \delta_3(q_0, \sigma) = \text{true if } L_0 = 0;
- \delta_3(q_0, \sigma) = \text{false if } L_0 \neq 0 \text{ and } L_2 \cap N_c \text{ inconsistent with } T_2, \text{c.f. Point (iii)};
\]
if \( L_0 \neq \emptyset \) and \( L_2 \cap N_C \) consistent with \( \mathcal{T}_2 \), then
\[
\delta_3(q_0, \sigma) = \Box q_0 \land q_{0r} \land \bigwedge_{A \in L_2 \cap N_C} q_A \land \bigwedge_{A \in (\mathcal{T}_2 \cap N_C) \setminus L_2} \neg q_A.
\]
- \( \delta_3(q_{0r}, \sigma) = \text{true if } L_0 = \emptyset; \)
- \( \delta_3(q_r, \sigma) = \bigwedge_{\text{func}(r) \in \mathcal{T}_2} q_r^i \land \bigwedge_{r \in L_2 \cap N_C} q_r \land \bigwedge_{r \in (\mathcal{T}_2 \cap N_C) \setminus L_2} \neg q_r. \)
- \( \delta_3(q_r^i, \sigma) = \begin{cases} \Box q_r & \text{if } r^- \in L_0 \\
\Box q_{r-} & \text{if } r^- \notin L_0 \end{cases} \)
- \( \delta_3(q_{r-}, \sigma) = \begin{cases} \text{true if } r \notin L_0 \\
\text{false otherwise} \end{cases} \)

Now, for states \( q_A \), we directly implement the conditions of a derivation tree. Finiteness of the derivation is ensured by the priority of states of the form \( q_A \). The relevant transitions are as follows:
- \( \delta_3(q_A, \sigma) = \text{false if } L_0 = \emptyset; \)
- \( \delta_3(q_A, \sigma) = \text{true if } A \in L_0; \)
- \( \delta_3(q_A, \sigma) = \bigvee_{T_2 = A_1 \cap \cdots \cap A_n \subseteq A} (q_{A_1} \land \cdots \land q_{A_n}) \lor \\
\bigvee_{\exists r. B \subseteq A \in T_2, T_2 = s \subseteq r} (q_A^r \land \neg q_B) \lor q_{s,B} \lor \\
\bigvee_{B \subseteq \exists r. A \in T_2, \text{func}(r) \in T_2} (q_A^r \land \neg q_B) \lor q_{s,-B} \lor \\
\bigvee_{\exists r. A \in T_2, \text{func}(r) \in T_2} (q_A^r \land \neg q_B) \lor q_{s,-B}. \)
- \( \delta_3(q_{r}, \sigma) = \begin{cases} \text{true if } r \in L_0; \\
\text{false otherwise} \end{cases} \)
- \( \delta_3(q_{r-}, \sigma) = q_A^r \lor q_B. \)

The transitions for \( \overline{q}_A \) are obtained by taking the “complement” of the ones for \( q_A \). More precisely, we define \( \delta_3(\overline{q}_A, \sigma) = \overline{\delta_3(q_A, \sigma)} \), where \( \overline{\sigma} \) is obtained from \( \sigma \) by exchanging \( \land \) and \( \lor \), \( \Box \) and \( \Box \), \( \neg \) and \( \neg \), and true and false, and replacing every state \( p \) with \( \overline{p} \); see the following set of transitions.
- \( \delta_3(\overline{q}_A, \sigma) = \text{true if } L_1 = \emptyset; \)
- \( \delta_3(\overline{q}_A, \sigma) = \text{false if } A \in L_1; \)
- \( \delta_3(\overline{q}_A, \sigma) = \bigvee_{T_2 = A_1 \cap \cdots \cap A_n \subseteq A} (\overline{q}_{A_1} \lor \cdots \lor \overline{q}_{A_n}) \land \\
\bigvee_{\exists r. B \subseteq A \in T_2, T_2 = s \subseteq r} (\overline{q}_A^r \land \neg \overline{q}_B) \land \overline{q}_{s,B} \land \\
\bigvee_{B \subseteq \exists r. A \in T_2, \text{func}(r) \in T_2} (\overline{q}_A^r \lor \neg \overline{q}_B) \land \overline{q}_{s,-B} \land \\
\bigvee_{B \subseteq \exists r. A \in T_2, \text{func}(r) \in T_2} (\overline{q}_A^r \lor \neg \overline{q}_B) \land \overline{q}_{s,-B}. \)
- \( \delta_3(\overline{q}_{r}, \sigma) = \begin{cases} \text{false if } r \in L_0; \\
\text{true otherwise} \end{cases} \)
- \( \delta_3(\overline{q}_{r^t}, \sigma) \neq q_A^r \lor \overline{q}_B. \)

Finally, states \( q_r \) and \( \overline{q}_r \) at some node \( n \) represent the obligation to verify that the role atom \( r(n \cdot -1, n) \) follows, respectively does not follow, from \( T \) and \( A_2 \). This is realized by the following transitions which implement Point (ii) above.
\[
\delta_3(q_r, \sigma) = \bigvee_{T_2 = s \subseteq r} q_A^s \land \delta_3(\overline{q}_r, \sigma) = \bigvee_{T_2 = s \subseteq r} \overline{q}_A^s.
\]

For the automaton \( \mathfrak{A}_4 \), we take \( \mathfrak{A}_4 = (Q_4, \Theta, q_0, \delta_4, \Omega_4) \) where
\[
Q_4 = \{ q_0, q_1, q_r \} \cup \{ q_i, q_i^3, q_i^{30} \mid t \in \text{tp}(T_2) \} \cup \\
\{ q_{r,t}, q_{r,t}^i \mid t \in \text{tp}(T_2), \rho \text{ set of } \text{sig}(T_2)\text{-roles} \},
\]
and \( \Omega_4 \) assigns zero to all states, except for states of the form \( q_i, t \in \text{tp}(T_2) \), to which it assigns one. For some \( n \in \text{ind}(A) \), denote with \( \mathcal{J}_n \) the universal model of the type \( \{ A(n) \mid A \in L_2(n) \} \) and \( \mathcal{T}_2 \). The automaton ensures that indeed (1) or (2) from Theorem 11 is not satisfied, by verifying that there is some \( n \in \text{ind}(A) \) such that one of the following conditions holds:
1. there is \( r \in \Sigma \) and \( n' \in \text{ind}(A) \) such that \( (n, n') \in r^2 \), but \( (n, n') \notin r^{21} \);
2. \( \mathcal{J}_n \not\models \mathcal{I}_1 \);
3. there is a \( \Sigma \)-subtree \( \mathcal{J} \) of \( \mathcal{J}_n \) such that
   \( \mathcal{J} \not\models \mathcal{I}_1 \), and
   \( \mathcal{J} \not\models \mathcal{I}_{1', \text{tp}_{21}(m)} \), for all \( m \) with \( L_0(m) \neq \emptyset \).

Condition 1 is straightforward (realized in state \( q_r \)). For Condition 2, we use states \( q_t \) with \( t \in \text{tp}(T_2) \). A state \( q_t \) assigned to a node \( n \) represents the obligation to verify that there is no \( \Sigma \)-homomorphism from the universal model of \( t \) and \( T_2 \) to \( \mathcal{I}_1 \) that maps the root to \( n \). This is the case if either the root cannot be mapped to \( n \), or, recursively, there is some \( \rho \)-successor \( t' \) of \( t \) in the universal model such that the universal model of \( t' \) and \( T_2 \) cannot be mapped to any \( \rho \)-neighbor of \( n \). This process is finite because of priority 1 for all \( q_t \) with \( t \in \text{tp}(T_2) \). For Condition 3, we precompute the set \( \mathcal{R}_2(t) \) of all types of roots of \( \Sigma \)-subtrees which appear in the universal model of \( t \) and \( T_2 \), and the relation \( \rightarrow^{\text{fin}}_{\Sigma} \) according to Theorem 12. Thus, the sets \( \mathcal{R}_2(t) \) and the test for finite homomorphisms can be used directly in the transition condition, see states \( q_3 \) and \( q_{l,30} \), respectively. Using states \( q_{l,30} \), the automaton ensures that a given root \( t \) of a \( \Sigma \)-subtree satisfies 3(a) and 3(b).

Let \( t \Sigma \) and \( \rho |_{\Sigma} \) denote the restriction of \( t \) and \( \rho \), respectively, to symbols from \( \Sigma \). For \( \sigma = (L_1, L_2, L_3) \), we take the
Theorem 18

In any DL between $\mathcal{ELI}$ and $\mathcal{ELHI}_F$, deductive conservative extensions, deductive $\Sigma$-entailment, and deductive $\Sigma$-inseparability are $\text{coNEXPTIME}$-hard.

The proof is by reduction of a NEXPTIME-complete tiling problem, where the aim is to tile a $2^n \times 2^n$-grid, to the complement of stCQ-conservative extensions. This tiling problem was introduced as a special case of the origin constrained domino problem by Grädel [1989], and its NEXPTIME-hardness follows from Grädel’s Theorem 3.3. An instance is given by a tuple $P = (\Sigma, \Sigma_0, H, V)$, where $\Sigma$ is a finite set of tile types, $\Sigma_0 \subseteq \Sigma$ is a set of distinguished tiles to be placed on position $(0,0)$ of the grid, and $H$ and $V$ are horizontal and vertical matching conditions. Let $|\Sigma| = n$. A solution to $P$ is a function $\tau : 2^n \times 2^n \to \Sigma$ such that

- if $\tau(i, j) = t$ and $\tau(i, j + 1) = t'$ then $(t, t') \in V$, for all $i < 2^n, j < 2^n$,
- $\tau(0, 0) \in \Sigma_0$. We can assume w.l.o.g. that for every tile $t \in \Sigma$, there is a $t'$ with $(t', t) \in V$.

Let $P = (\Sigma, \Sigma_0, H, V)$. We show how to construct $\mathcal{ELI}$ TBoxes $T_1$ and $T_2$ such that $T_1 \cup T_2$ is a $(\text{sig}(T_1), \text{sig}(T_2))$-stCQ-conservative extension of $T_1$ iff there is no solution for $P$. Hence, stCQ-conservative extensions, $(T, \Sigma)$-stCQ entailment, and $(T, \Sigma)$-stCQ inseparability are $\text{coNEXPTIME}$-hard in $\mathcal{ELI}$ (and any DL that contains it as a fragment). Since $T_1$ and $T_2$ are formulated in $\mathcal{ELI}$, we trivially have $T_1 \models_{\text{sig}(T_1)} T_1 \cup T_2$. Thus, hardness of deductive conservative extensions follows from Lemma 24 in all DLs between $\mathcal{ELI}$ and $\mathcal{ELHI}_F$ since $T_1$ and $T_2$ are formulated in $\mathcal{ELI}$ and Lemma 24 covers deductive conservative extensions in $\mathcal{ELHI}_F$. This also implies hardness of deductive $\Sigma$-entailment and of deductive $\Sigma$-inseparability in the mentioned DLs.

The intuitions and correctness proofs are based on the following characterization of stCQ-conservative extensions.

Lemma 33 Let $T_1$ and $T_2$ be $\mathcal{ELI}$ TBoxes such that all role names in $T_2$ are in $\text{sig}(T_1)$. Then $T_1 \cup T_2$ is a $(\text{sig}(T_1), \text{sig}(T_2))$-stCQ-conservative extension of $T_1$ iff $\mathcal{I}_{T_1 \cup T_2, A} \models_{\text{sig}(T_1)} \mathcal{I}_{T_1, A}$ for all tree-shaped $(\text{sig}(T_1))$-ABoxes $A$.

Proof. An interpretation is strongly tree-shaped if it is tree-shaped and does not contain multi-edges, that is, any $d, d' \in \Delta^T$ are involved in at most one role edge. Since $T_1$ and $T_2$ are formulated in $\mathcal{ELI}$ (and thus do not contain role inclusions), for any tree-shaped ABox $A$ the universal models $\mathcal{I}_{T_1 \cup T_2, A}$ and $\mathcal{I}_{T_1, A}$ are strongly tree-shaped. The assumption on role names in $T_2$ made in the lemma implies that every element in $\mathcal{I}_{T_1 \cup T_2, A}$ can be reached from an ABox individual by traveling along only $\text{sig}(T_1)$-roles. Together, this implies the following: (*) there is a $\text{sig}(T_1)$-simulation from $\mathcal{I}_{T_1 \cup T_2, A}$ to $\mathcal{I}_{T_1, A}$ iff there is a $\text{sig}(T_1)$-homomorphism from $\mathcal{I}_{T_1 \cup T_2, A}$ to $\mathcal{I}_{T_1, A}$.

From Theorem 11 we get that $T_1 \cup T_2$ is a $(\text{sig}(T_1), \text{sig}(T_2))$-stCQ-conservative extension of $T_1$ iff $\mathcal{I}_{T_1 \cup T_2, A} \models_{\text{con}(\text{sig}(T_1))} \mathcal{I}_{T_1, A}$. But $\mathcal{I}_{T_1 \cup T_2, A} \models_{\text{con}(\text{sig}(T_1))} \mathcal{I}_{T_1, A}$ by the assumption on role names in $T_2$ made in the lemma and simulations can be replaced with homomorphisms by (*).

We will build $T_1$ and $T_2$ such that the same single role name $r$ is used in $T_1$ and $T_2$, thus the assumption in Lemma 33 will be satisfied.

For a clearer presentation, we proceed in two steps. We first define $T_1$ and $T_2$ to be an $\mathcal{ELHI}$-TBox, i.e., on both sides of CIs we allow concepts of the following form:

$$L, L' ::= T | A | L \cap L' | L \cup L' | \exists r.L.$$ The only non-trivial use of disjunction will be on the right-hand side of a CI in $T_2$. In a second step, we show how to remove disjunction.

We use $S$ to abbreviate the role composition $r;r^\ast$, writing for example $\exists S.C$ for $\exists r.\exists r^\ast.C$. Note that $S$ behaves like...
We are going to enforce (\(T_1\)) to be empty (except introducing the required symbols) and \(T_2\) to verify the existence of an \(S\)-path in the input ABox whose individuals represent the grid positions along with a tiling, row by row from left to right, starting at the lower left corner and ending at the upper right corner. The positions in the grid are represented in binary by the concept names \(X_1, \ldots, X_{2n}\) in the ABox where \(X_1, \ldots, X_n\) indicate the horizontal position and \(X_{n+1}, \ldots, X_{2n}\) the vertical position. The tiling is represented by concept names \(T_t, t \in \Sigma\). The verification is done by propagating a concept name as a marker bottom up and while doing this, verifying the horizontal matching condition. Under the assumption that an additional labeling with concept names \(T'_t, t \in \Sigma\), is such that

\(\text{(+) every point in the path is labeled with } T'_t \text{ if its descendant at distance exactly } 2^\text{th} \text{ (that is, the grid position immediately below it) is labeled with } T_t,\)

the vertical matching condition is also verified.

For several reasons, this program cannot quite be implemented in the desired way. First, we still have to make sure that (\(+)\) actually holds. This is done as follows. We install yet another labeling with concept names \(T_t, t \in \Sigma\), such that a node is labeled with \(T'_t\) if it is not labeled with \(T_t\). Then \(T_2\) checks for a violation of (\(+)\) in the following way: when the propagation reaches the final individual the verified path, the generation of a finite anonymous \(S\)-path is triggered. That path homomorphically embeds into the \(S\)-path in the ABox in many different ways since \(S\) is reflexive-symmetric. In fact, for any individual on the ABox path we can find a homomorphism such that the endpoint of the anonymous path maps to that individual because (a) the anonymous path is long enough to reach the first individual on the ABox path and (b) the homomorphism can always ‘fold’ the reflexive-symmetric role \(S\) in a suitable way. At the end of the anonymous path, we then guess (using disjunction) a tile \(t\), make \(T'_t\) true, continue building the anonymous path for another \(2^n\) steps (in a way such that it cannot fold), and finally make \(T_t\) true. Let us pretend for a second that our TBoxes are formulated in \(\mathcal{ELC}\). If (\(+)\) is violated, then the guess can be made so that the anonymous path homomorphically maps into the ABox path. Otherwise, this is not the case. Clearly, the latter can occur only if \(P\) has a solution.

The fact that \(S\) is reflexive-symmetric allows the mentioned folding of the existential path. However, it poses some complications in the verification of the \(S\)-path in the ABox because we must be careful not to confuse successors with predecessors. To this end, every grid position is actually represented by three consecutive individuals labeled with the concept names \(B_0, B_1, B_2\), respectively. All these individuals are labeled identically regarding the \(X\)-counter and the concept names \(T_t\). We are going to enforce (\(+)\) for the \(B_2\)-individuals and only these individuals also receive \(T'_t\) and \(T_t\) labels (any other \(B_i\) would work as well). Another problem is that \(T_2\) cannot check all possible kinds of defects. In particular, it cannot detect the defect that an element is labeled with more than one tile or that there are multiple successors in the ABox that have an incompatible labeling with the counter concept names. We thus use \(T_1\) to check for such defects. If found, it will generate a defect of the kind that \(T_2\) cannot verify, that is, a violation of (\(+)\).

We start with assembling \(T_2\), which uses a single role name \(r\) via the abbreviation \(S\) introduced above and the following concept names.

- jointly with \(T_1\):
  - \(X_1, \ldots, X_{2n}, \overline{X}_1, \ldots, \overline{X}_{2n}\) for the binary representation of the horizontal and vertical grid positions on the ABox path
  - \(B_0, B_1, B_2\) for distinguishing successors and predecessors on the ABox path (these concept names implement a unary counter that counts modulo three)
  - \(T_t, \overline{T}_t, t \in \Sigma\), representing tile types present/not present at individuals on the ABox path
  - \(T'_t, t \in \Sigma\), representing tile types present at the descendant at distance exactly \(3 \cdot 2^n\) from the given individual on the ABox path
- additionally:
  - \(L\) as a verification marker to be propagated along the ABox path
  - \(\overline{t}_i, 1 \leq i \leq 2n\), to indicate that the incrementation of the counter values at an ABox individual is correct regarding the \(i\)-th bit (the 1st bit being that of least value)
  - \(Y_1, \ldots, Y_{2n}, \overline{Y}_1, \ldots, \overline{Y}_{2n}\) for counting the length of the anonymous path
  - \(Y'_1, \ldots, Y'_n, \overline{Y}'_1, \ldots, \overline{Y}'_n\) implement another counter on the anonymous path, used to continue extending the path by exactly \(3 \cdot 2^n\) positions to reach the grid position immediately below
  - \(M_t, t \in \Sigma\), for memorizing a tile type on the anonymous path

\(T_2\) consists of the following CIs.

1. The initial grid position starts the propagation:
\[
X_1 \sqcap \cdots \sqcap X_{2n} \sqcap B_0 \sqcap \bigcup_{t \in \Sigma} T_t \sqsubseteq L
\]
2. The verification proceeds upwards. We first verify that the counter is incremented properly when moving upwards along the \(S\)-path in the ABox:
\[
B_0 \sqcap X_i \sqcap \exists S.(B_2 \sqcap X_i) \sqcap \bigcup_{1 \leq j < i} \exists S.(B_2 \sqcap X_j) \sqsubseteq \overline{t}_i \\
B_0 \sqcap \overline{X}_i \sqcap \exists S.(B_2 \sqcap \overline{X}_i) \sqcap \bigcup_{1 \leq j < i} \exists S.(B_2 \sqcap \overline{X}_j) \sqsubseteq \overline{t}_i \\
B_0 \sqcap X_i \sqcap \exists S.(B_2 \sqcap X_i) \sqcap \bigcup_{1 \leq j < i} \exists S.(B_2 \sqcap \overline{X}_j) \sqsubseteq \overline{t}_i \\
B_0 \sqcap \overline{X}_i \sqcap \exists S.(B_2 \sqcap \overline{X}_i) \sqcap \bigcup_{1 \leq j < i} \exists S.(B_2 \sqcap \overline{X}_j) \sqsubseteq \overline{t}_i
\]
where \(i\) ranges over \(1..2n\) and \(j\) over \(\{0, 1\}\). These inclusions only work under the assumption that no individual

---

*a We will make sure that all ‘relevant domain elements’ have an \(r\)-successor, which guarantees reflexivity.*
has two $S$-neighbors that are labeled with the same $B_i$ but are labeled differently regarding $X_i$ and $X_j$ for some $j$. We shall prevent this situation later using $T_1$.

3. We next make a verification step inside a row of the grid:

$$B_0 \cap \text{ok} \cap \cdots \cap \text{ok}_n \cap T_i \cap \nabla X_i \cap \exists S.(B_2 \sqcap L \sqcap T_i) \sqsubseteq L$$

$$B_1 \cap \text{ok} \cap \cdots \cap \text{ok}_n \cap T_i \cap \nabla X_i \cap \exists S.(B_2 \sqcap L \sqcap T_i) \sqsubseteq L$$

$$B_2 \cap \text{ok} \cap \cdots \cap \text{ok}_n \cap T_i \cap \bigcap_{\ell \in [1..T_i]} T_i \cap T_i \cap \nabla X_i \cap \exists S.(B_2 \sqcap L \sqcap T_i) \sqsubseteq L$$

where $t_1, t_2, t_3$ range over $\nabla$ such that $(t_1, t_2) \in H$ and $(t_3, t_2) \in V$ and $i$ ranges over $1..n$. The use of $\nabla X_i$ on the left-hand sides ensures that we move inside a row. In the first line, we make a move between horizontally neighboring grid positions, verifying the horizontal matching condition. In the other lines, we move along the three points representing the same grid position, ensuring that they are all labeled by the same tile.

4. We also have to consider the case where we jump from one grid row to the next, ignoring the tiling condition:

$$B_0 \cap \text{ok} \cap \cdots \cap \text{ok}_n \cap T_i \cap \nabla X_i \cap \exists S.(B_2 \sqcap L) \sqsubseteq L$$

where $t$ ranges over $\nabla$ and $i$ ranges over $n + 1..2n$. The use of $\nabla X_i$ on the left-hand side ensures that we are not yet in the topmost row.

5. When the final individual of the ABox path is reached (maximum counter value and $B_2$-label), we make an extra step in the ABox to a $B_0$-labeled individual and then generate the first object of an existential path:

$$B_0 \cap \exists S.(B_2 \sqcap X_1 \sqcap \cdots \sqcap X_{2n} \sqcap L) \sqsubseteq \exists S.C$$

where $C = Y_1 \sqcap \cdots \sqcap Y_{2n} \sqcap B'_2$

The purpose of the extra step will be explained later on.

6. Here and in the following, we use the abbreviation

$$\exists S^{(3)}(X_1, X_2, X_3) := \exists S.X_1 \cap \exists S.X_2 \cap \exists S.X_3$$

for concept names $X_1, X_2, X_3$.

We continue building the path, decrementing the $Y$-counter:

$$Y_i \cap B'_2 \sqsubseteq \exists S^{(3)}(B'_1, B'_0, B'_2)$$

$$B'_0 \cap \exists S.(B'_2 \sqcap Y_i) \sqsubseteq Y_i$$

$$B'_2 \cap \exists S.(B'_1 \sqcap Y'_i) \sqsubseteq Y'_i$$

$$B'_2 \cap \exists S.(B'_1 \sqcap \bigcap_{i \leq j < i} Y'_i) \sqsubseteq Y'_i$$

$$B'_2 \cap \exists S.(B'_1 \sqcap \bigcup_{i \leq j < i} Y'_i) \sqsubseteq Y'_i$$

where $i$ ranges over $1..2n$ and $j$ over $\{0, 1\}$. It is essential to use different $B'_i$ and counter concepts $Y_i$, $Y'_i$ than in the ABox; otherwise the anonymous path could not homomorphically embed into the ABox path in a folded way. It would actually suffice to build a path of length $3 \cdot 2^{2n-1}$ because no violation of $(\ast)$ can start in the bottommost row. However, overcounting does not compromise correctness.

7. At the end of the anonymous path, we implement a violation of $(\ast)$ as described above: we guess a tile $t$ involved in the violation, make sure that $T'_i$ holds at the current point, start a new counter, travel exactly $3 \cdot 2^n$ steps (without any folding), and verify that $\overline{T}_i$ holds where we arrive:

$$\overline{Y}_1 \cap \cdots \cap \overline{Y}_{2n} \cap B'_2 \sqsubseteq \overline{Y}'_1 \cap \cdots \cap \overline{Y}'_{n} \sqcap \bigcup_{\ell \in T'_i} (T'_i \sqcap M_i)$$

$$Y'_i \sqcap B_2 \sqsubseteq \exists S^{(3)}(B_0, B_1, B_2)$$

$$B_{j+1} \sqcap \exists S.(B_j \sqcap \overline{Y}'_i) \sqsubseteq \overline{Y}'_i$$

$$B_{j+1} \sqcap \exists S.(B_j \sqcap \overline{Y}'_i) \sqsubseteq \overline{Y}'_i$$

$$B_0 \sqcap \exists S.(B_2 \sqcap \bigcap_{i \leq j < i} \overline{Y}'_i) \sqsubseteq \overline{Y}'_i$$

where $i$ ranges over $1..n$ and $t$ over $\nabla$. Here we use the same $B_i$ as in the ABox to avoid folding.

This finishes the definition of $T_2$. We now define $T_1$, which uses the following additional concept names.

- $D$ for indicating the occurrence of a defect
- $Z_1, \ldots, Z_n$ for an additional counter.

$T_1$ consists of the following CIs.

1. Tiles are mutually exclusive: for all distinct $t, t' \in \nabla$:

$$T_i \cap T_{i'} \subseteq D$$

where $D$ starts a path that implements a violation of $(\ast)$, to be implemented below.

2. The problematic situation described at the end of Item 2 in the definition of $T_2$ cannot occur:

$$\exists S.(B_k \sqcap X_i) \sqcap \exists S.(B_k \sqcap \overline{X}_i) \sqsubseteq D$$

where $k$ ranges over $\{0, 1, 2\}$ and $i$ over $1..2n$.
3. We next implement the path triggered by $D$:

\[
\begin{align*}
B_2 \cap (D \cup S.D \cup S.S.D) &\subseteq Z_1 \cap \cdots \cap Z_n \cap T_i', \\
Z_i \cap B_2 &\subseteq \exists S^{(3)}. (B_0, B_1, B_2) \\
B_{j+1} \cap \exists S. (B_j \cap Z_i) &\subseteq Z_i, \\
B_{j+1} &\subseteq \exists S. (B_j \cap Z_i), \\
B_0 \cap \exists S. (B_2 \cap \bigcap_{1 \leq j < i} Z_i) &\subseteq Z_i, \\
B_0 \cap \exists S. (B_2 \cap Z_i) &\subseteq Z_i, \\
B_0 \cap \exists S. (B_2 \cap \bigcup_{1 \leq j < i} Z_i) &\subseteq Z_i, \\
B_0 \cap \exists S. (B_2 \cap Z_i) &\subseteq Z_i, \\
Z_1 \cap \cdots \cap Z_n &\subseteq T_i,
\end{align*}
\]

where $i$ ranges over $1..n$ and $t \in \Sigma$ is fixed.

Before we eliminate the disjunction used in $T_2$, let us mention the central property of $T_1$ and $T_2$ that can be used to show correctness of the reduction. Since $T_2$ uses disjunction, a universal model for $T_2$ and an ABox $\mathcal{A}$ is not guaranteed to exist. Instead, there is a set of models for $T_2$ and $\mathcal{A}$ that is universal in the sense that for every model $\mathcal{I}$ of $T_2$ and $\mathcal{A}$, there is a model in the set that admits a homomorphism into $\mathcal{I}$. We refrain from giving a formal definition. Now, the central property of $T_1$ and $T_2$ is as follows: $P$ has a solution iff there is a tree-shaped $\mathrm{sig}((T_1))$-$\mathcal{A}$ consistent with $T_1 \cup T_2$ such that, for every $I$ in the universal set of models for $T_1 \cup T_2$ and $\mathcal{A}$, there is no $\mathrm{sig}(T_1)$-homomorphism from $I$ to $\mathcal{I}_{T_1, \mathcal{A}}$.

We now show how to get rid of disjunction. The central property of $T_1$ and $T_2$ will essentially be preserved, with a single universal model playing the role of the universal set of models. The disjunctions on the left-hand sides of Cs only serve as abbreviations and can easily be removed with only a polynomial blowup of TBox sizes. What remains is the disjunction in Item 7 of the definition of $T_2$. To get rid of it, we need to modify both $T_1$ and $T_2$:

- In the first line of Point 7 of the definition of $T_2$, the disjunction is replaced with a conjunction, generating $|\Sigma|$ many defective chains at once:

\[
\bigwedge_{t \in \Sigma} \bigvee_{t \in \Sigma} \exists S. (Y_1 \cap \cdots \cap Y_n \cap B_2 \cap T_i' \cap M_i)
\]

The resulting universal model is illustrated in Figure 1 where we assume $\Sigma = \{t_1, t_2, t_3\}$, showing the final individual on the ABox path, the extra step from Item 5 of the definition of $T_2$, the anonymous path, and the branching gadget attached to the end of it.

- We have now generated too many paths and thus the desired homomorphism may not exist even if $(\ast)$ is violated in the ABox. We compensate by enforcing in $T_1$ that when a $B_2$-individual in the ABox path is labeled with $T_i'$, then it’s $B_0$-predecessor on that path rooted at $\Sigma - 1$ many additional paths, realizing every possible violation of $(\ast)$ except the one induced by $t \in \Sigma$. This is illustrated in Figure 2 where we again assume $\Sigma = \{t_1, t_2, t_3\}$, showing a $B_2$-individual labeled with $T_i'$, and the extra successors of it’s $B_0$-predecessor generated by $T_i'$.

Note that this explains the extra step in Item 5 of the definition of $T_2$ since also the final $B_2$-element of the ABox path must have a $B_0$-predecessor.

We add the following to $T_1$, using fresh concept names $M_i', t \in \Sigma$:

\[
\begin{align*}
B_0 \cap \exists S. (B_2 \cap T_i') &\subseteq \bigcap_{1 \leq j < i} \exists S. (Y_1' \cap \cdots \cap Y_n' \cap B_2 \cap M_i') \\
B_0 \cap \exists S. (B_2 \cap Y_i') &\subseteq Y_i', \\
B_0 \cap \exists S. (B_2 \cap Y_i') &\subseteq Y_i', \\
B_0 \cap \exists S. (B_2 \cap \bigcup_{1 \leq j < i} Y_i') &\subseteq Y_i', \\
B_0 \cap \exists S. (B_2 \cap \bigcup_{1 \leq j < i} Y_i') &\subseteq Y_i', \\
B_0 \cap \exists S. (B_2 \cap Y_i') &\subseteq Y_i', \\
B_0 \cap \exists S. (B_2 \cap Y_i') &\subseteq Y_i', \\
Y_1' \cap \cdots \cap Y_n' \cap M_i' &\subseteq T_i
\end{align*}
\]

where $i$ ranges over $1..n$ and $t$ over $\Sigma$.

**Lemma 34** $T_1 \cup T_2$ is a $(\mathrm{sig}(T_1), \mathrm{sig}(T_2))$-stCQ-conservative extension of $T_1$ if there is no solution for $P$.

**Proof.** By Lemma 33, it suffices to show the following.

**Claim.** $P$ has a solution iff there is a tree-shaped $\mathrm{sig}(T_1)$-$\mathcal{A}$ such that there is no $\mathrm{sig}(T_1)$-homomorphism from $\mathcal{I}_{T_1 \cup T_2, \mathcal{A}}$ to $\mathcal{I}_{T_1, \mathcal{A}}$.

We now sketch a proof of the claim.

![Figure 1: Part of the universal model of $T_2$.](image-url)
Figure 2: Extra successors to eliminate disjunction

“⇒”. Assume that \( P \) has a solution. Let \( \mathcal{A} \) be the ABox that contains a single \( S \)-path of length \( 3 \cdot 2^n \) which correctly encodes the solution to \( P \) via the concept names \( X_i, \overline{X}_i, B_i, T_i, T'_i \). Since \( \mathcal{A} \) correctly encodes the tiling, there is no \( \text{sig}(T_1) \)-homomorphism from \( I_{T_1 \cup T_2, \mathcal{A}} \) to \( I_{T_1, \mathcal{A}} \); in particular, the anonymous path described by \( T_2 \) that ends in \( |T| \) many violations of (\( \ast \)), each represented via a path, cannot be mapped to \( I_{T_1, \mathcal{A}} \).

“⇐”. Assume \( P \) has no solution and let \( \mathcal{A} \) be a tree-shaped \( \text{sig}(T_1) \)-ABox. If \( \mathcal{A} \) contains no path of length \( 3 \cdot 2^n \) that is labeled in the desired way with the concept names \( X_i, \overline{X}_i, B_i, T_i, T'_i \) (and which does not necessarily satisfy (\( \ast \))), then the generation of the anonymous path in \( I_{T_1 \cup T_2, \mathcal{A}} \) is not triggered and the identity is a \( \text{sig}(T_1) \)-homomorphism from \( I_{T_1 \cup T_2, \mathcal{A}} \) to \( I_{T_1, \mathcal{A}} \). If there is such a path, then it violates (\( \ast \)) since \( P \) has no solution. Consequently, the anonymous path in \( I_{T_1 \cup T_2, \mathcal{A}} \) homomorphically maps to \( I_{T_1, \mathcal{A}} \), which is sufficient to show that there is a homomorphism from \( I_{T_1 \cup T_2, \mathcal{A}} \) to \( I_{T_1, \mathcal{A}} \).  \( \square \)