

Theorem Proving in Isabelle

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Try this in Isabelle

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i.e.

$$\text{Suc}(?m) = \text{Suc}(?n) \implies ?m = ?n.$$

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(first two steps by $\forall E$)

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Check it out!

Lifting

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What about resolution with

$$(\rightarrow I) \frac{P \implies Q}{P \rightarrow Q} \quad \text{or} \quad (\forall I) \frac{\bigwedge x. P}{\forall x. P} \quad ?$$

Problem: The premises contain meta-logical symbols (\implies , \bigwedge), hence do not match conclusion of any other rule!

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Meta-rule:

$$\frac{\phi \implies \psi}{(\theta \implies \phi) \implies (\theta \implies \psi)}$$

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Resolution with $\rightarrow I$, i.e. with $(?P \Longrightarrow ?Q) \Longrightarrow ?P \rightarrow ?Q$:

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Solution for $\forall I$ and the like:

- Introduce quantification over **parameter** x in premises and conclusion
- Make all unknowns $?a$ depend on x : replace by $?a(x)$.
- Meta-rule:

$$\frac{\phi \Longrightarrow \psi}{\bigwedge x. \phi^x \Longrightarrow \bigwedge x. \psi^x}$$

(ϕ^x is ϕ with parametrized unknowns)

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Want to resolve $\wedge E1$, i.e. $?P \wedge ?Q \implies ?P$, with $\forall I$, i.e. $\bigwedge x. ?P(x) \implies \forall x. ?P(x)$.

Lift $\wedge E1$ over parameter x :

$$\bigwedge x. ?P(x) \wedge ?Q(x) \implies \bigwedge x. ?P(x)$$

Resolve with $\forall I$:

$$\bigwedge x. ?P(x) \wedge ?Q(x) \implies \forall x. ?P(x)$$