

Modularity in Ontologies: Introduction (Part B)

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Plan for the rest of today's lecture

- 1 A case for modularity of ontologies
- 2 Overview and comparison of modularisation approaches
- 3 Overview of the remainder of this course



And now . . .

- 1 A case for modularity of ontologies
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What can I do with my ontology?

Ontology users and engineers use ontologies to

- represent and archive knowledge
- compute inferences from that knowledge (quickly)
e.g., classification, query answering, explanation

Modularity can help with these tasks.



What can I do with my ontology?

Building and using an ontology can be eased by

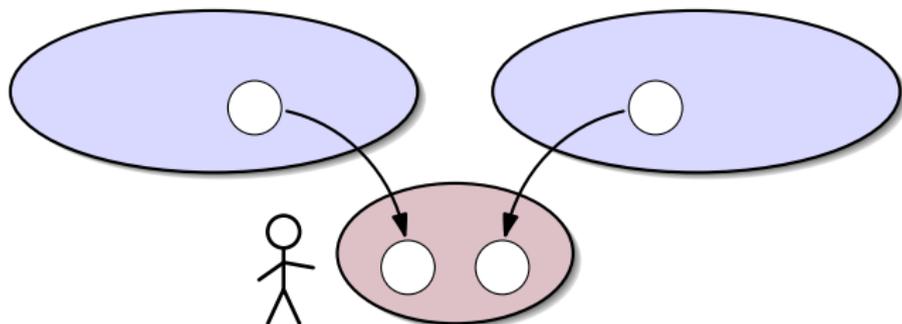
- frequent and fast reasoning
classification, explanations;
expressivity \leftrightarrow complexity, optimisations, incremental reasoning
- reusing knowledge from existing ontologies
efficient import
- exposing the logical structure of the represented knowledge
comprehension
- collaborative development
- version control

Modularity can help with these tasks.



An import/reuse scenario

“Borrow” knowledge from external ontologies



- Provides access to well-established knowledge
- Doesn't require expertise in external disciplines

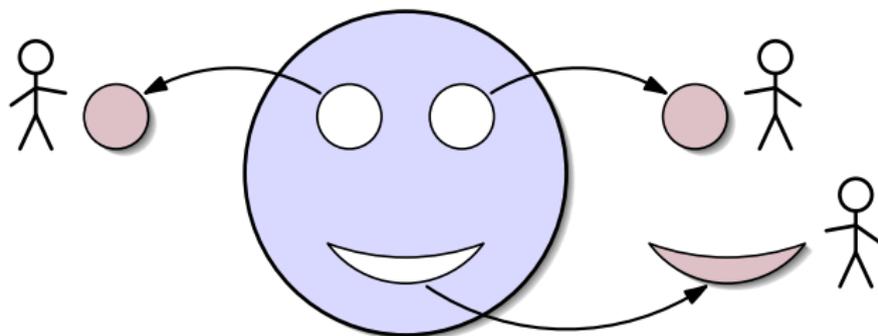
This scenario is well-understood and implemented.

⇒ *Tuesday + Wednesday*



A collaboration scenario

Collective ontology development



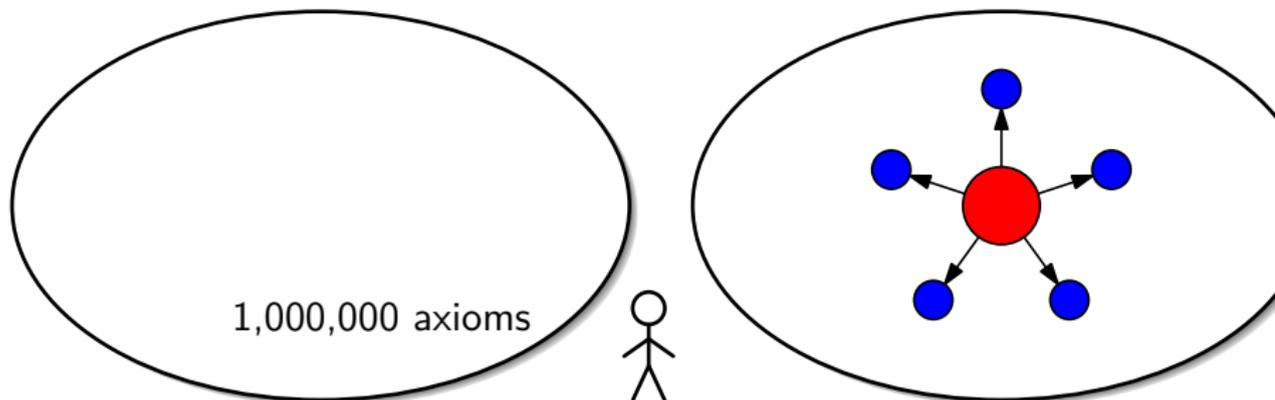
- Developers work (edit, invoke reasoning) locally
- Extra care at re-combination
- Prescriptive/analytic behaviour

This approach is mostly understood, but not implemented yet.



Understanding and/or structuring an ontology

Compute the logical structure of an ontology



This is work in progress. ➤ Thursday



And now . . .

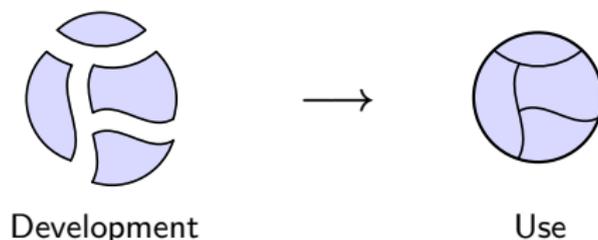
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A priori vs. a posteriori

A priori modularisation approaches

- First, a modular structure for the ontology \mathcal{O} is decided on.
- Then, \mathcal{O} is developed and used according to that structure.



A-priori modularisation approaches

- Provide a framework to develop an ontology modularly from the start
- Provide means to “bridge” between the modules
dependency of modules/signature, flow of knowledge
- Often consist of extensions of (description) logics
- Sometimes allow for distributed reasoning
- Generally, don't guarantee that modules are logically closed
in some cases, this is deliberately so



A-priori: different files with imports

- Used to develop large ontologies about different domains
- Each domain expert (team) maintains “their” file \mathcal{F}_i
- The overall ontology \mathcal{O} imports all files:
$$\mathcal{O} = \mathcal{F}_1 \cup \dots \cup \mathcal{F}_n$$
- Example: $\mathcal{F}_1, \mathcal{F}_2, \mathcal{F}_3$ about diseases, anatomy and drugs
- Problems?
 - The \mathcal{F}_i are not necessarily logically closed
 - Experts’ knowledge interferes with each other,
e.g.: diseases are located in body parts and treated by drugs
- \rightsquigarrow Maintenance of \mathcal{O} as difficult as in the monolithic case
 - Reasoning or reuse might still require the whole ontology
- Still used to develop and maintain ontologies!
see e.g. <http://bioportal.bioontology.org>



Package-based description logics (PB-DLs)

[Bao et al. 2006, 2009]

- Extension of standard DLs
- Domain-specific files are called *packages*
- Semantic import links between packages (explicit dependency)
- Terms annotated with “home package”
- Semantics local w.r.t. each package
- Reasoning controlled by the links
- Translation to “plain” DLs yields implicit decision procedures
- Problems?
 - Reasoning or reuse might still require the whole ontology



Distributed description logics

[Borgida and Serafini 2003] [Serafini and Taminin 2009]

- Similar to PB-DLs
- Replace import links by “bridge rules”:
subconcept relations between (complex) concepts from
different packages
- Distributed decision procedures exist

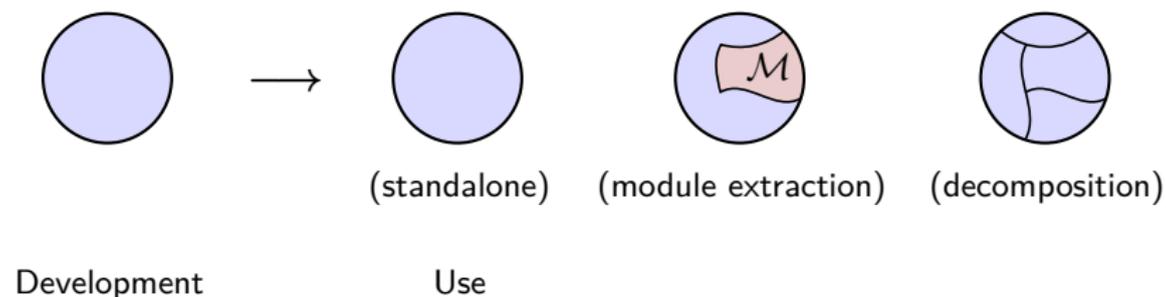
Related notion: E-connections
[Kutz et al. 2001]



A priori vs. a posteriori

A posteriori modularisation approaches

- The ontology \mathcal{O} is built and used as a monolithic entity.
- A module is extracted or \mathcal{O} is decomposed into modules.



A-posteriori modularisation approaches

- Regard an ontology \mathcal{O} as a monolithic entity
remember: \mathcal{O} is a set of axioms
- Module: subset $\mathcal{M} \subseteq \mathcal{O}$
- Extract one module (e.g., for reuse) or
decompose \mathcal{O} into several modules (e.g., for comprehension)
- Often, a **signature** (set of terms) $\Sigma \subseteq \text{sig}(\mathcal{O})$ is specified
and the module extracted using Σ as a parameter
- Ideally, modules **encapsulate** knowledge in some form
e.g., all consequences of \mathcal{O} in Σ
- Not all a-posteriori module notions guarantee encapsulation



Graph-based a-posteriori modularisation approaches

- Are based on a graph representation of the ontology
usually concept/role hierarchy, sometimes enriched with disjointness
- Start with a signature Σ
- Traverse the graph and “harvest” entities and axioms
follow subconcept relation and/or restrictions (\exists , domain, range)
- Resulting module = set of harvested axioms
- Examples
 - Ontology segmentation [Seidenberg and Rector 2006, 2009]
 - Traversals [Noy and Musen 2003, 2009]
 - More general framework [d’Aquin et al. 2007]



Pro and contra graph-based approaches

Pro

- Modules can usually be extracted efficiently
time polynomial in the size of $\mathcal{O} \rightsquigarrow$ robustly scalable
- Easy to implement
- Applicable to many logics

Contra

- Heuristic, no characterisation of the expected module contents
 - In particular, no logical guarantees such as entailment preservation
- \rightsquigarrow Modules typically lose knowledge from \mathcal{O}



A-posteriori approaches with coverage

Coverage

$\mathcal{M} \subseteq \mathcal{O}$ covers \mathcal{O} for Σ if

all Σ -consequences of \mathcal{O} already follow from \mathcal{M} .

- i.e., \mathcal{M} preserves all knowledge in \mathcal{O} about α
 \Rightarrow *Tuesday*
- This guarantee is needed, e.g., for ontology reuse or reasoning
 \Rightarrow *Tuesday + Wednesday*
- Of course, \mathcal{O} is always covering

Problems

- *Minimal* covering modules are, in general, hard to extract
 \Rightarrow *Tuesday*



Coverage-providing module notions

- Restricted to logics where coverage can be decided efficiently
e.g., MEX for acyclic \mathcal{EL} \Rightarrow *Wednesday*
[Konev et al. 2008]
- Or use a tractable condition sufficient for coverage,
leading to modules that always contain minimal modules

Examples:

- Modules obtained from partitions based on E-connections
[Cuenca Grau et al. 2006]
- Locality-based modules \Rightarrow *Wednesday*
[Cuenca Grau et al. 2007, 2009]
- Reachability-based modules \Rightarrow *Friday*
[Suntisrivaraporn 2008]



Comparison of a-posteriori module extraction approaches

Module notion	Covrg.	Min.	Covered DLs	Complexity
All axioms referencing Σ	✗		any	easy
Graph-based	✗		any	easy
The whole ontology	✓	✗✗	any	easy
Minim. mod. with coverage*	✓	✓	few	hard
MEX*	✓	✓	acyclic \mathcal{EL}	easy
E-connections based mod.	✓	✗	OWL	easy
Locality-based mod.*	✓	✗	OWL	easy
Reachability-based mod.*	✓	✗	almost OWL	easy
(Modules with rewriting)	✓?	✓✓?	few?	hard?

*Will be covered here \Rightarrow *Tuesday, Wednesday, Friday*



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Course overview

- ② Module extraction and its formal foundations
 - A case scenario: modular reuse
 - Logical guarantees required
 - Conservative extensions, inseparability, robustness
- ③ Module extraction
 - Efficient module notions (locality, MEX)
 - Module extraction algorithms and tools
- ④ Decomposing ontologies
 - Atomic decomposition
- ⑤ Recent advances/current work
 - Forgetting and interpolation
 - Reachability-based modules
 - Incremental reasoning
 - Modular reasoning

